

iip: an integratable TCP/IP stack

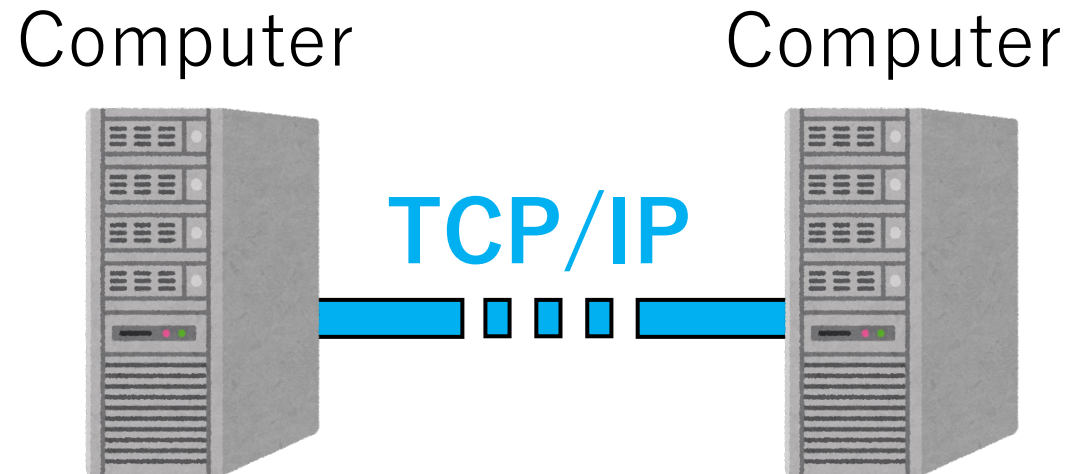
Kenichi Yasukata



Internet Initiative Japan

TCP/IP and TCP/IP Stacks

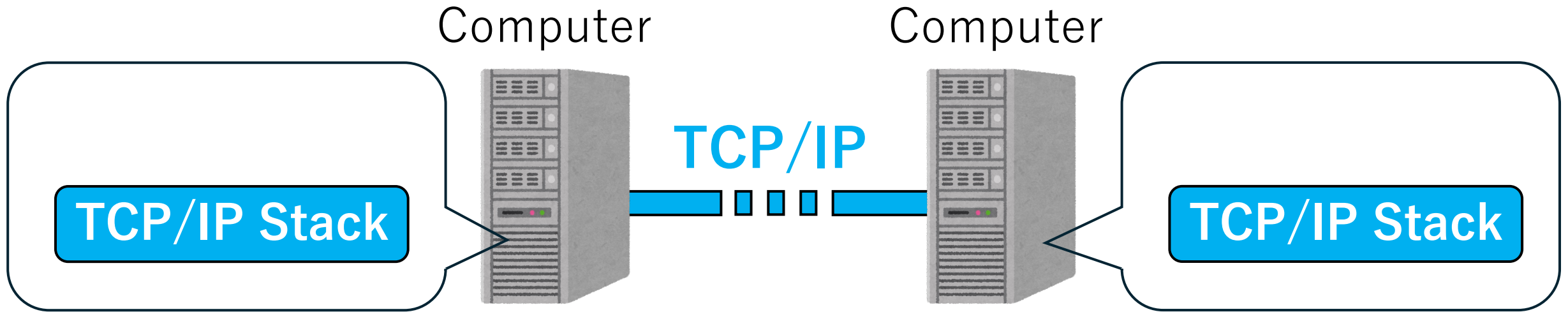
- TCP/IP is a standardized protocol suite commonly used for communication in computer networks



<https://github.com/yasukata/iip>

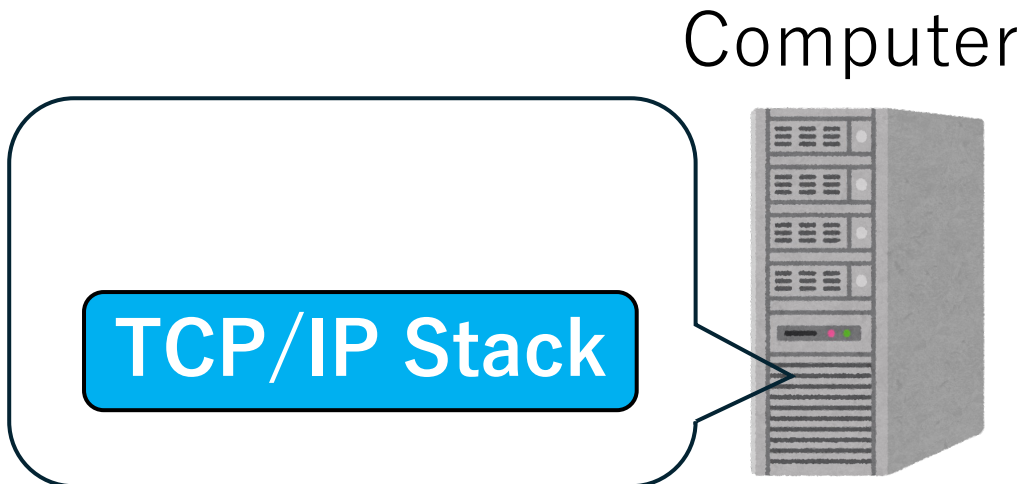
TCP/IP and TCP/IP Stacks

- TCP/IP is a standardized protocol suite commonly used for communication in computer networks
- TCP/IP stacks are typically software that implements procedures to comply with the TCP/IP standard



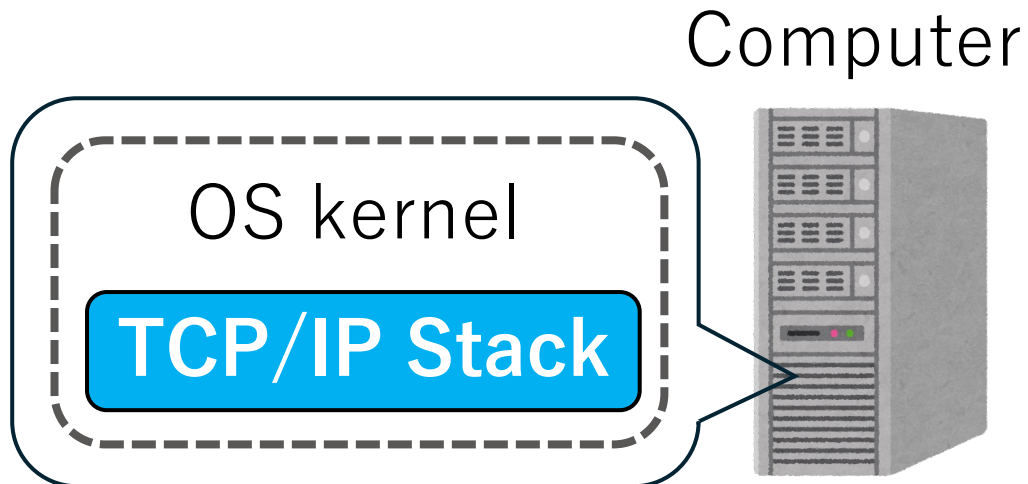
TCP/IP Stacks in Legacy OS Kernels

- TCP/IP stacks are typically maintained as part of OS kernels



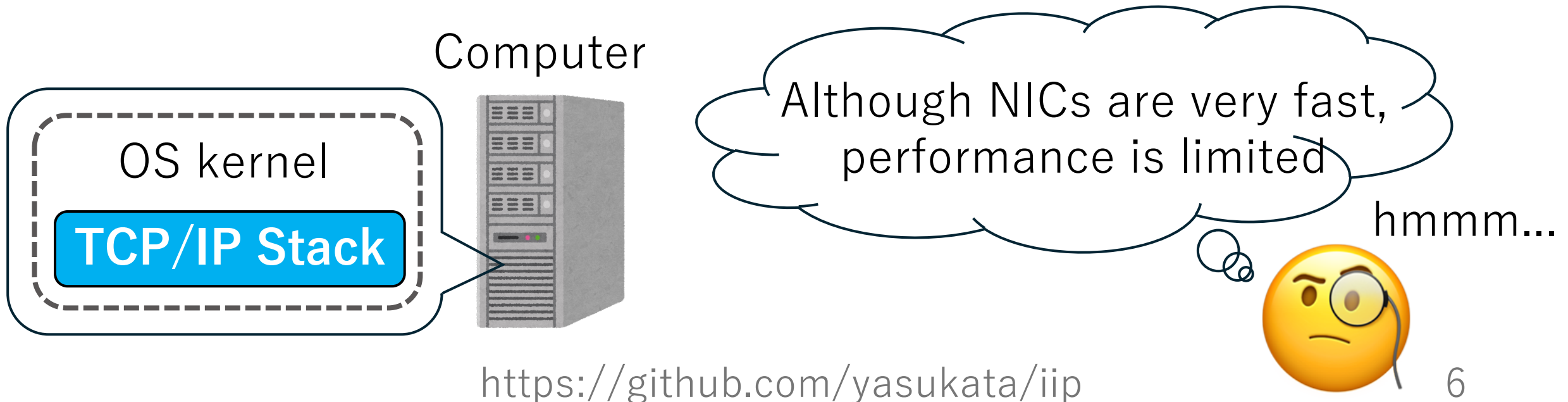
TCP/IP Stacks in Legacy OS Kernels

- TCP/IP stacks are typically maintained as part of OS kernels



TCP/IP Stacks in Legacy OS Kernels

- TCP/IP stacks are typically maintained as part of OS kernels
- People found it is hard for TCP/IP stacks in legacy OS kernels to effectively utilize the benefits of high-speed NICs



Performance-optimized TCP/IP Stacks

- To address this issue, research and industry communities have invented many performance-optimized TCP/IP stacks
 - e.g., Sandstorm (SIGCOMM'14), mTCP (NSDI'14)

Computer

**Performance-
optimized
TCP/IP Stack**



Power of NICs is unleashed
by fast TCP/IP stacks!



<https://github.com/yasukata/iip>

Performance-optimized TCP/IP Stacks



Their performance is excellent

e.g., Sandstorm (SIGCOMM'14), mTCP (NSDI'14)

Performance-optimized TCP/IP Stacks



Their performance is excellent

e.g., Sandstorm (SIGCOMM'14), mTCP (NSDI'14)



They often incur high integration complexity

Category	Performance	Integration
Performance-optimized	✓	

Portability-aware TCP/IP Stacks



There are TCP/IP stacks that allow for easy integration
e.g., lwIP, FNET, picoTCP

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Portability-aware TCP/IP Stacks



There are TCP/IP stacks that allow for easy integration
e.g., lwIP, FNET, picoTCP



They often lack the care for performance-critical factors

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- None of previous TCP/IP stack implementations allow for easy integration and good performance simultaneously

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- As a result, developers only had limited and laborious options

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- As a result, developers only had limited and laborious options
 - intensively modifying one of the existing TCP/IP stacks

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- As a result, developers only had limited and laborious options
 - intensively modifying one of the existing TCP/IP stacks
 - building a new TCP/IP stack from scratch

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- As a result, developers only had limited and laborious options
 - intensively modifying one of the existing TCP/IP stacks
 - building a new TCP/IP stack from scratch
 - accepting performance limitations of an applicable TCP/IP stack

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

Problem

- As a result, developers only had limited and laborious options
 - intensively modifying one of the existing TCP/IP stacks
 - building a new TCP/IP stack from scratch
 - accepting performance limitations of an applicable TCP/IP stack
 - giving up the integration

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓

This Work

- We develop iip, an integratable TCP/IP stack, that allows for easy integration and good performance simultaneously

Category	Performance	Integration
Performance-optimized	✓	
Portability-aware		✓
This work	✓	✓

Issues of Existing TCP/IP Stacks

- Performance-optimized TCP/IP stacks
 - Dependencies on other components
 - Functionality conflicts
 - Limited choices for CPU core assignment models
- Portability-aware TCP/IP stacks
 - Unaware of NIC offloading features
 - Lack of zero-copy I/O capability
 - Lack of multi-core scalability

Issues of Existing TCP/IP Stacks

Performance-optimized
TCP/IP stack

Performance-optimized
TCP/IP stacks often consist of
various components

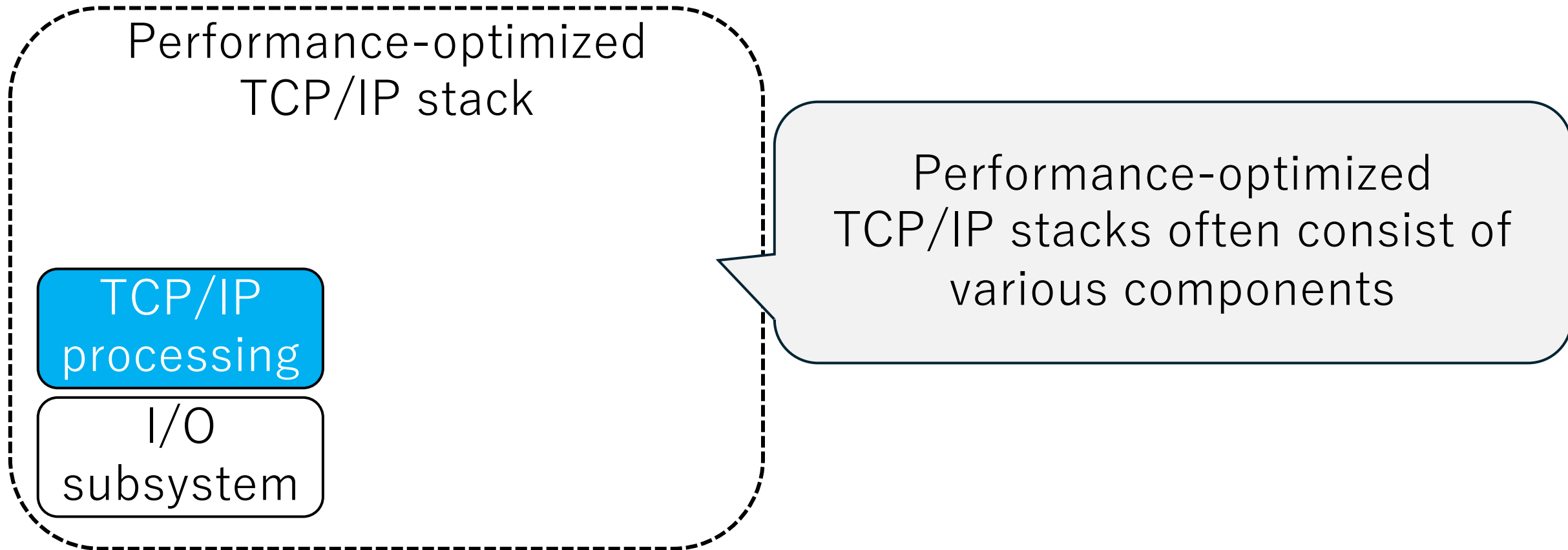
Issues of Existing TCP/IP Stacks

Performance-optimized
TCP/IP stack

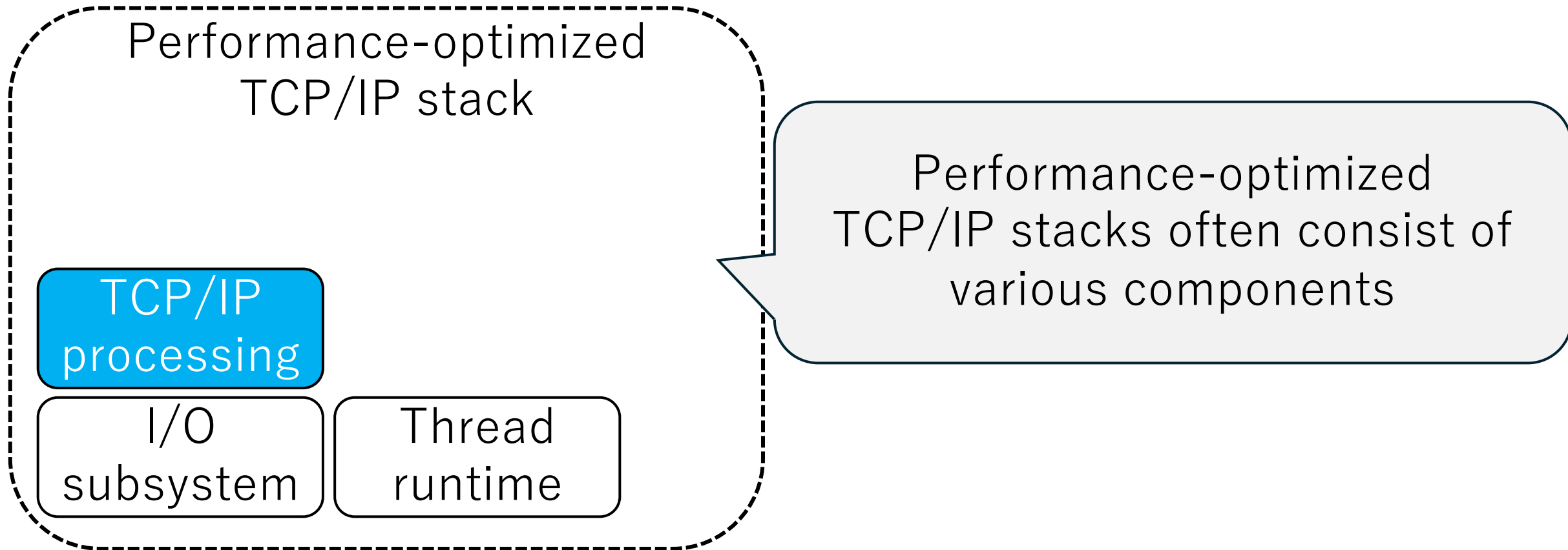
TCP/IP
processing

Performance-optimized
TCP/IP stacks often consist of
various components

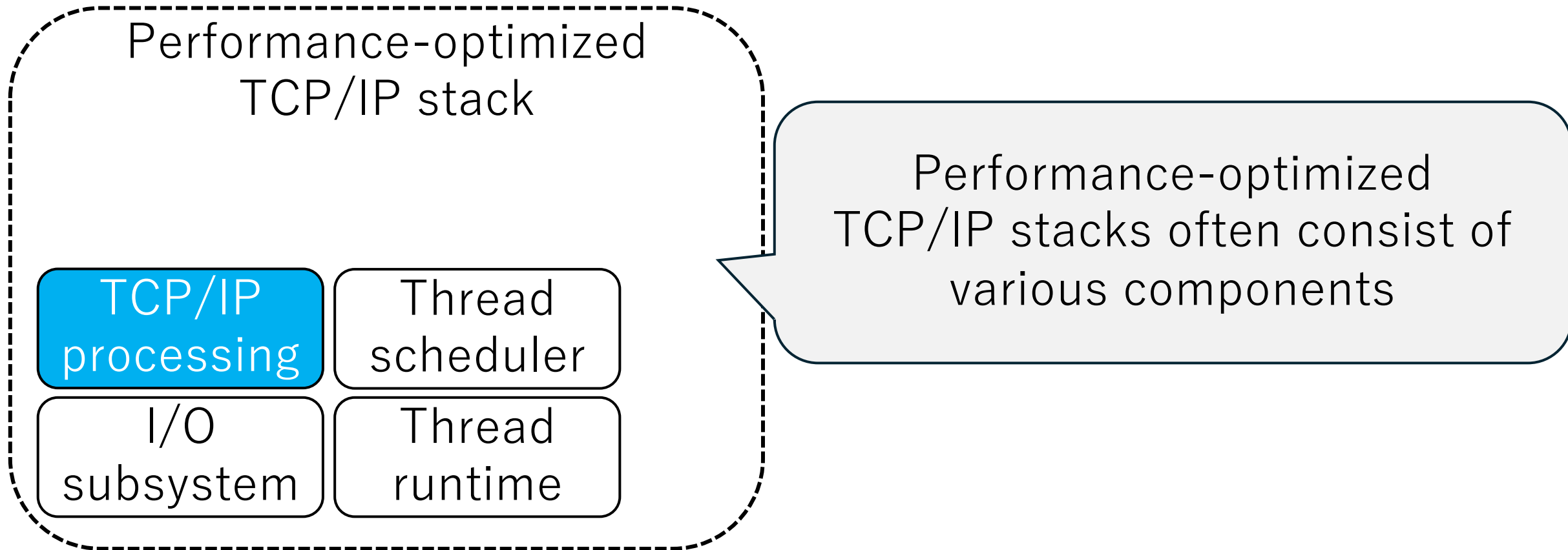
Issues of Existing TCP/IP Stacks



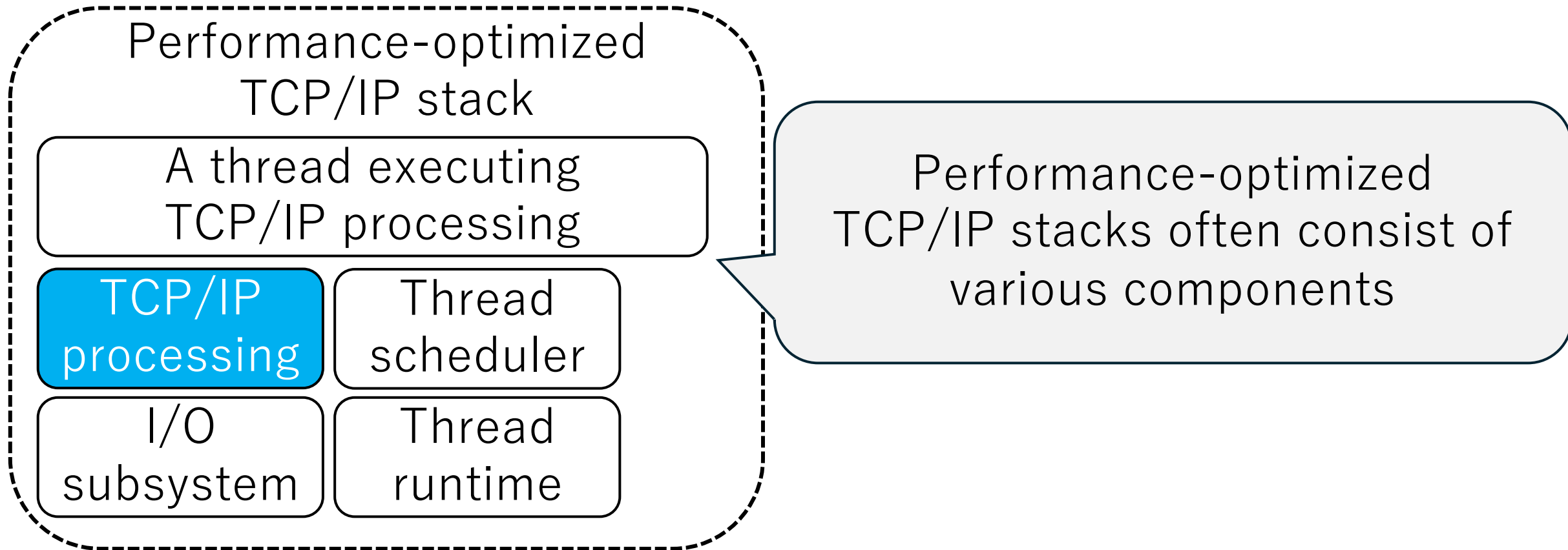
Issues of Existing TCP/IP Stacks



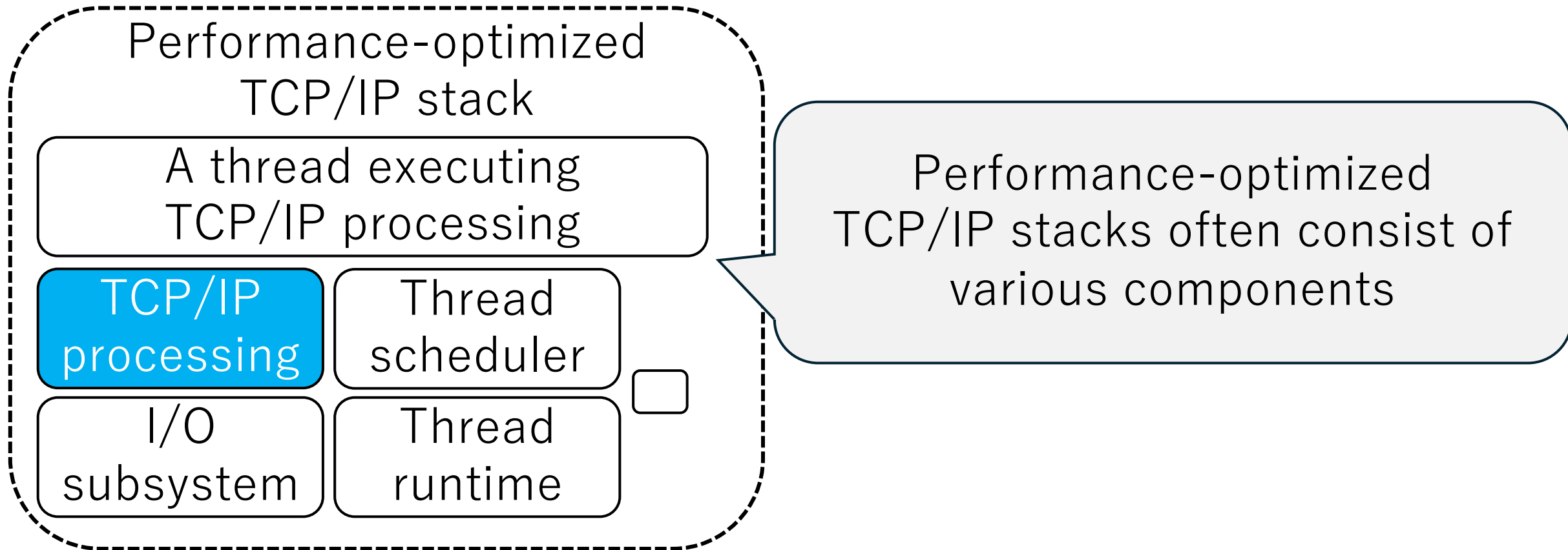
Issues of Existing TCP/IP Stacks



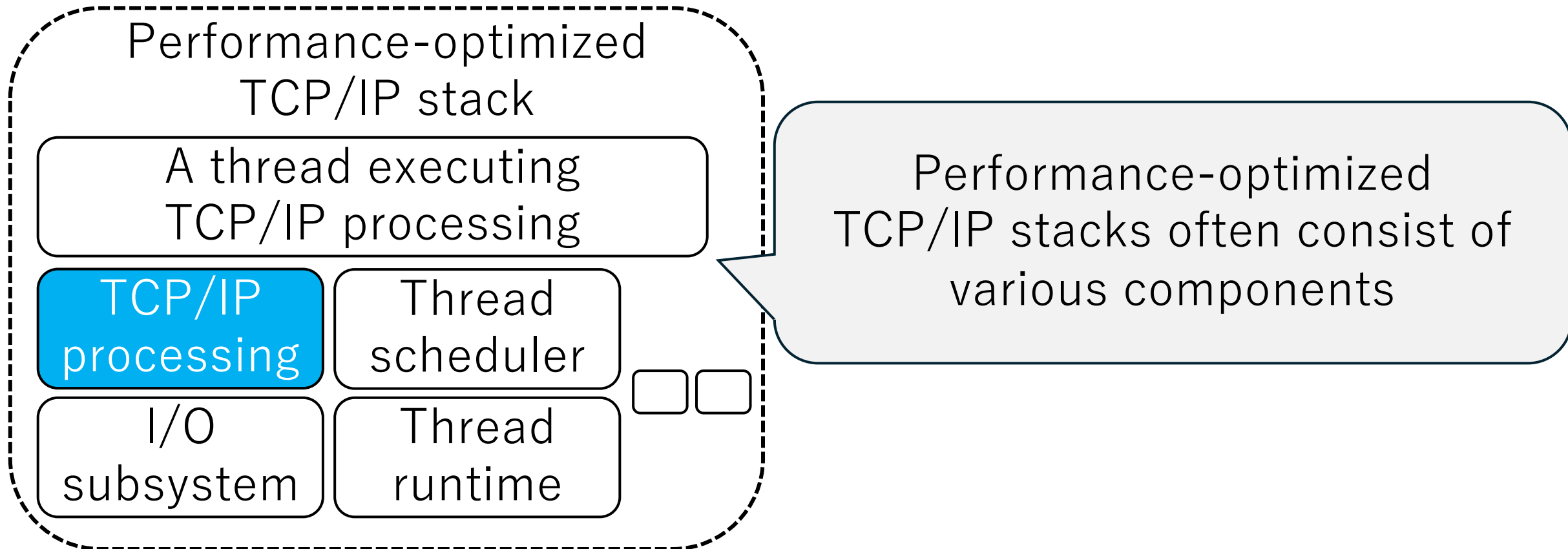
Issues of Existing TCP/IP Stacks



Issues of Existing TCP/IP Stacks

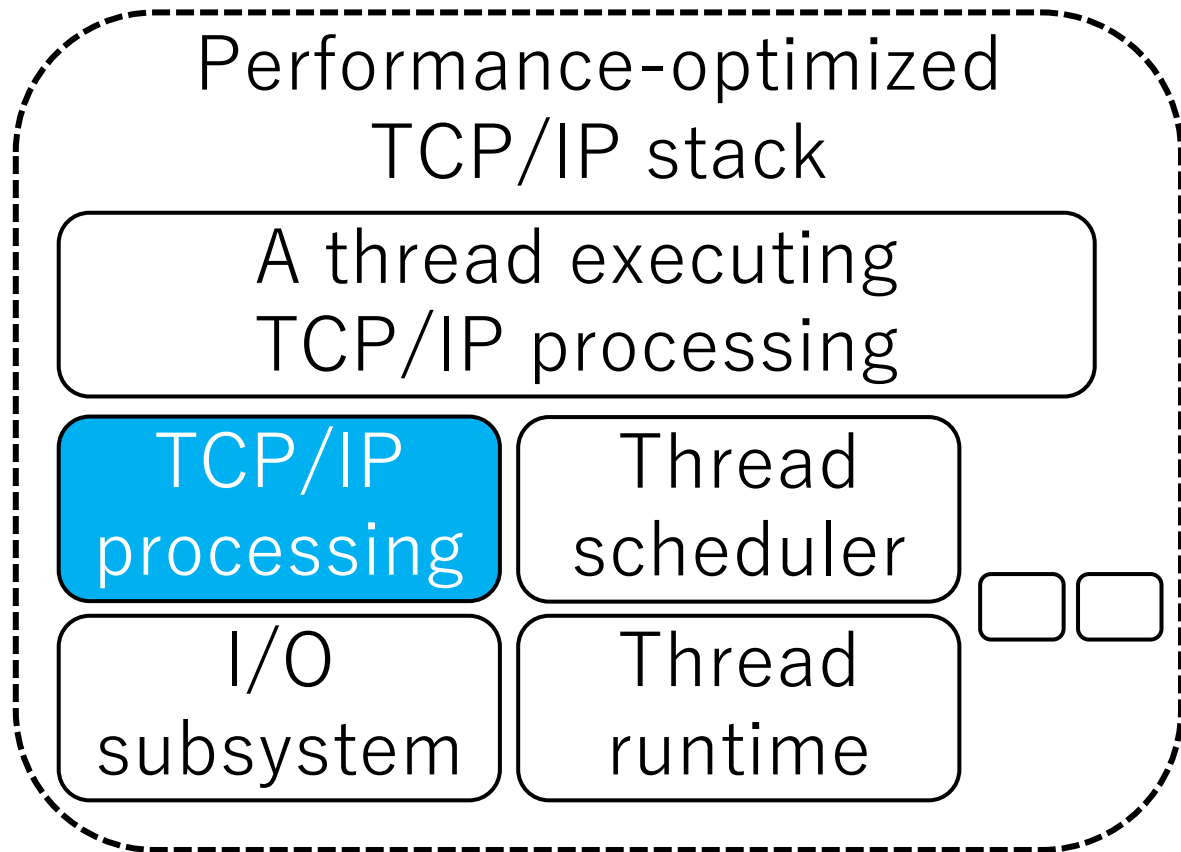


Issues of Existing TCP/IP Stacks



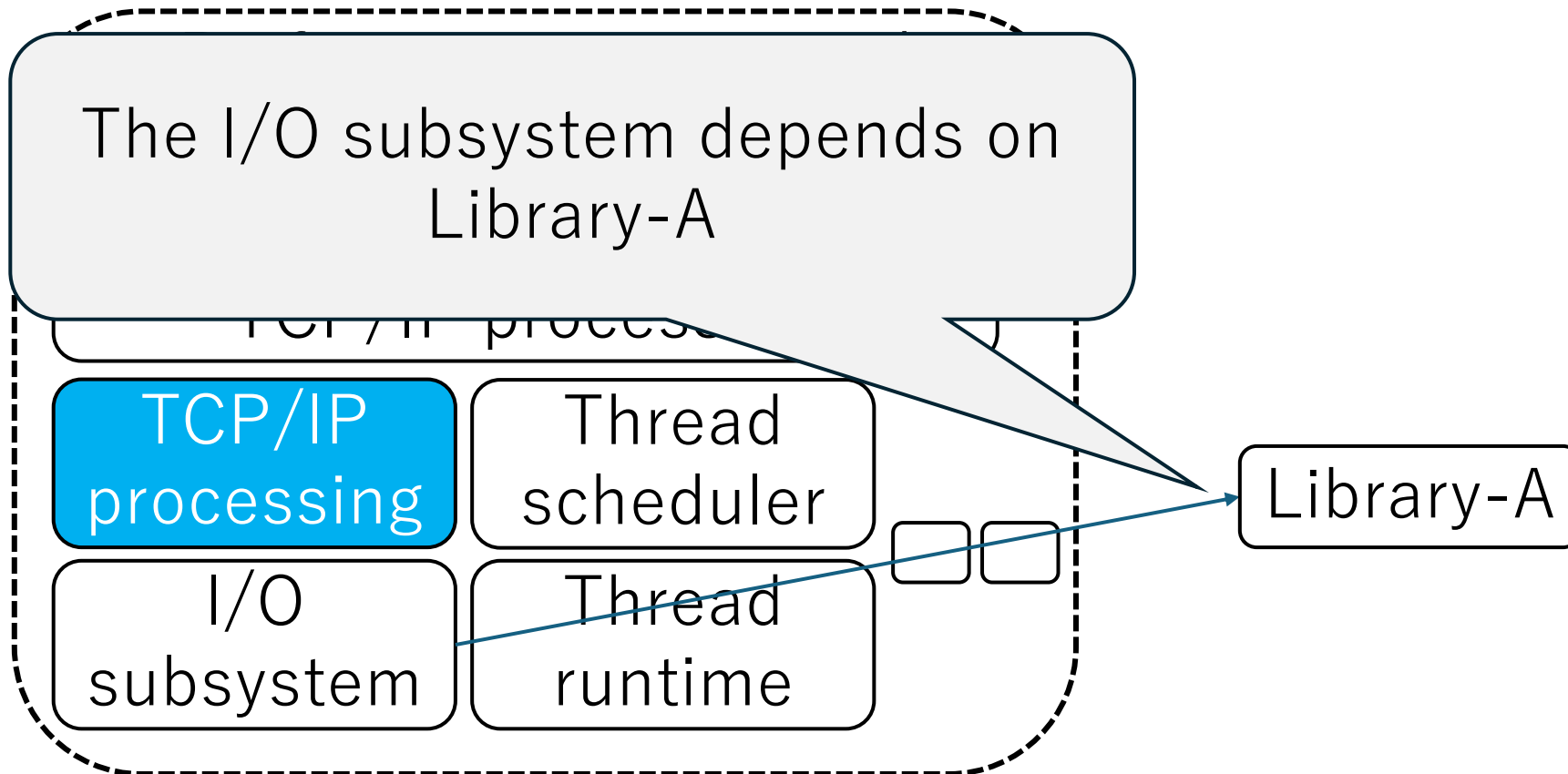
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



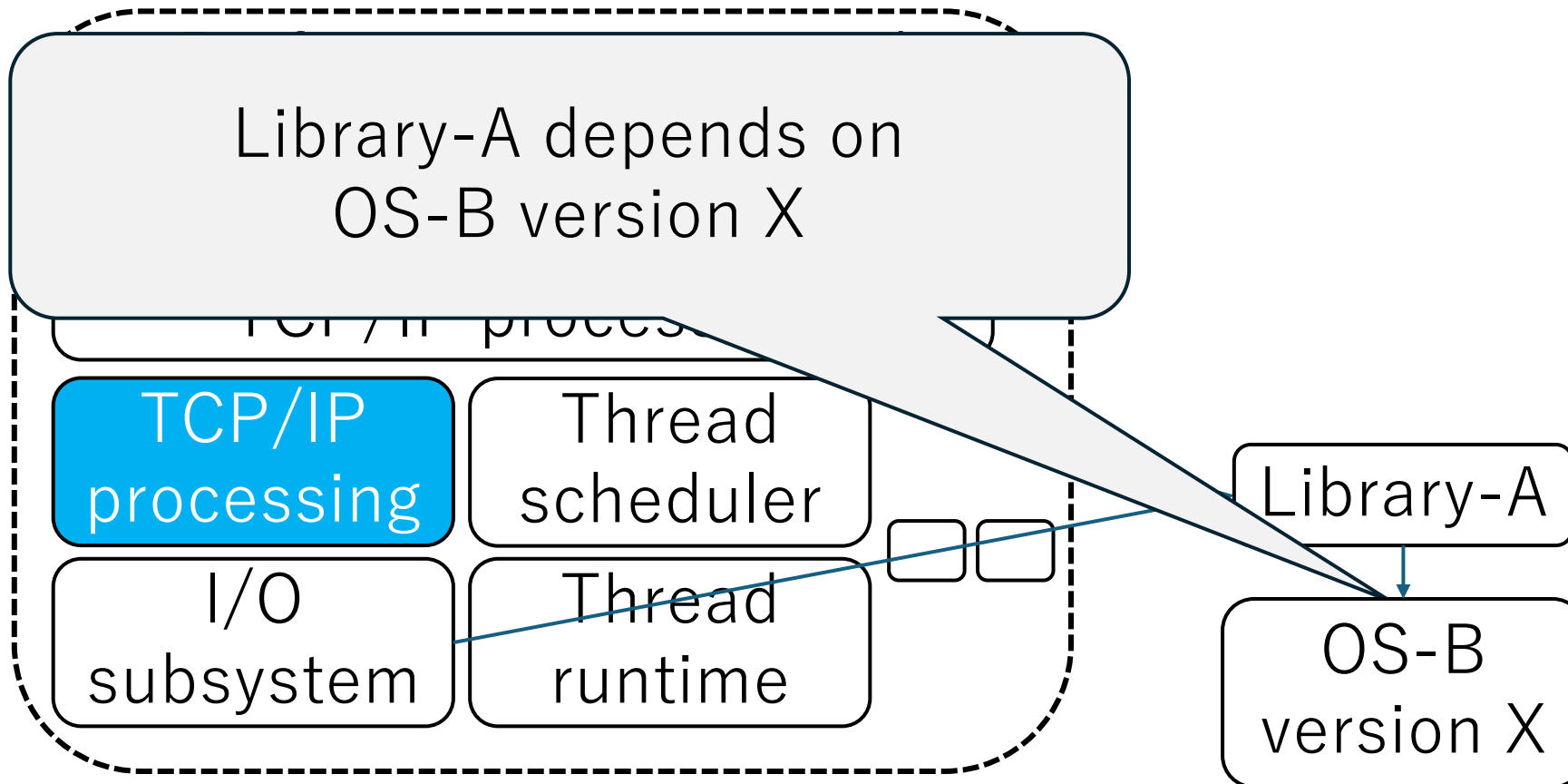
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



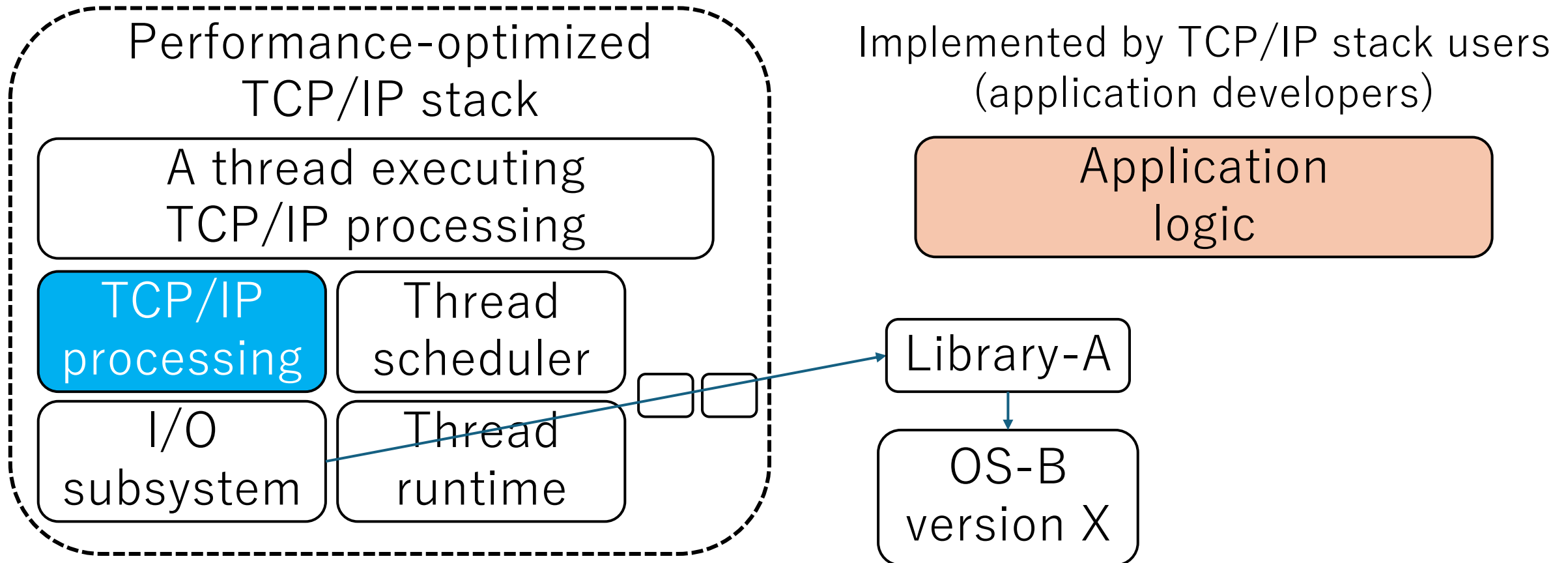
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



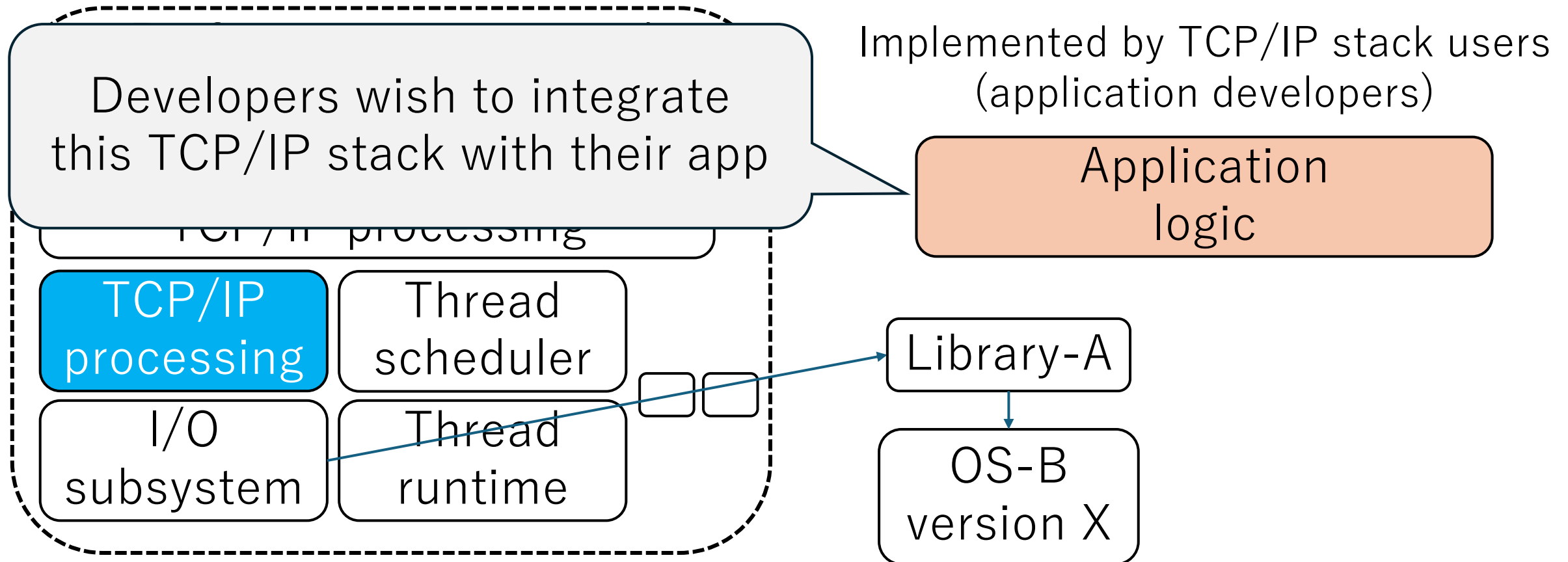
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



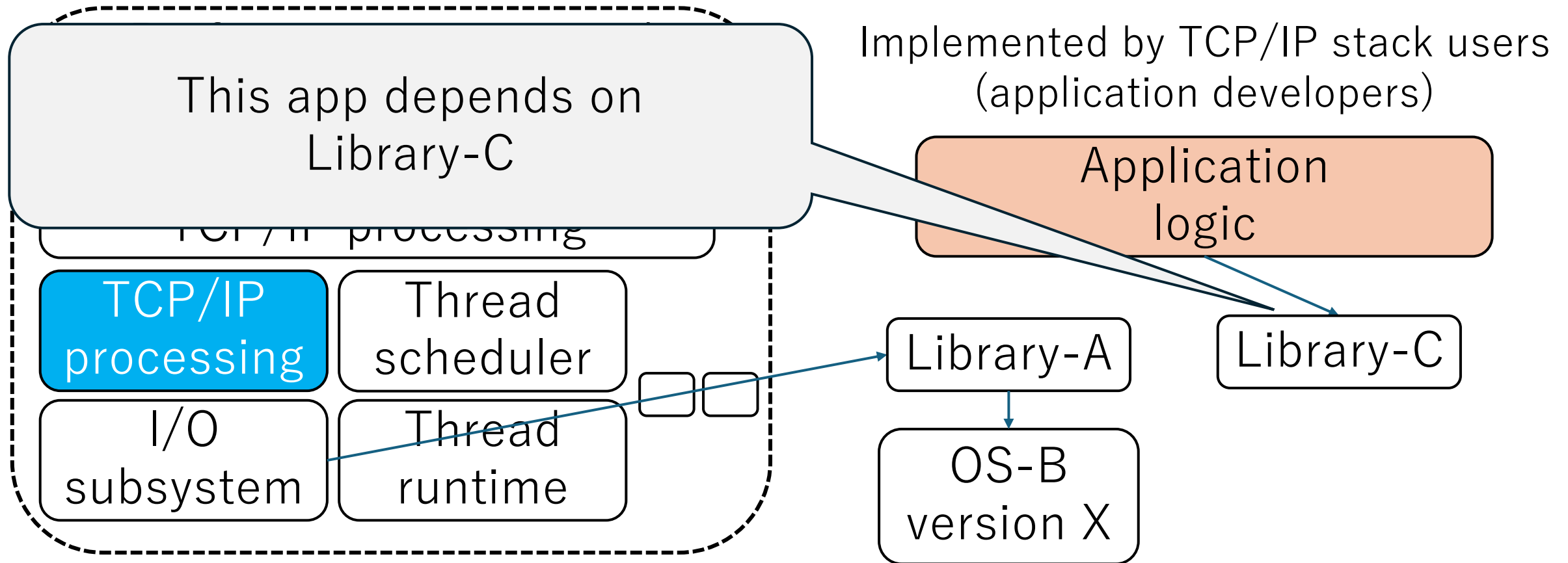
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



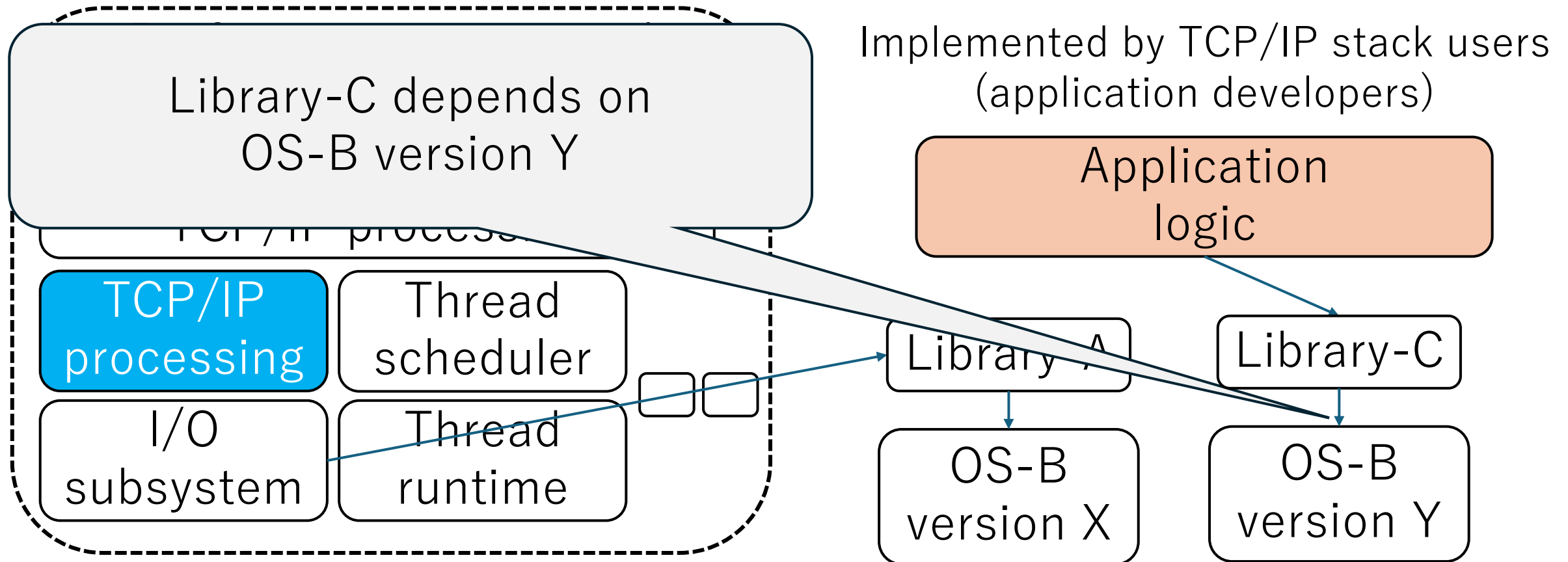
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



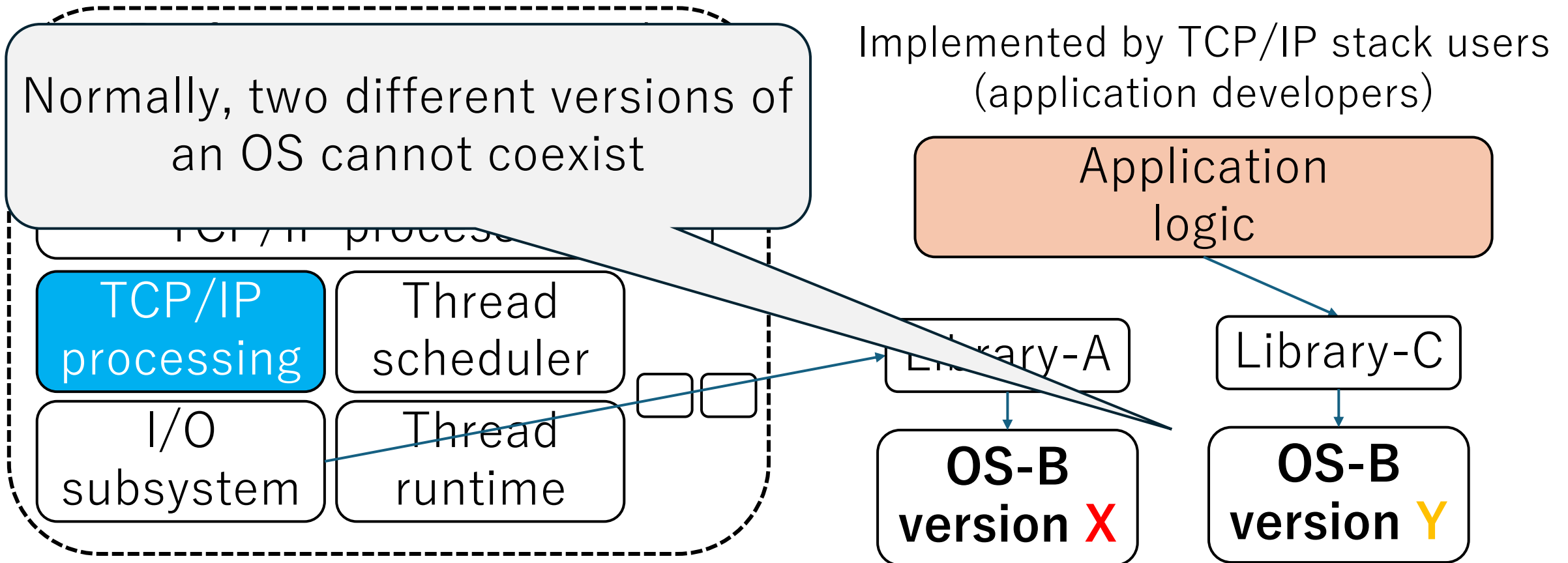
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



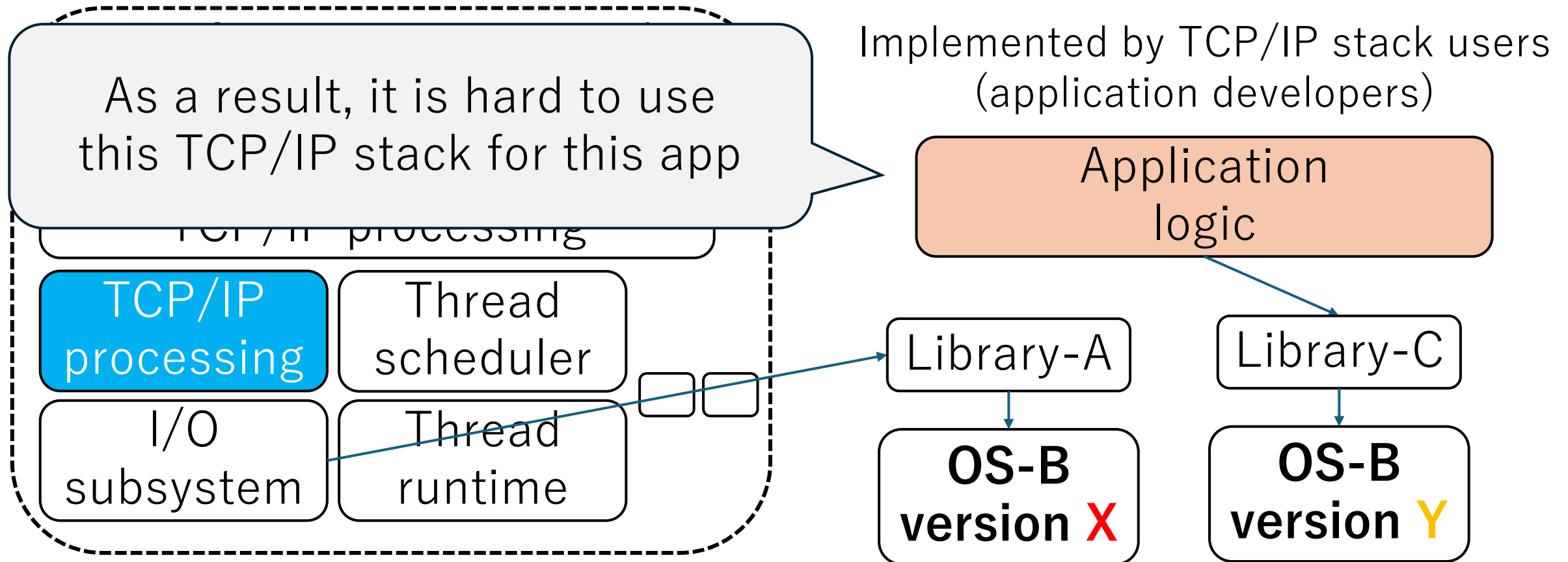
Issues of Existing TCP/IP Stacks

Dependencies on Other Components



Issues of Existing TCP/IP Stacks

Dependencies on Other Components



Issues of Existing TCP/IP Stacks

Dependencies on Other Components

Dependencies often increase integration complexity

Implemented by TCP/IP stack users
(application developers)

TCP/IP processing

Thread scheduler

I/O subsystem

Thread runtime

Library-A

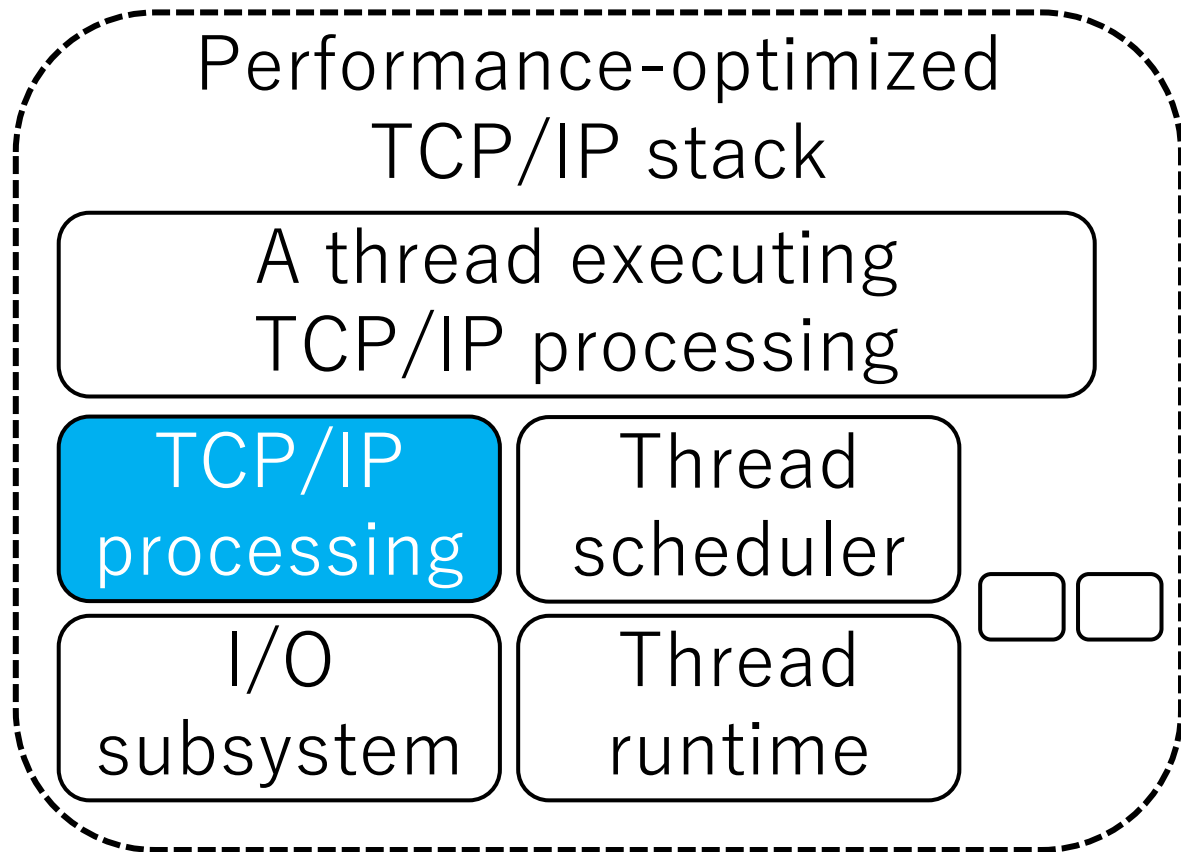
Library-C

**OS-B
version X**

**OS-B
version Y**

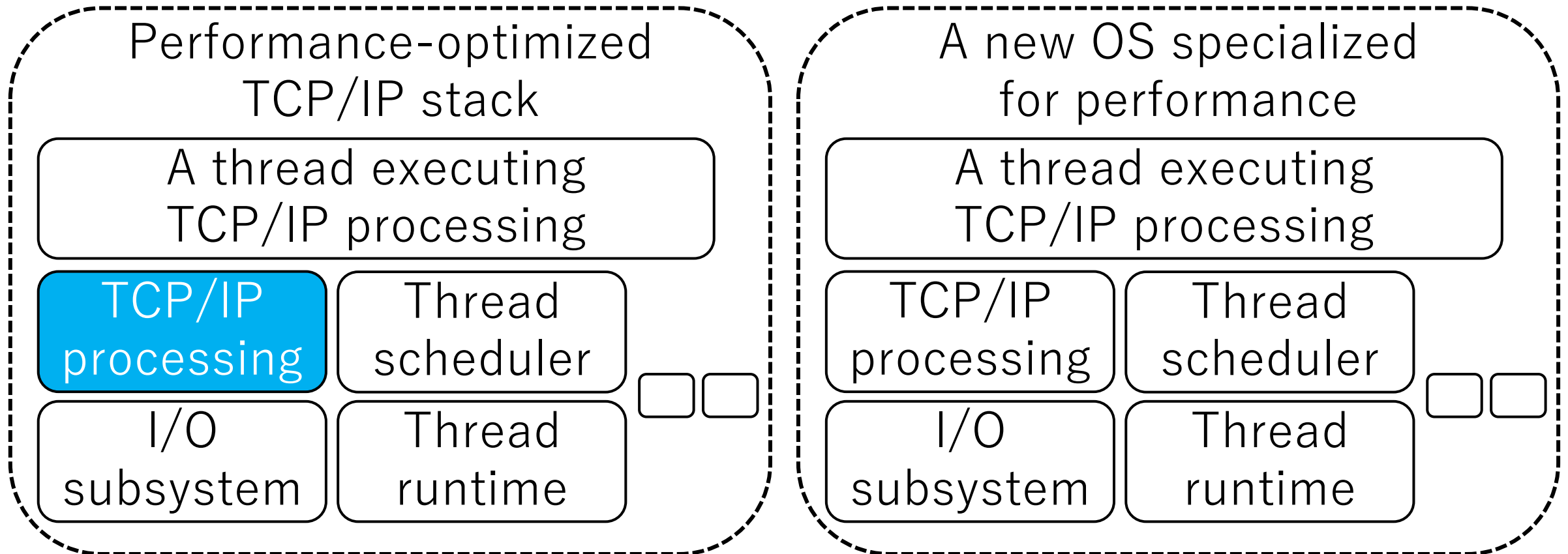
Issues of Existing TCP/IP Stacks

Functionality Conflicts



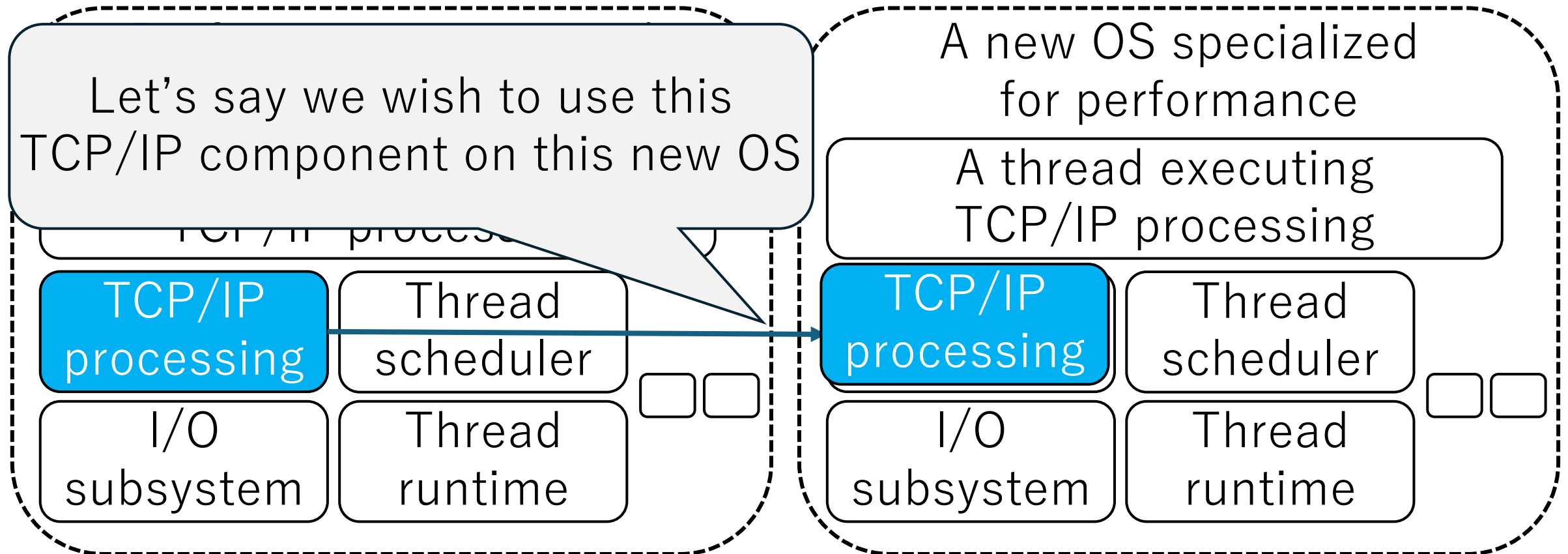
Issues of Existing TCP/IP Stacks

Functionality Conflicts



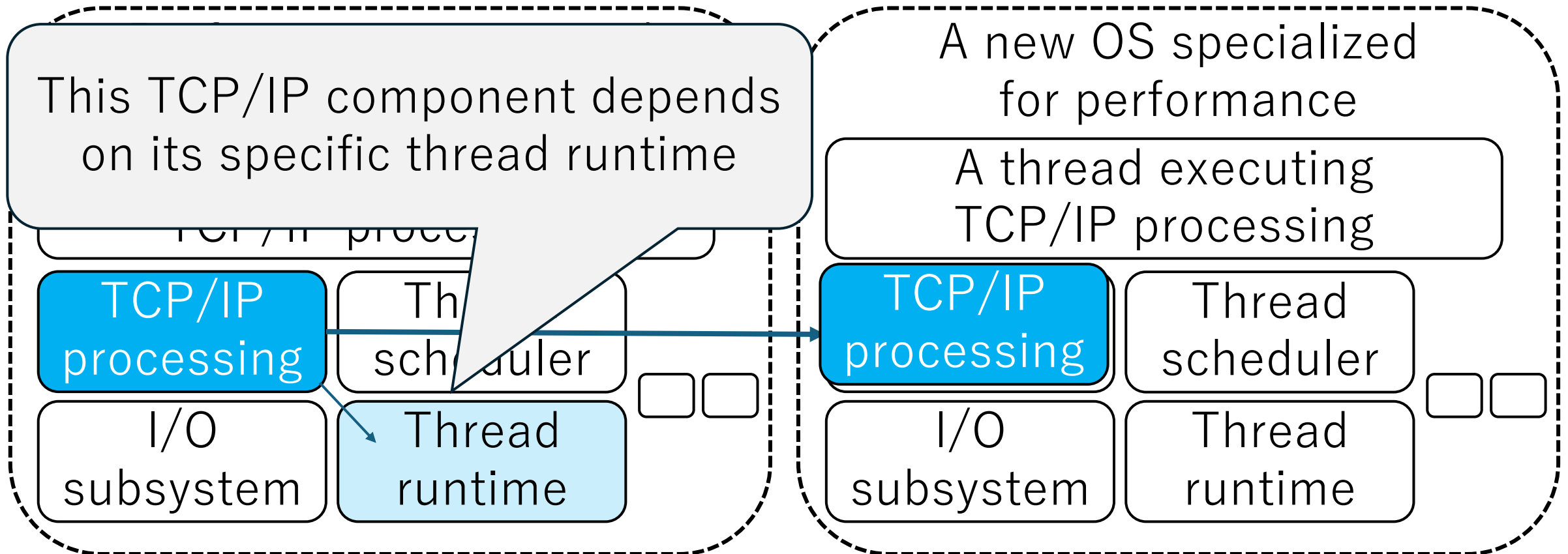
Issues of Existing TCP/IP Stacks

Functionality Conflicts



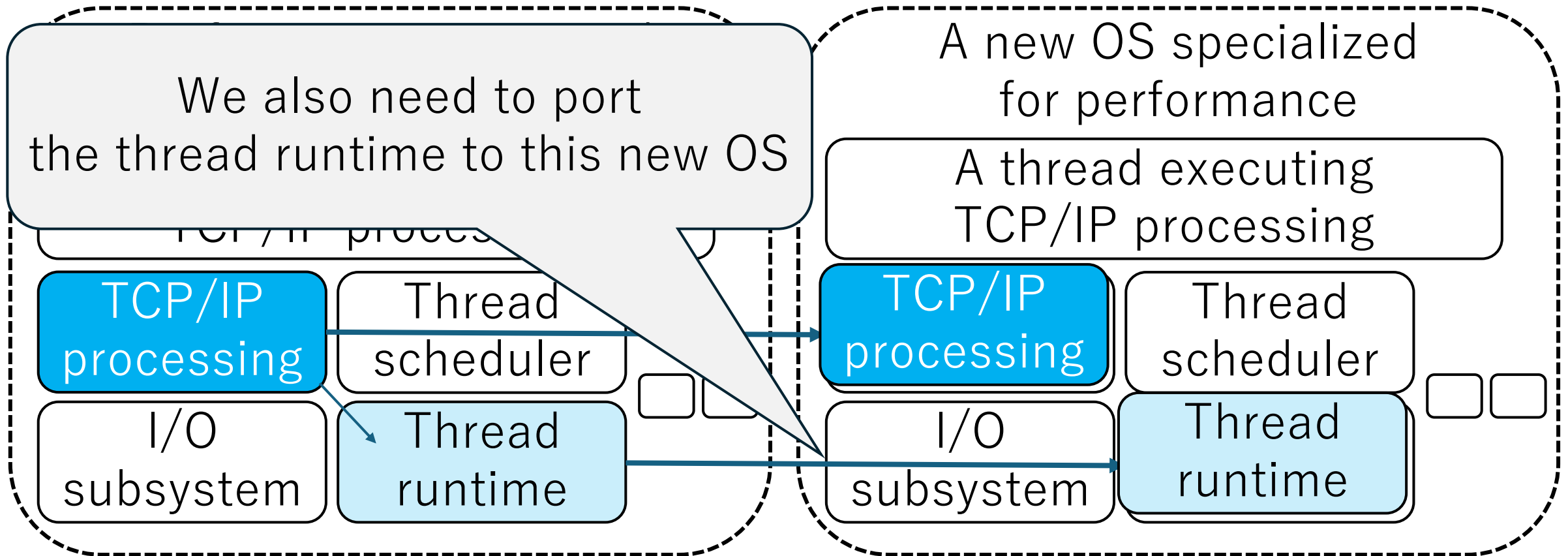
Issues of Existing TCP/IP Stacks

Functionality Conflicts



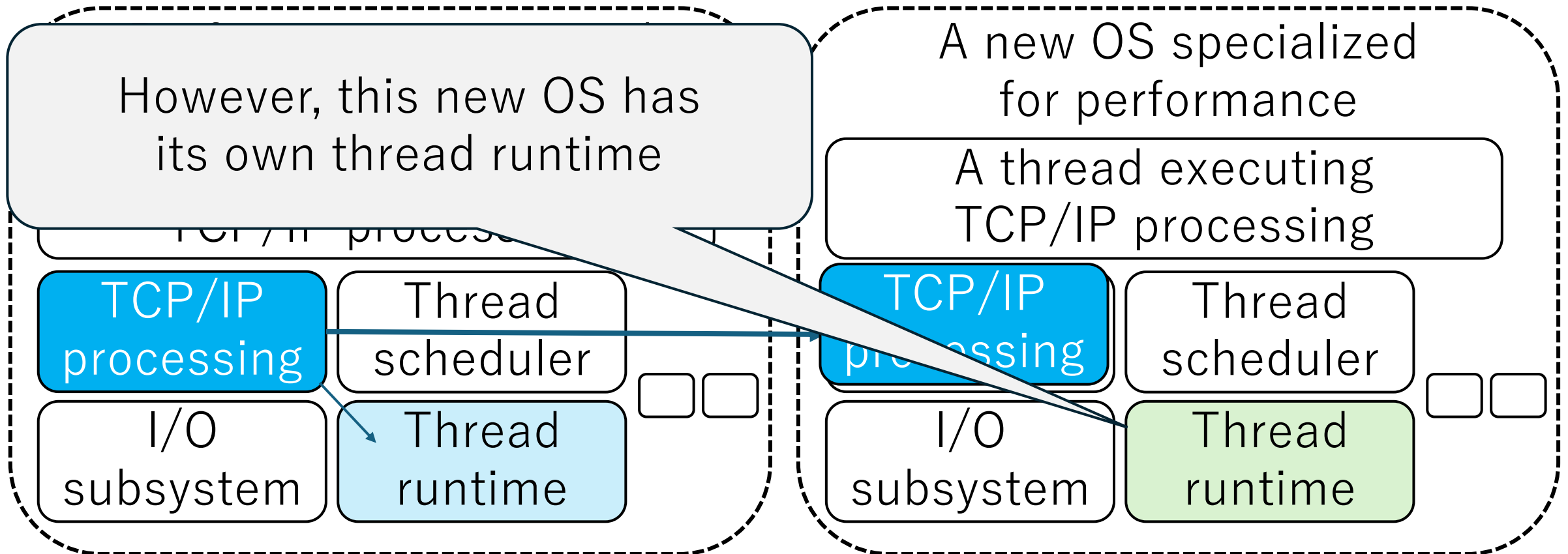
Issues of Existing TCP/IP Stacks

Functionality Conflicts



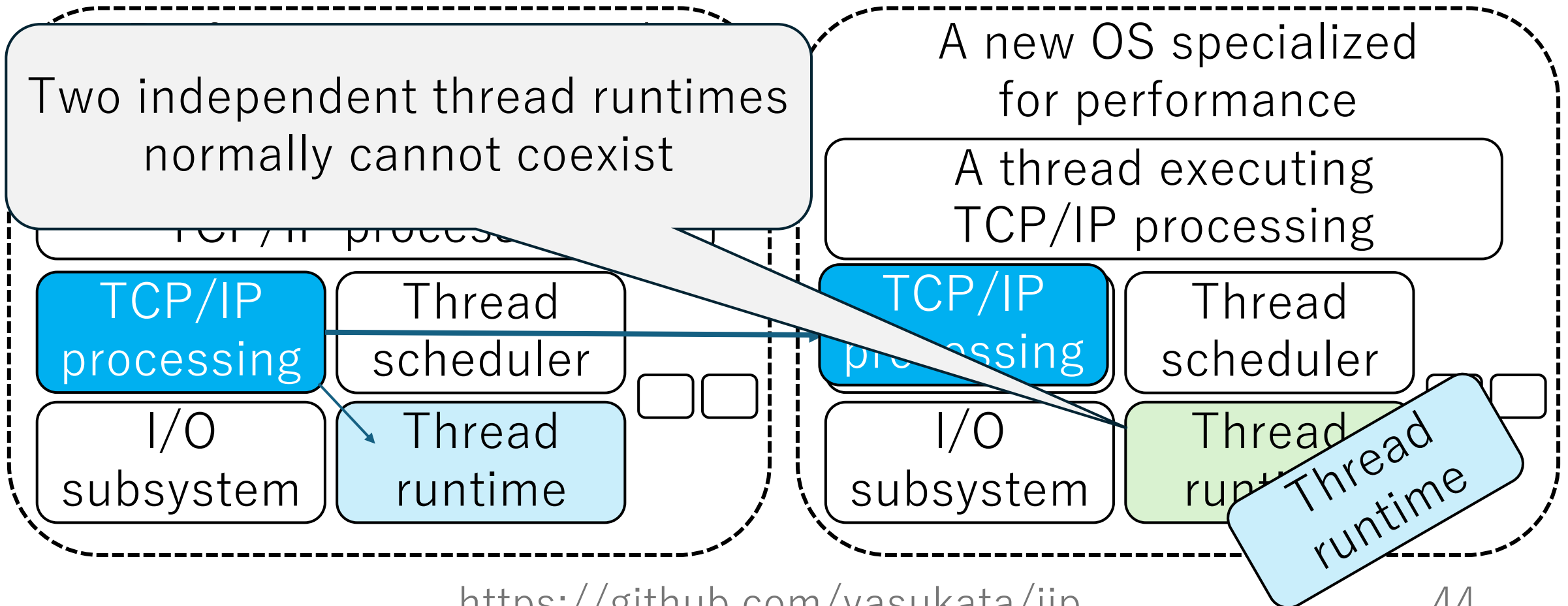
Issues of Existing TCP/IP Stacks

Functionality Conflicts



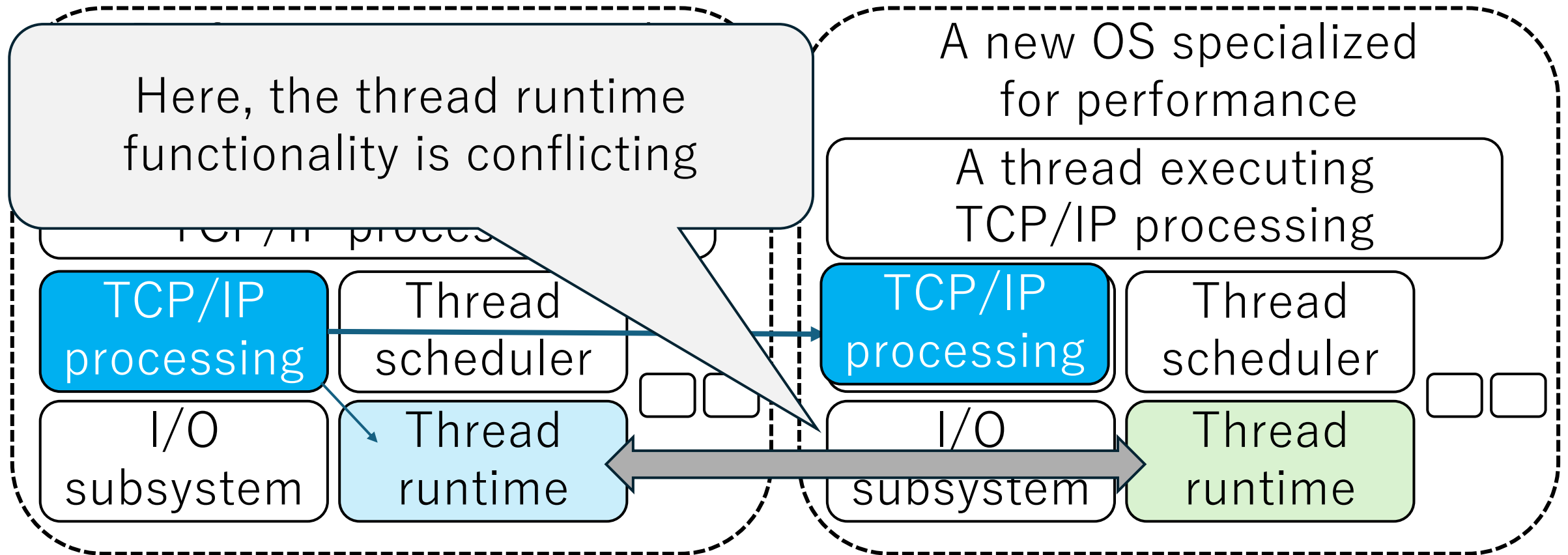
Issues of Existing TCP/IP Stacks

Functionality Conflicts



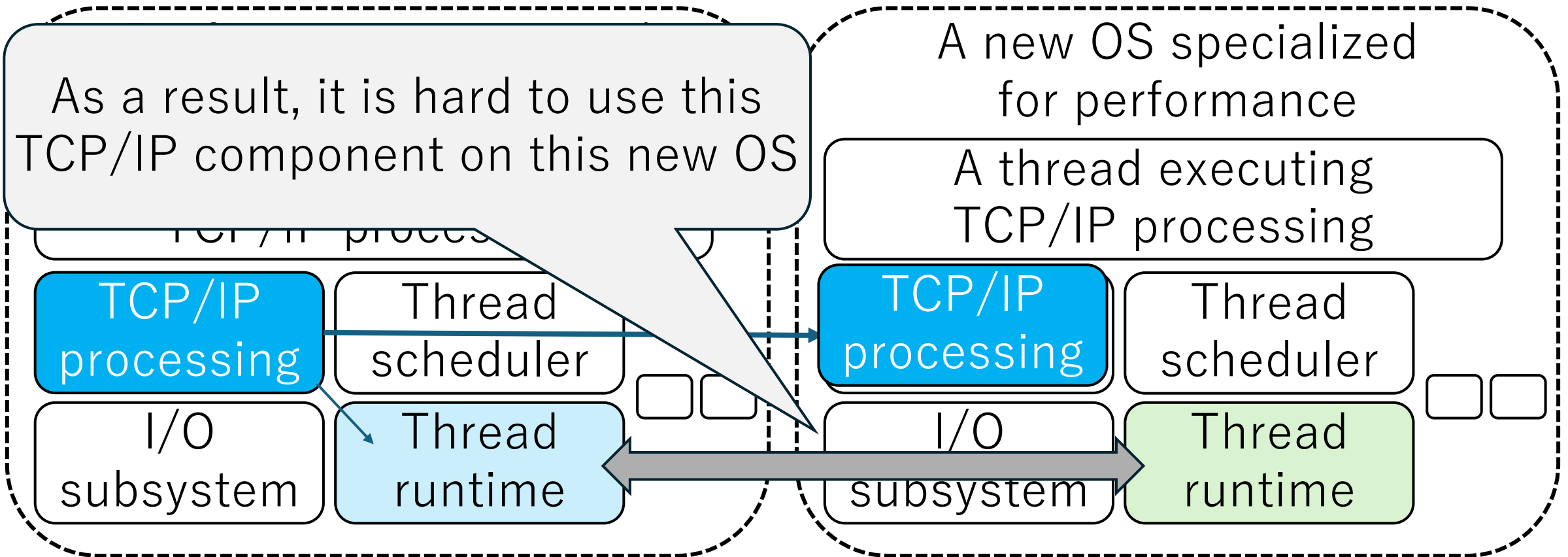
Issues of Existing TCP/IP Stacks

Functionality Conflicts



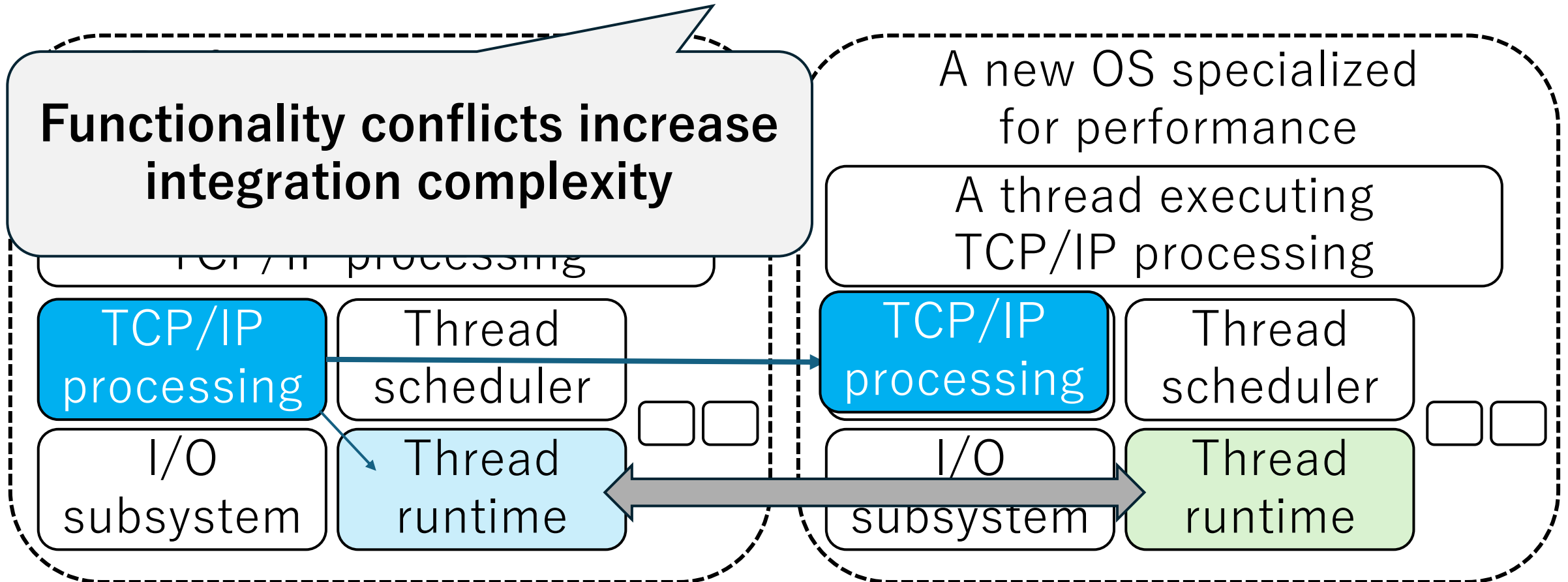
Issues of Existing TCP/IP Stacks

Functionality Conflicts



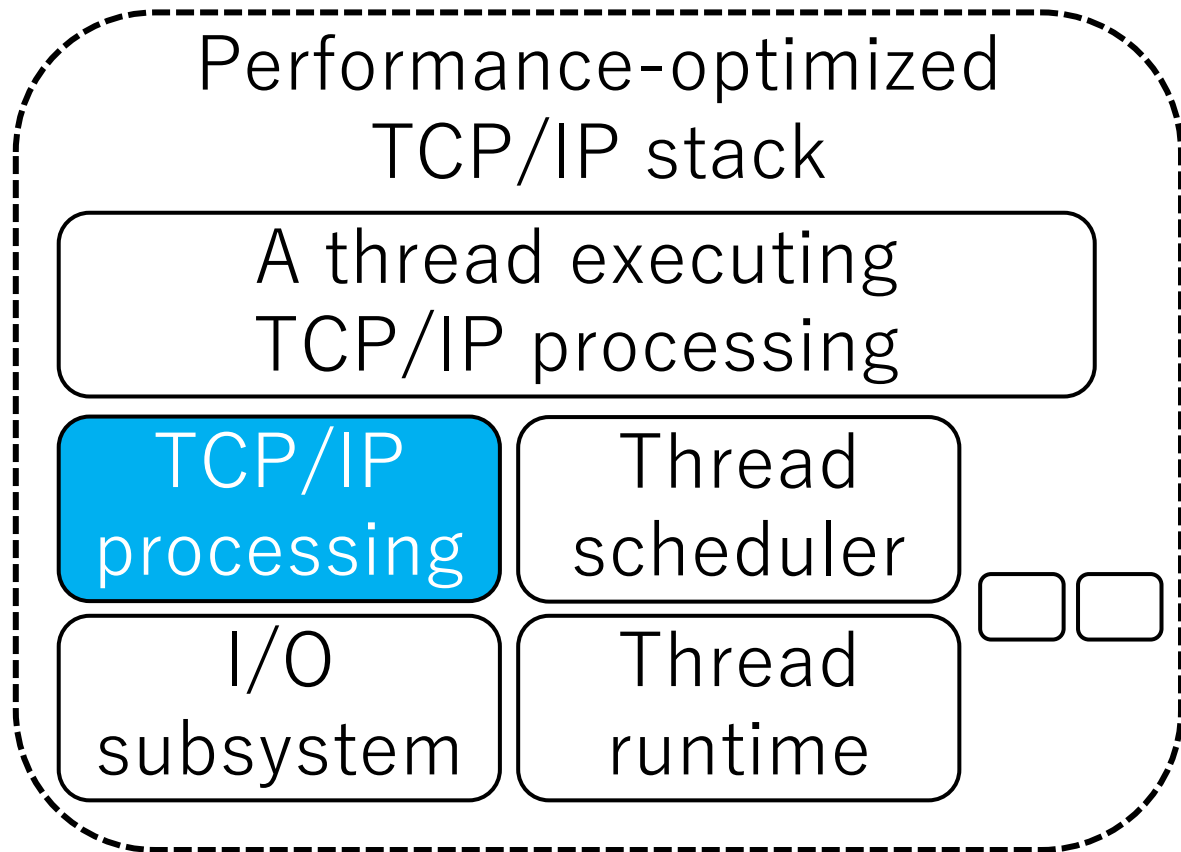
Issues of Existing TCP/IP Stacks

Functionality Conflicts



Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models



Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

Performance-optimized
TCP/IP stack

A thread executing
TCP/IP processing

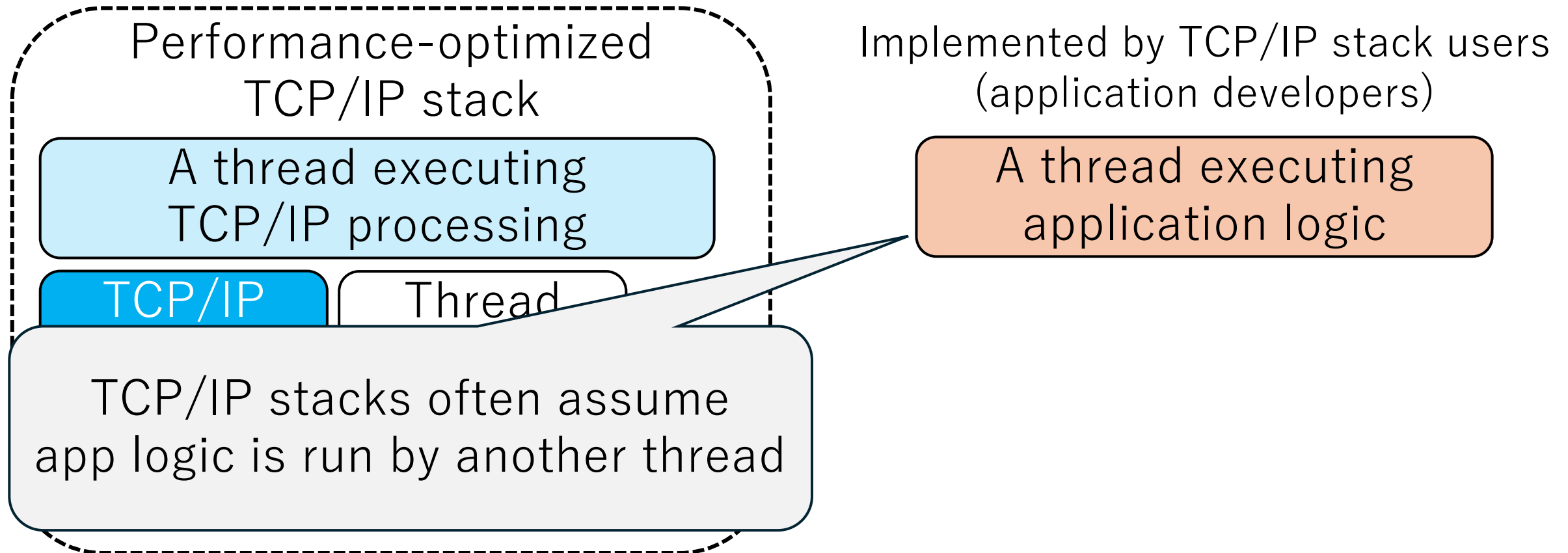
TCP/IP

Thre

TCP/IP stacks often include
threads for TCP/IP processing

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models



Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

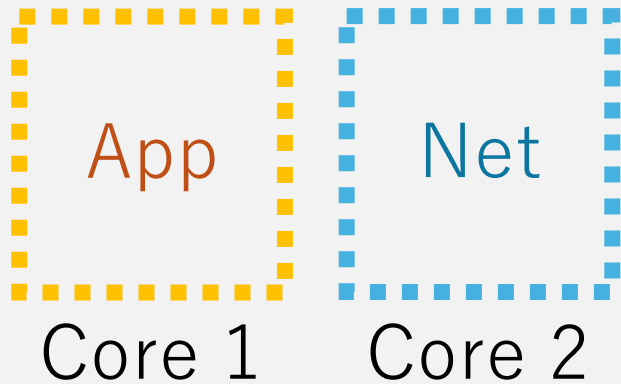
Three potential CPU core assignment models

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

Three potential CPU core assignment models

Split



App and Net threads run on different CPU cores

Issues of Existing TCP/IP Stacks

Limited

App and Net threads run on the same CPU core

Implementation Models

Thread

Implementation models

Split

Merge



Core 1

Core 2

Core 1

Issues of Existing TCP/IP Stacks

Limited

Duplicate the same setup
to available CPU cores

Deployment Models

Thre

Deployment models

Split

Merge



Core 1

Core 2



Core 1

Core 2

Issues of Existing TCP/IP Stacks

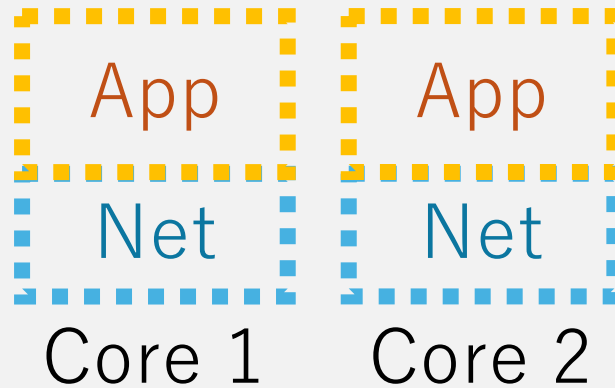
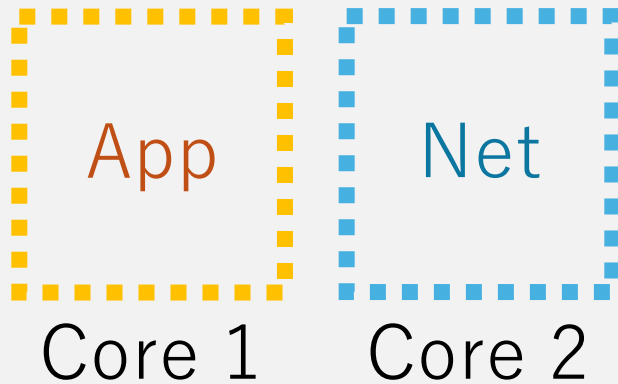
Limited Choices for CPU Core Assignment Models

A thread executes both App and Net logic

Split

Merge

Unified

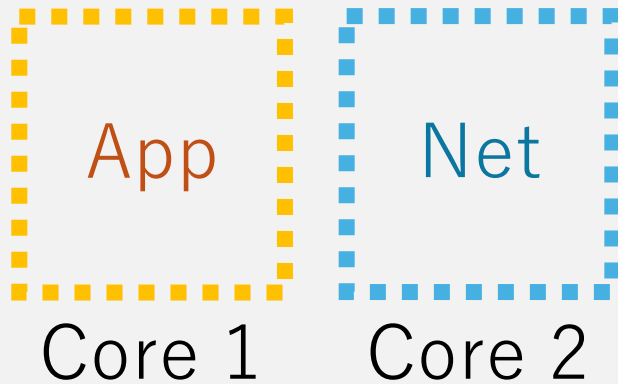


Issues of Existing TCP/IP Stacks

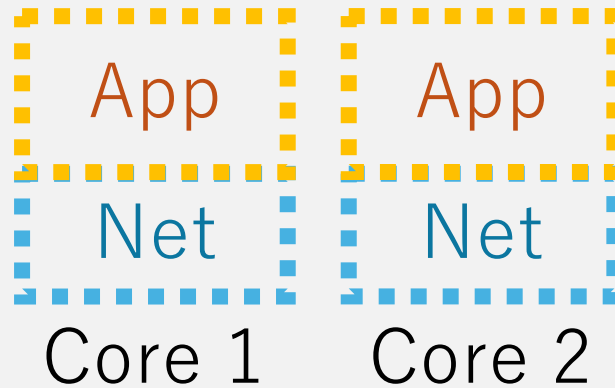
Limited Choices for CPU Core Assignment Models

Duplicate the setup to available CPU cores

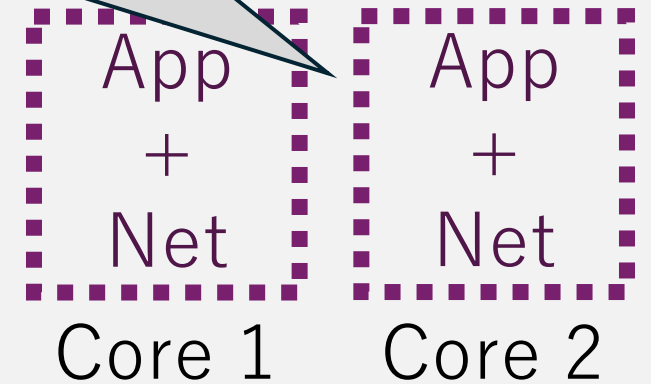
Split



Merge



Unified

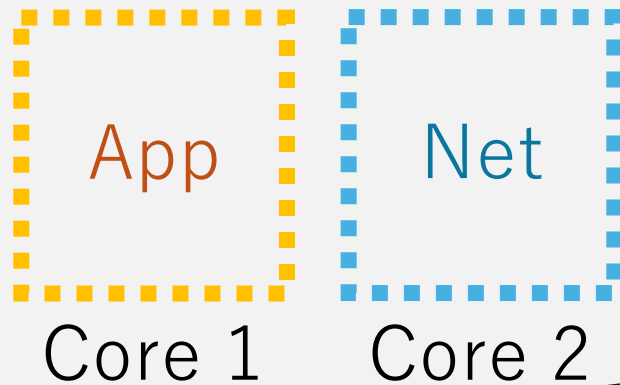


Issues of Existing TCP/IP Stacks

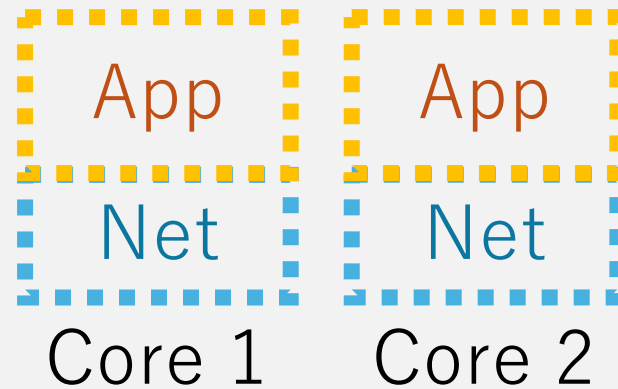
Limited Choices for CPU Core Assignment Models

Three potential CPU core assignment models

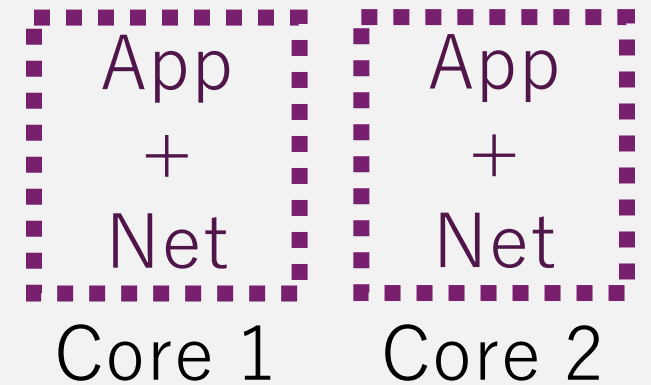
Split



Merge



Unified



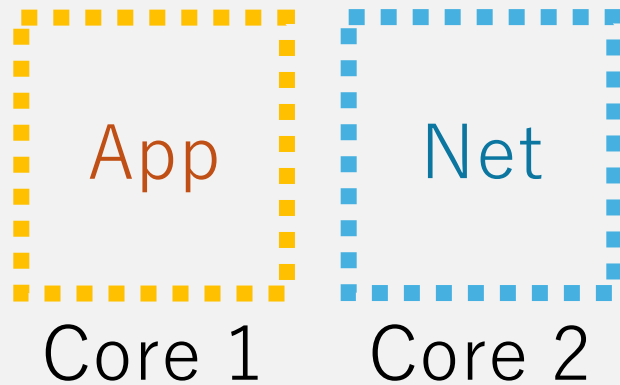
Each of them has different performance characteristics

Issues of Existing TCP/IP Stacks

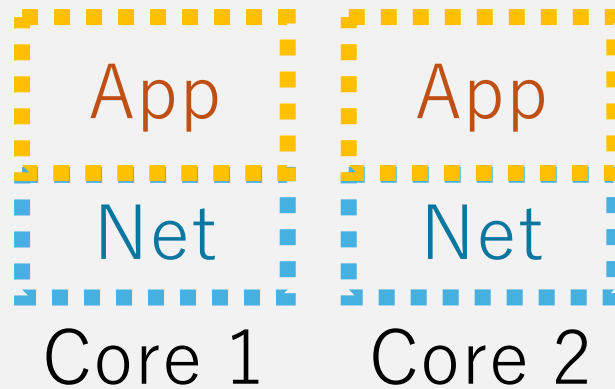
Limited Choices for CPU Core Assignment Models

Three potential CPU core assignment models

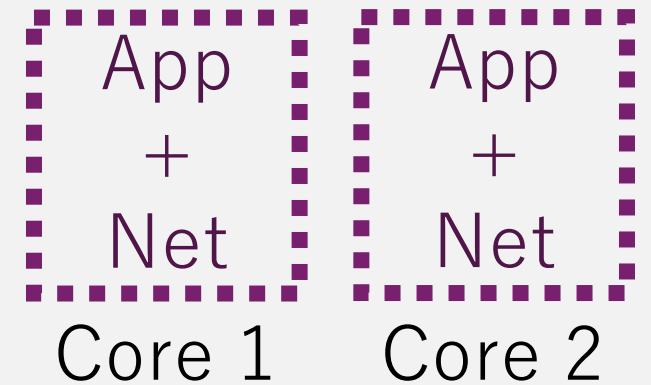
Split



Merge



Unified

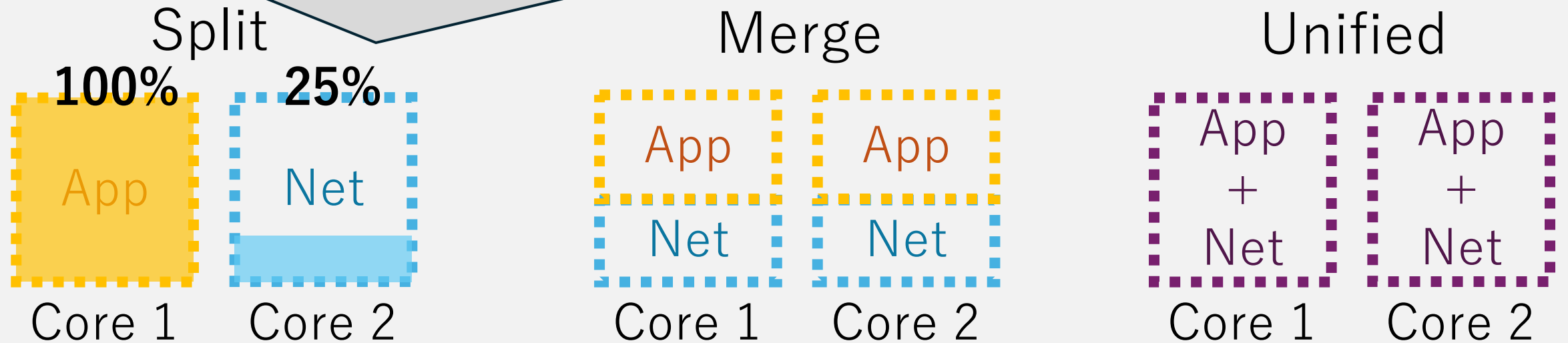


CPU utilization

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

A busy thread cannot use other CPU cores' unused cycles

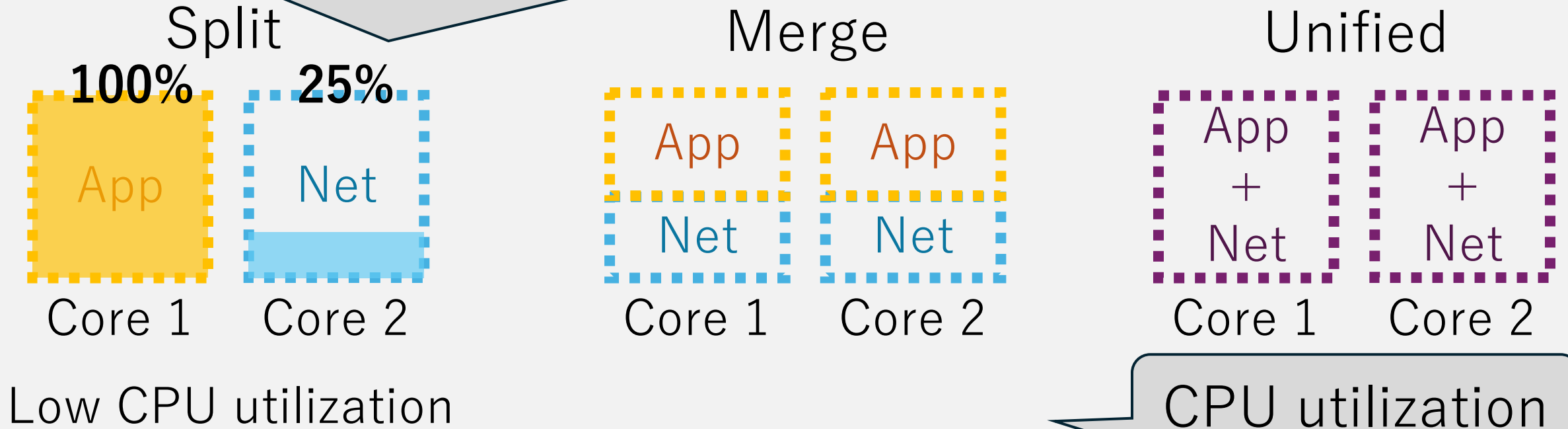


CPU utilization

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

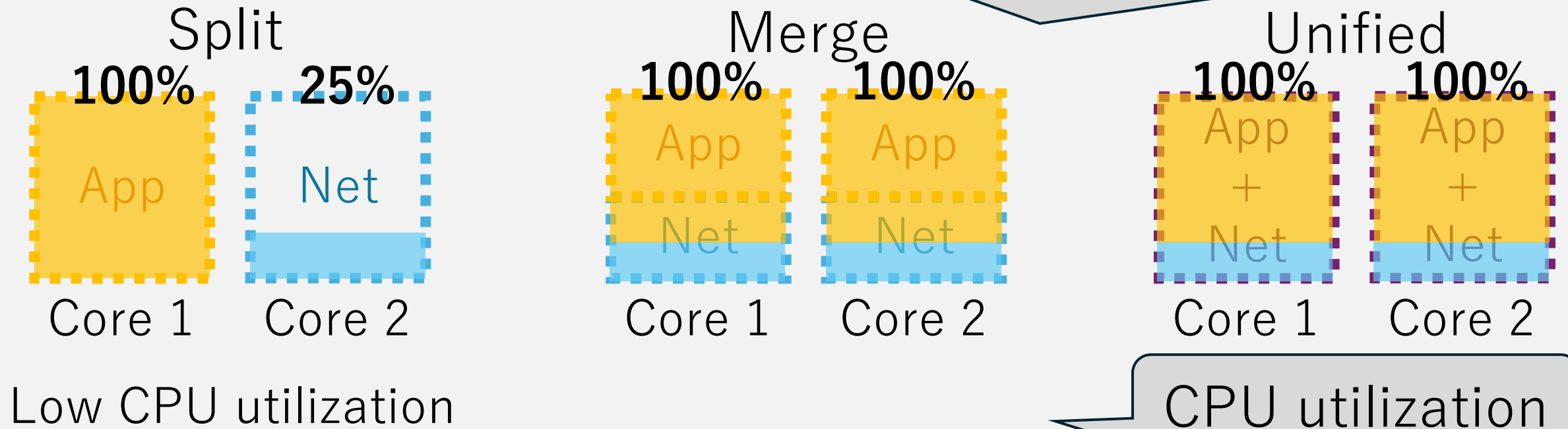
A busy thread cannot use other CPU cores' unused cycles



Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

The merge and unified models can fully use the CPU cycles

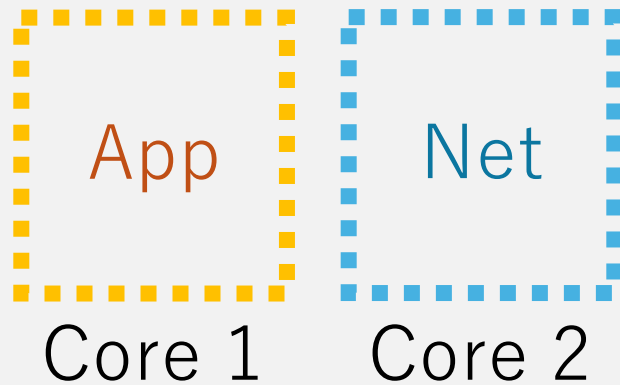


Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

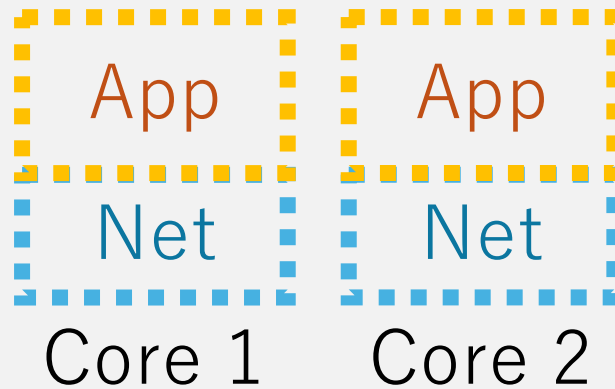
Three potential CPU core assignment models

Split

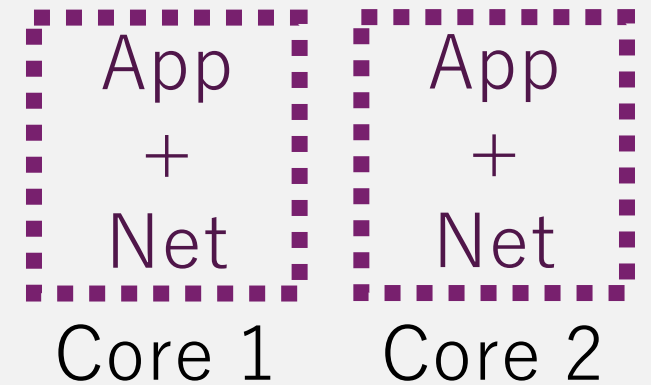


Low CPU utilization

Merge



Unified



Context switch

Issues of Existing TCP/IP Stacks

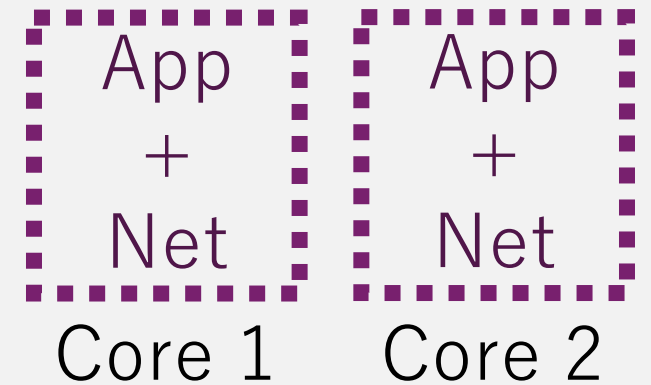
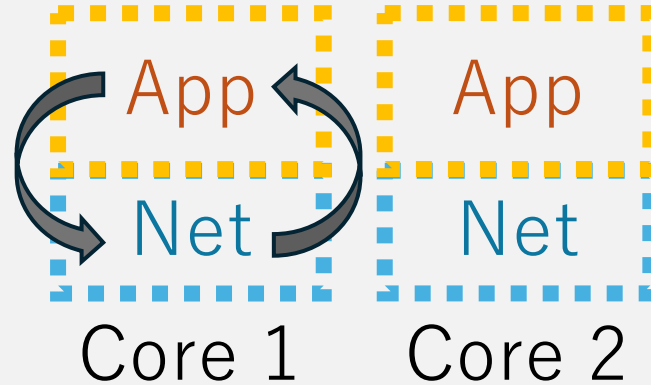
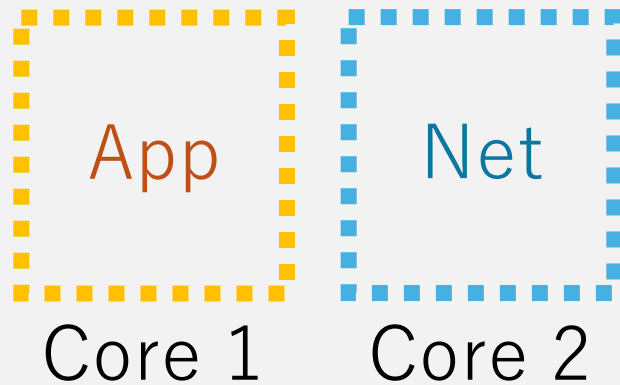
Limited Choices for CPU Core Assignment Models

A thread switching is necessary for every app/net transition

Split

Merge

Unified



Low CPU utilization

Context switch

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

A thread switching is necessary for every app/net transition

Split

Merge

Unified



Core 1

Core 2

Core 1

Core 2

Core 1

Core 2

Low CPU utilization

High transition cost

Context switch

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

The split and unified models are free from this transition cost

Split

Merge

Unified



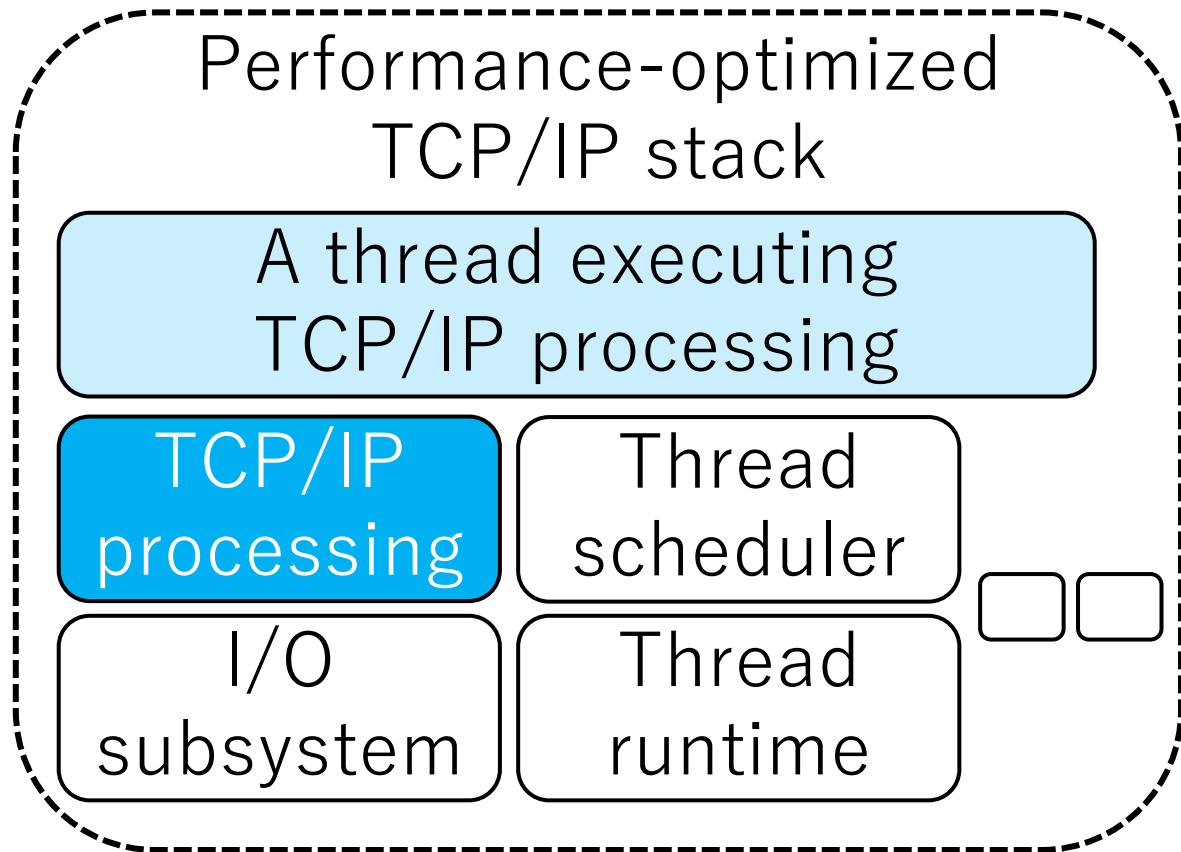
Core 1 Core 2
Low CPU utilization

Core 1 Core 2
High transition cost

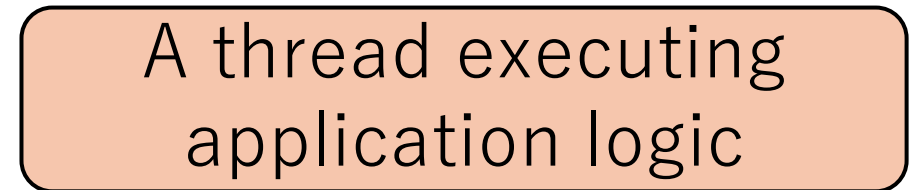
Core 1 Core 2
Context switch

Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

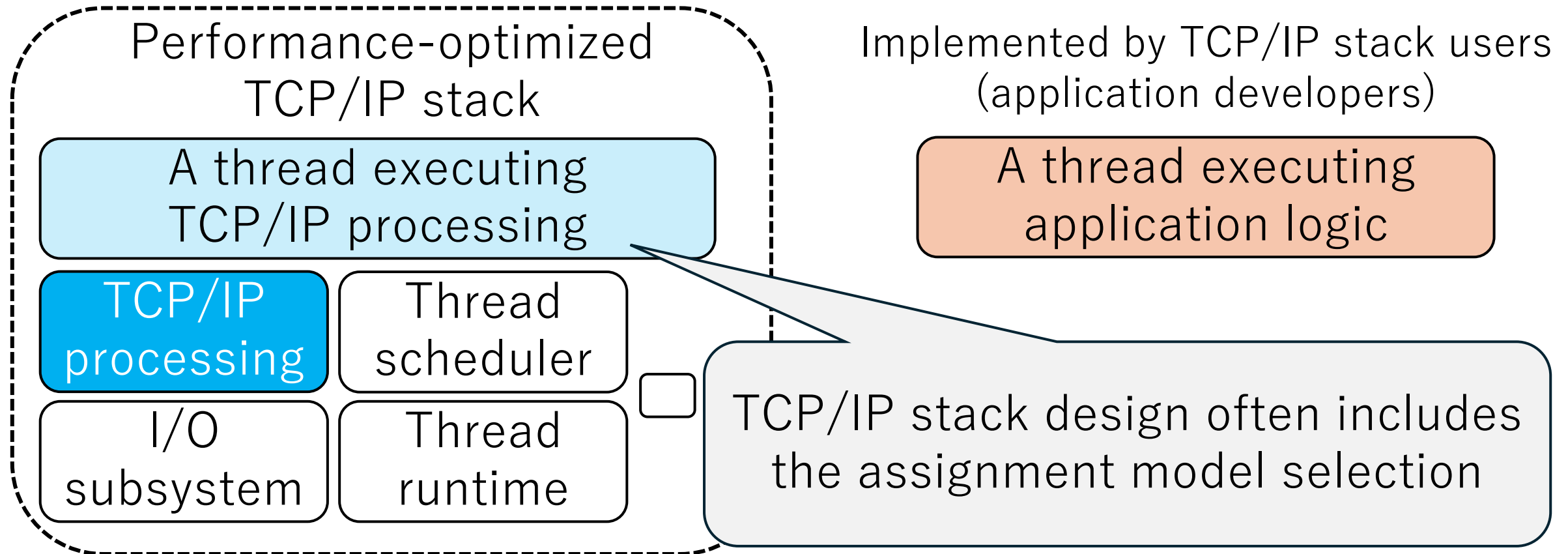


Implemented by TCP/IP stack users
(application developers)



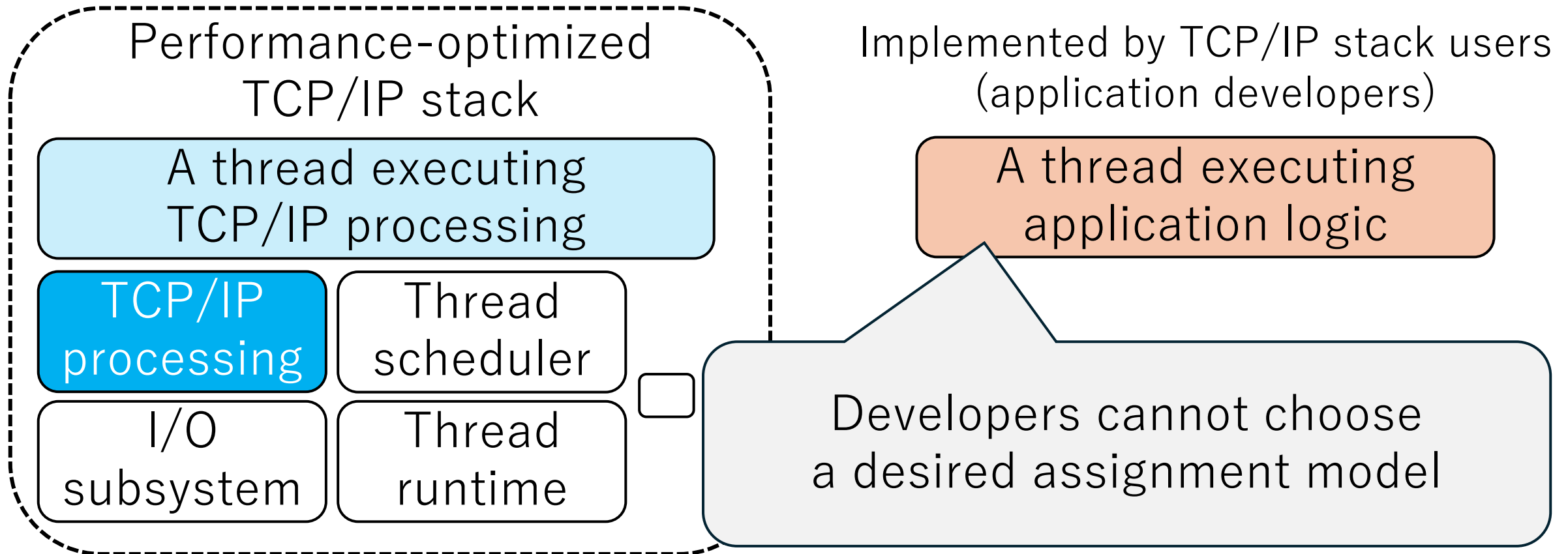
Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models



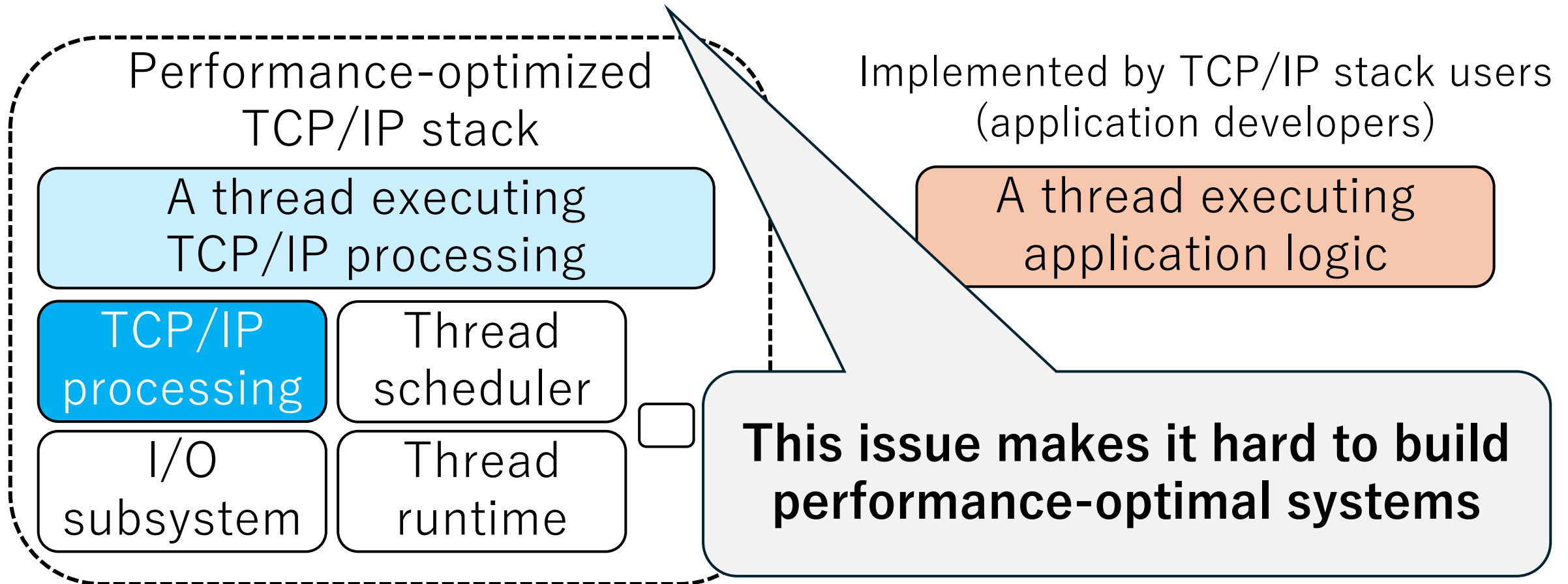
Issues of Existing TCP/IP Stacks

Limited Choices for CPU Core Assignment Models

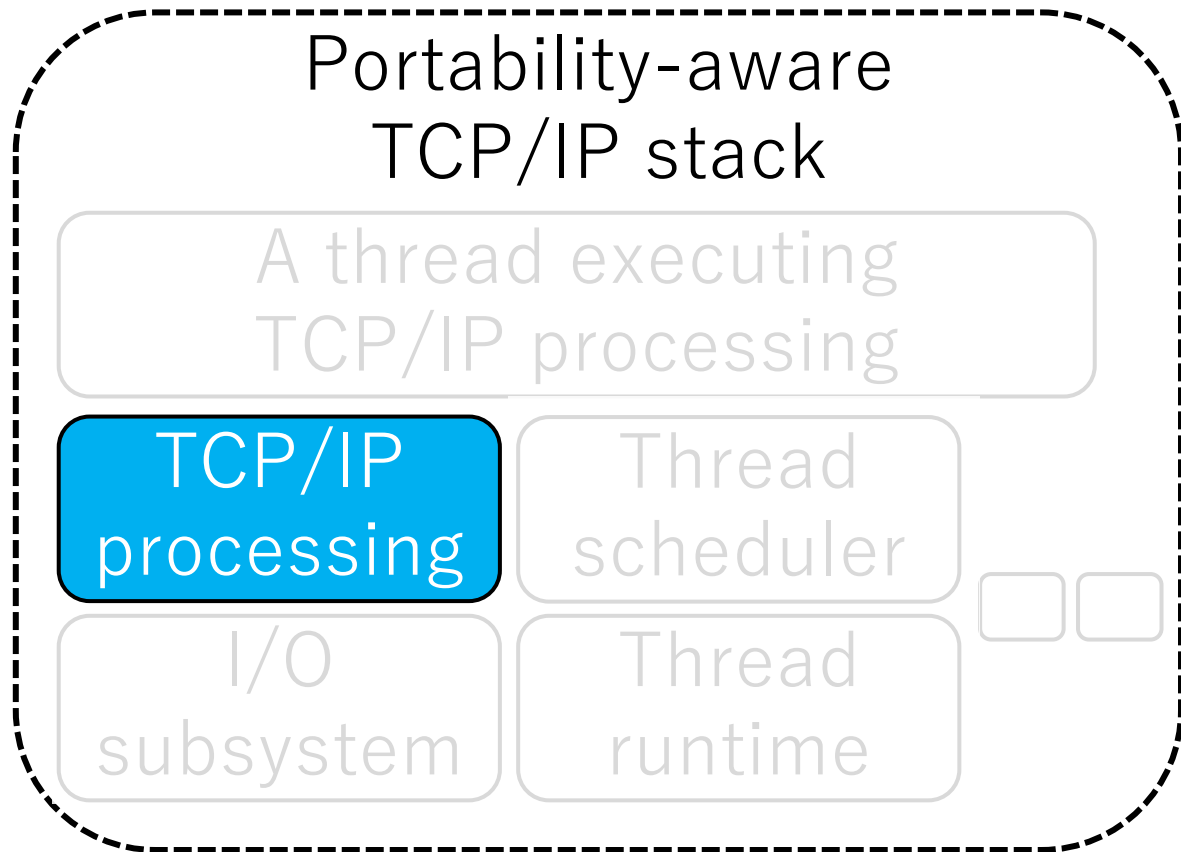


Issues of Existing TCP/IP Stacks

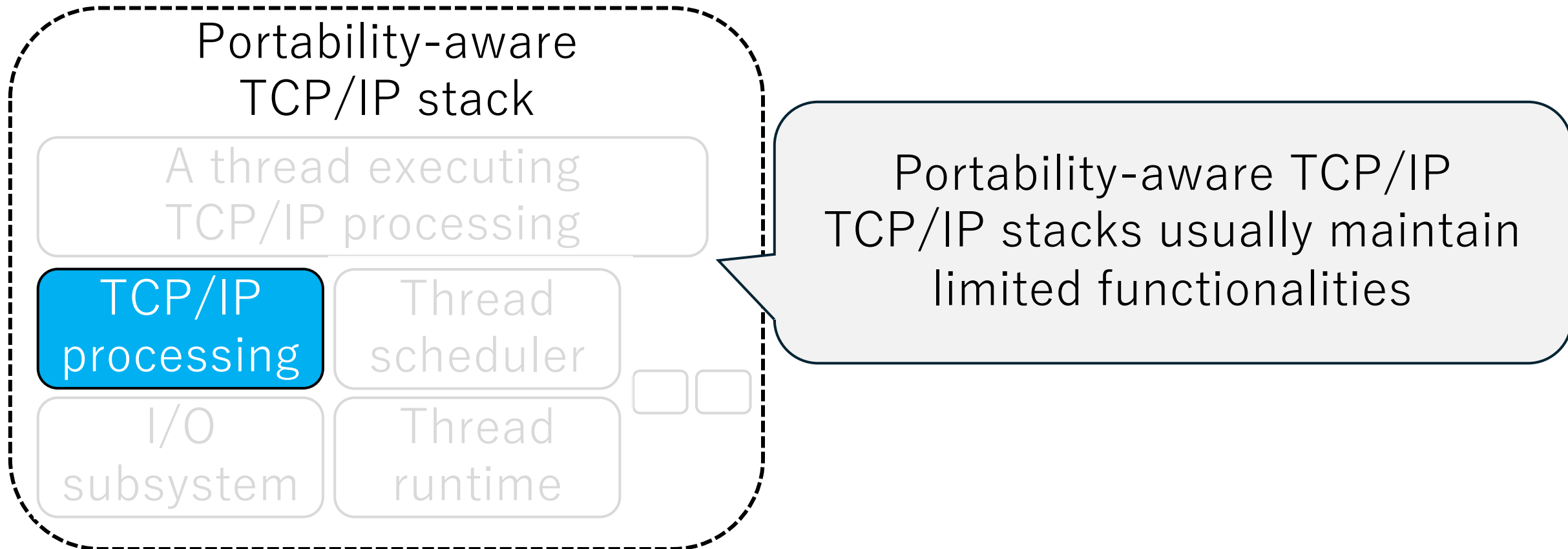
Limited Choices for CPU Core Assignment Models



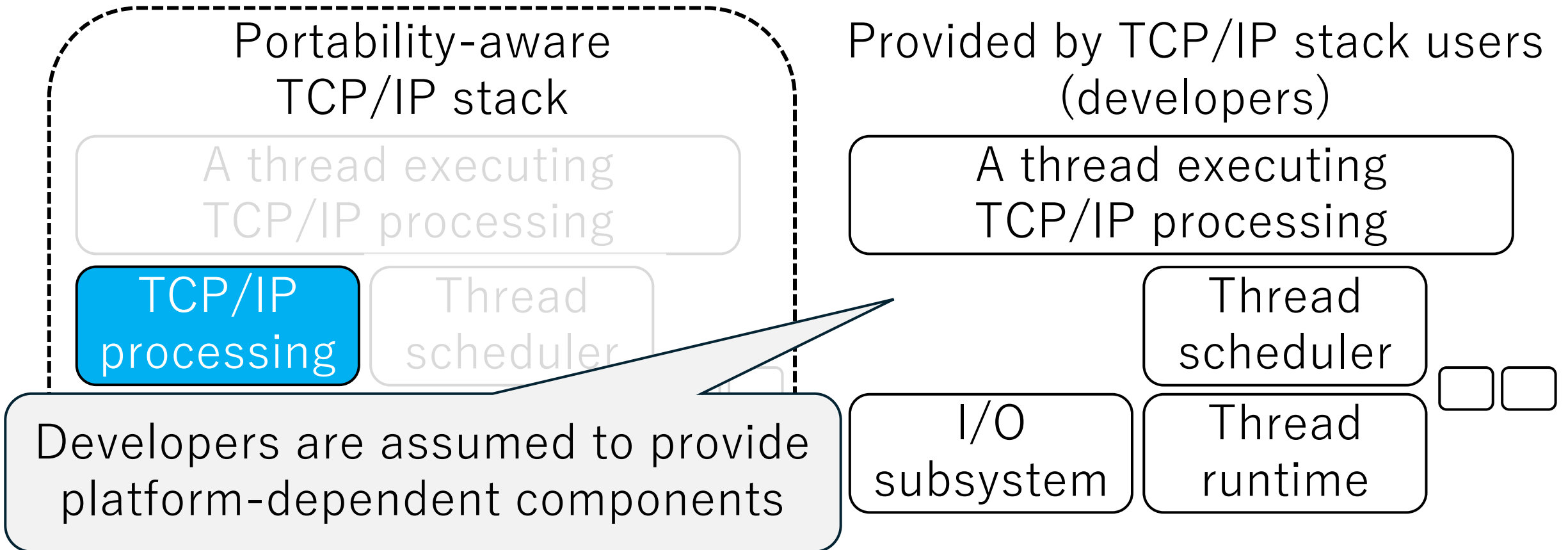
Issues of Existing TCP/IP Stacks



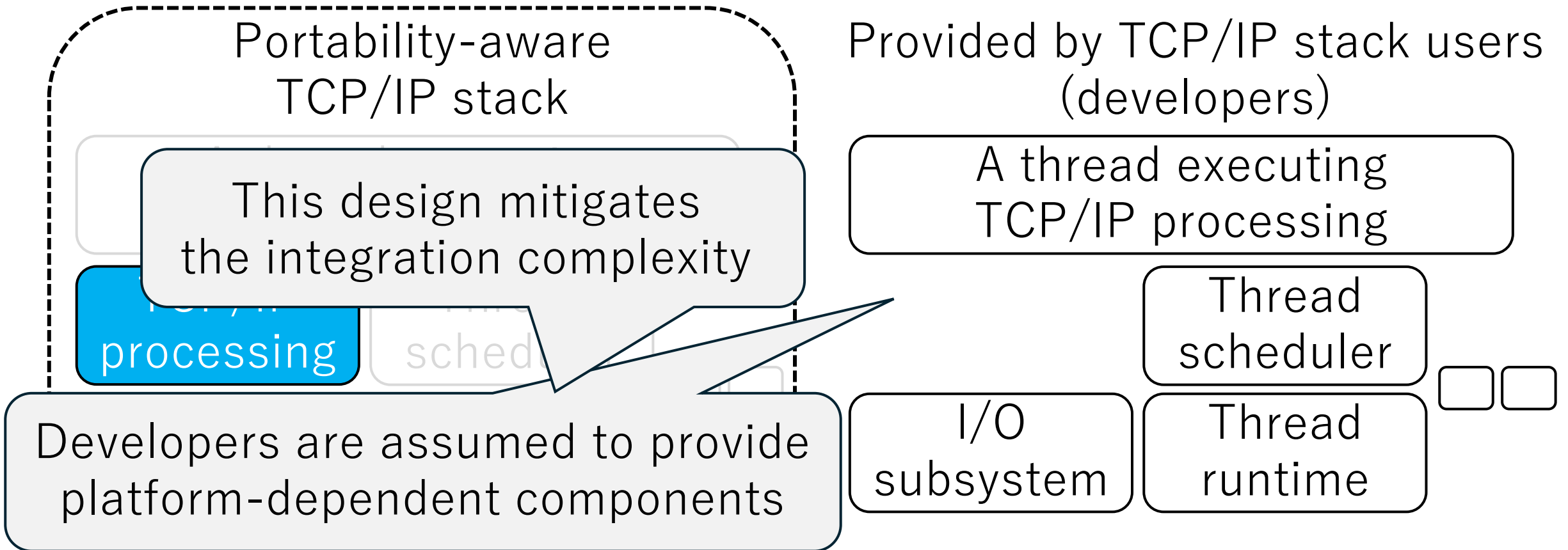
Issues of Existing TCP/IP Stacks



Issues of Existing TCP/IP Stacks

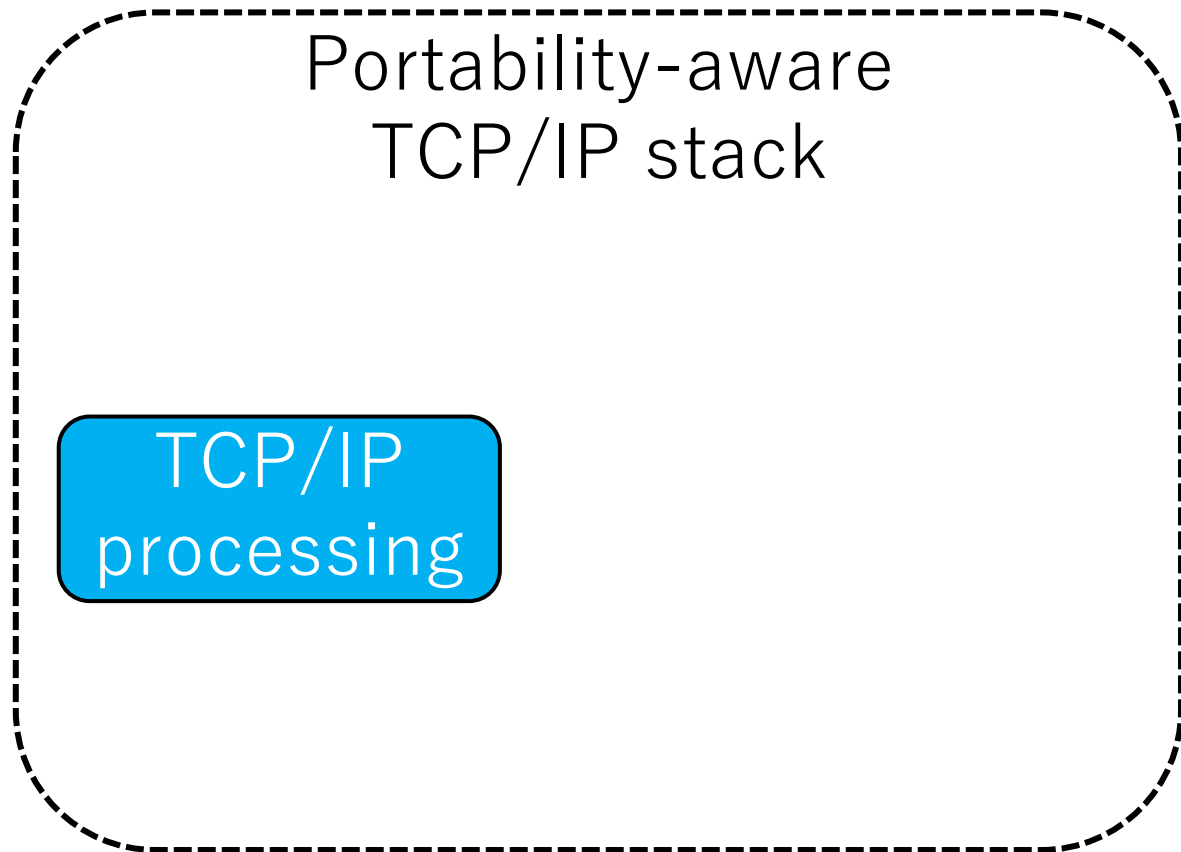


Issues of Existing TCP/IP Stacks

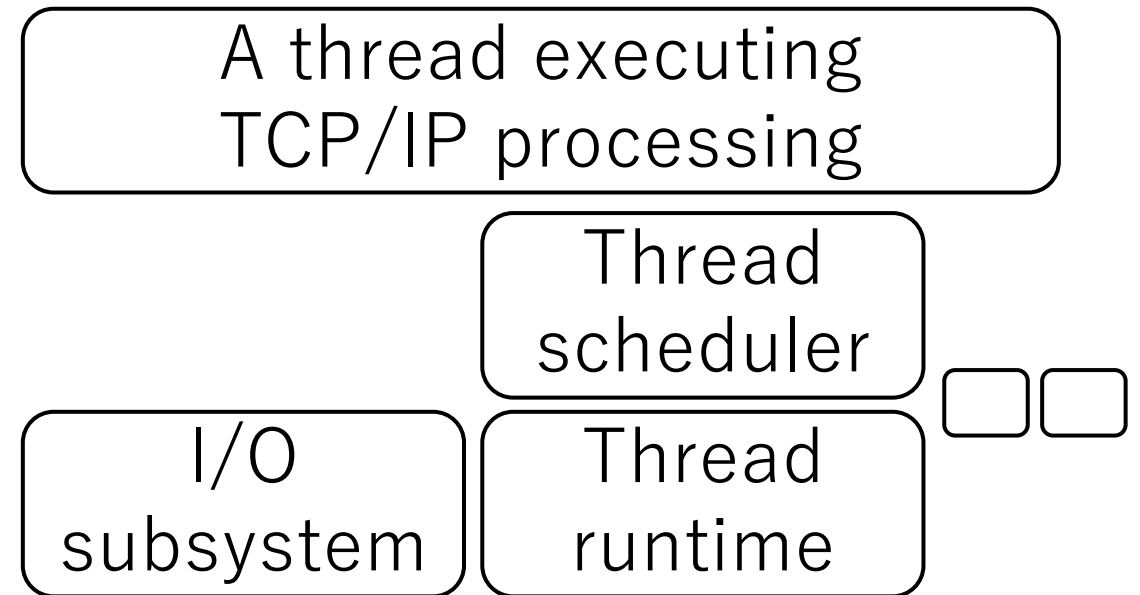


Issues of Existing TCP/IP Stacks

Not sufficiently aware of performance-critical factors

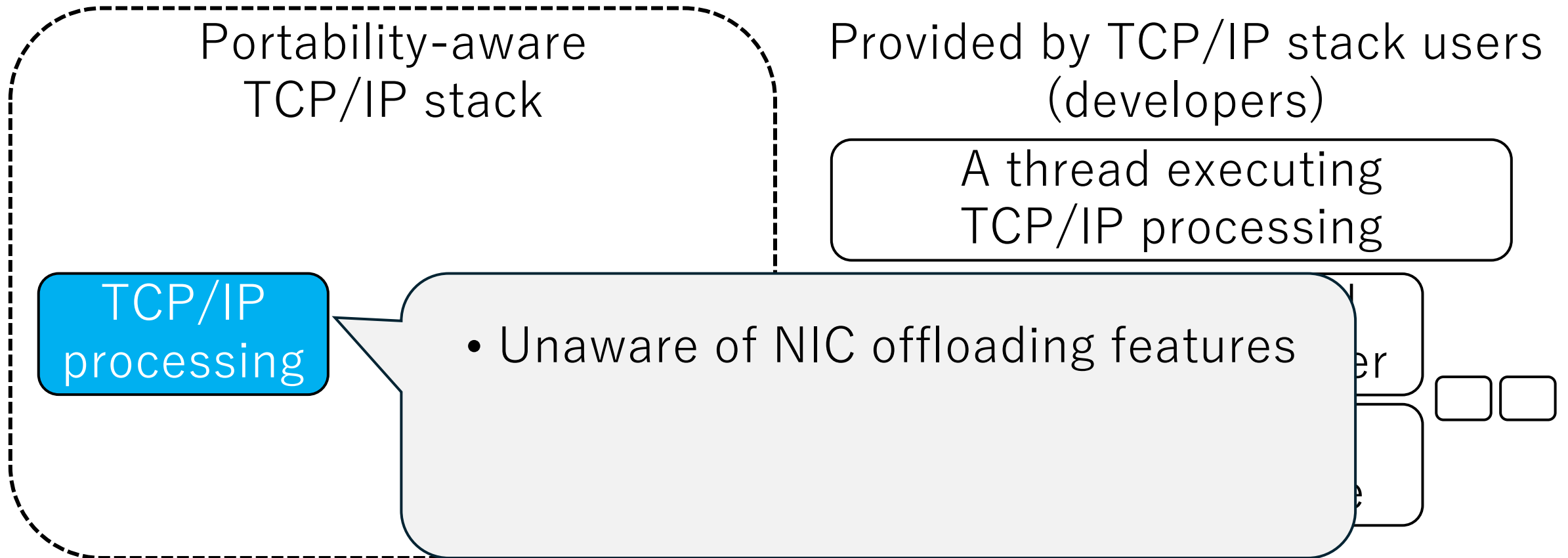


Provided by TCP/IP stack users
(developers)



Issues of Existing TCP/IP Stacks

Not sufficiently aware of performance-critical factors



Issues of Existing TCP/IP Stacks

Not sufficiently aware of performance-critical factors

Portability-aware
TCP/IP stack

Provided by TCP/IP stack users
(developers)

TCP/IP
processing

A thread executing
TCP/IP processing

- Unaware of NIC offloading features
- Lack of zero-copy I/O capability

Issues of Existing TCP/IP Stacks

Not sufficiently aware of performance-critical factors

Portability-aware
TCP/IP stack

Provided by TCP/IP stack users
(developers)

TCP/IP
processing

A thread executing
TCP/IP processing

- Unaware of NIC offloading features
- Lack of zero-copy I/O capability
- Lack of multi-core scalability

Issues of Existing TCP/IP Stacks

Not sufficiently aware of performance-critical factors

Portability

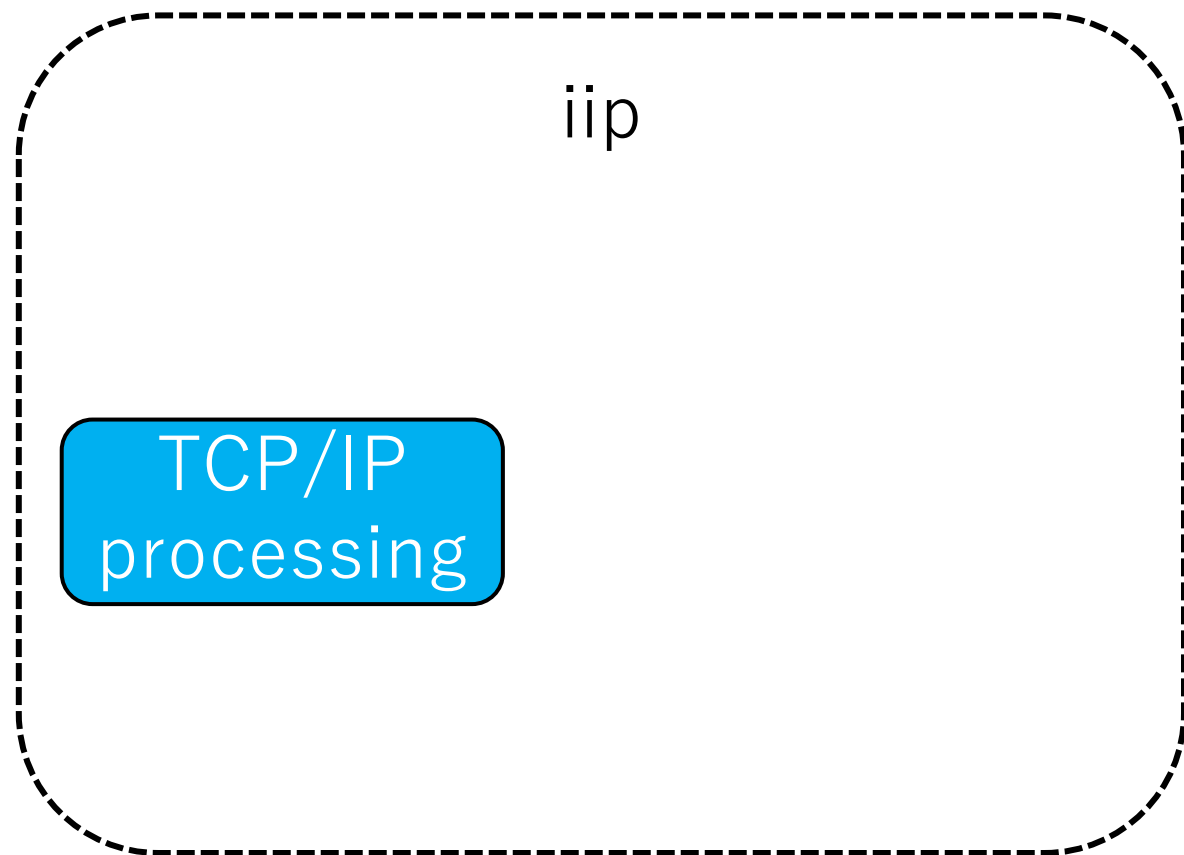
Provided by TCP/IP stack users

Portability-aware TCP/IP stacks cannot achieve good performance in various workloads

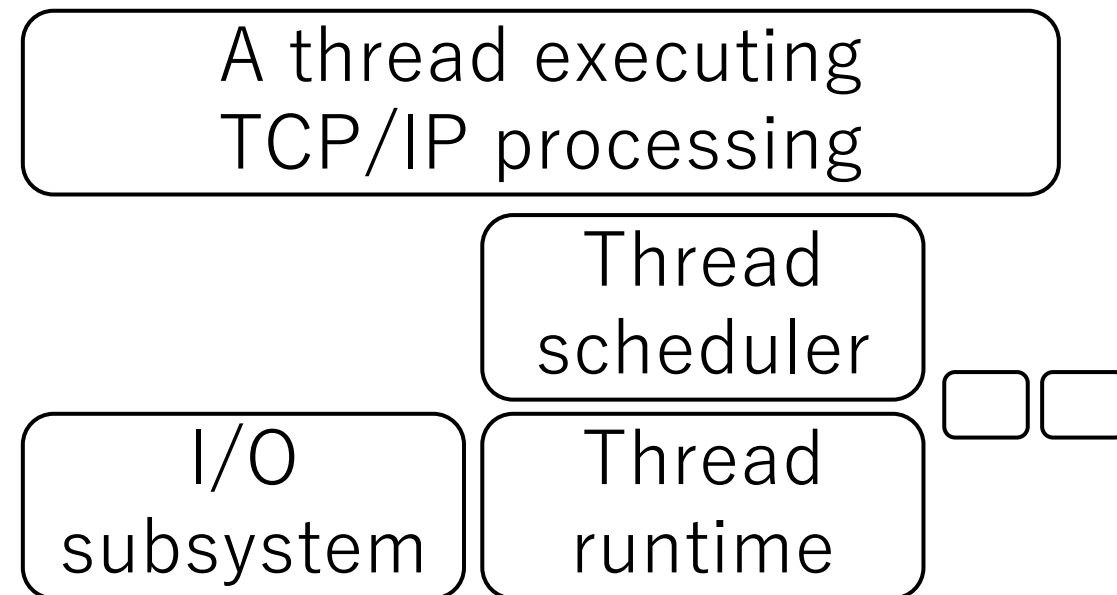
TCP/IP processing

- Unaware of NIC offloading features
- Lack of zero-copy I/O capability
- Lack of multi-core scalability

iip



Provided by TCP/IP stack users
(developers)



iip

iip only implements the TCP/IP processing functionality to minimize the chance for causing functionality conflicts

iip

TCP/IP
processing

Provided by TCP/IP stack users
(developers)

A thread executing
TCP/IP processing

Thread
scheduler

I/O
subsystem

Thread
runtime

iip

Platform-dependent functionalities are assumed to be provided by developers through the API

Quality conflicts

Provided by TCP/IP stack users (developers)

A thread executing TCP/IP processing

Thread scheduler

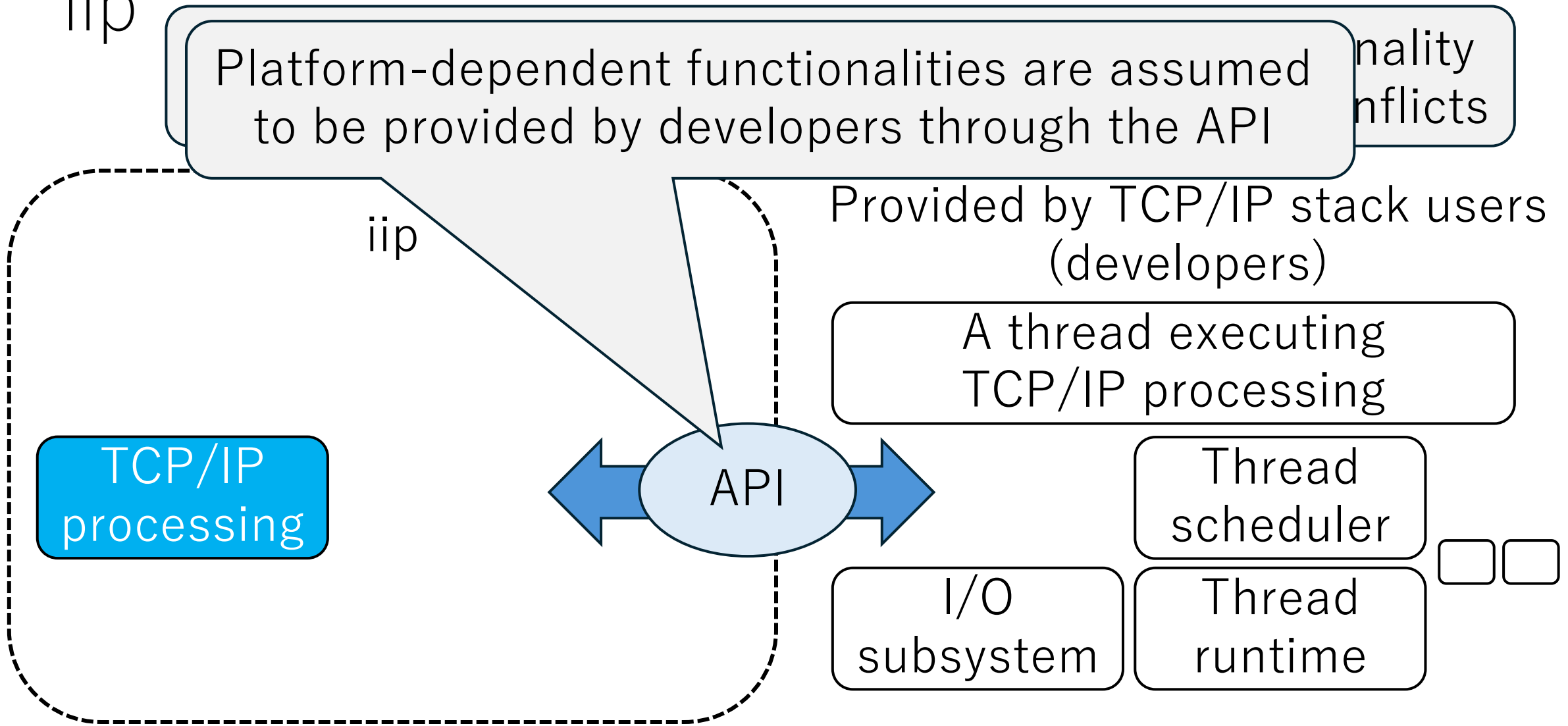
Thread runtime

I/O subsystem

API

TCP/IP processing

iip



iip

Platform-dependent functionalities are assumed to be provided by developers through the API

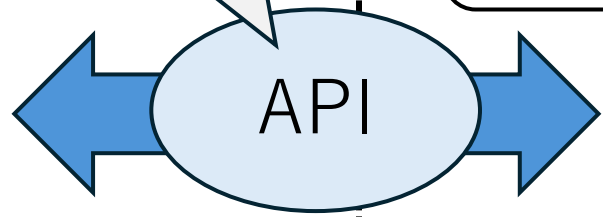
Quality conflicts

Provided by TCP/IP stack users (developers)

No dependency on CPUs, NICs, OSes, libraries, and compilers

A thread executing TCP/IP processing

TCP/IP processing



Thread scheduler

I/O subsystem

Thread runtime



iip

Platform-dependent functionalities are assumed to be provided by developers through the API

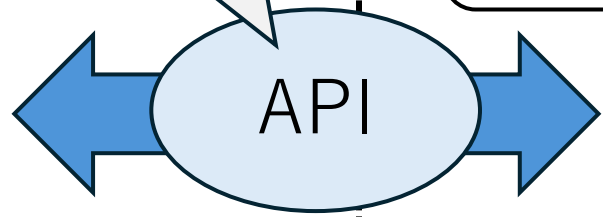
Quality conflicts

Provided by TCP/IP stack users (developers)

No dependency on CPUs, NICs, OSes, libraries, and compilers

A thread executing TCP/IP processing

TCP/IP processing



Thread scheduler

C89 / C++98 compliant for old and future compilers

I/O subsystem

Thread runtime



iip

Platform-dependent functionalities are assumed to be provided by developers through the API

Quality conflicts

Provided by TCP/IP stack users (developers)

iip

No dependency on CPUs, NICs, OSes, libraries, and compilers

A thread executing

iip pays attention to performance-critical factors

TCP/IP processing

API

C89 / C++98 compliant for old and future compilers

iip

Platform-dependent functionalities are assumed to be provided by developers through the API

Quality conflicts

Provided by TCP/IP stack users (developers)

No dependency on CPUs, NICs, OSes, libraries, and compilers

TCP/IP processing

API

A thread executing
iip pays attention to performance-critical factors
• uses NIC offloading features

C89 / C++98 compliant for old and future compilers

iip

Platform-dependent functionalities are assumed to be provided by developers through the API

Quality conflicts

Provided by TCP/IP stack users (developers)

iip

No dependency on CPUs, NICs, OSes, libraries, and compilers

A thread executing

iip pays attention to performance-critical factors
• uses NIC offloading features
• supports zero-copy I/O

TCP/IP processing

API

C89 / C++98 compliant for old and future compilers

iip

Platform-dependent functionalities are assumed to be provided by developers through the API

Quality
Conflicts

Provided by TCP/IP stack users (developers)

iip

No dependency on CPUs, NICs, OSes, libraries, and compilers

A thread executing

- iip pays attention to performance-critical factors
- uses NIC offloading features
- supports zero-copy I/O
- aware of multi-core scalability

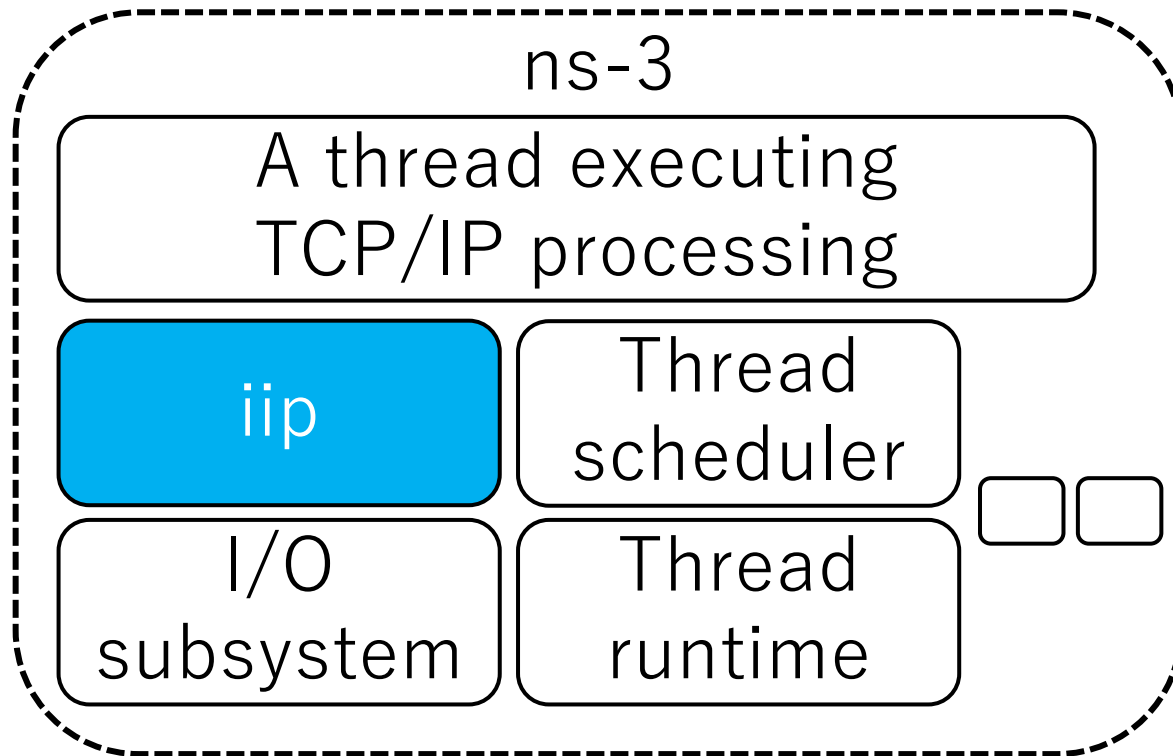
TCP/IP processing

API

C89 / C++98 compliant for old and future compilers

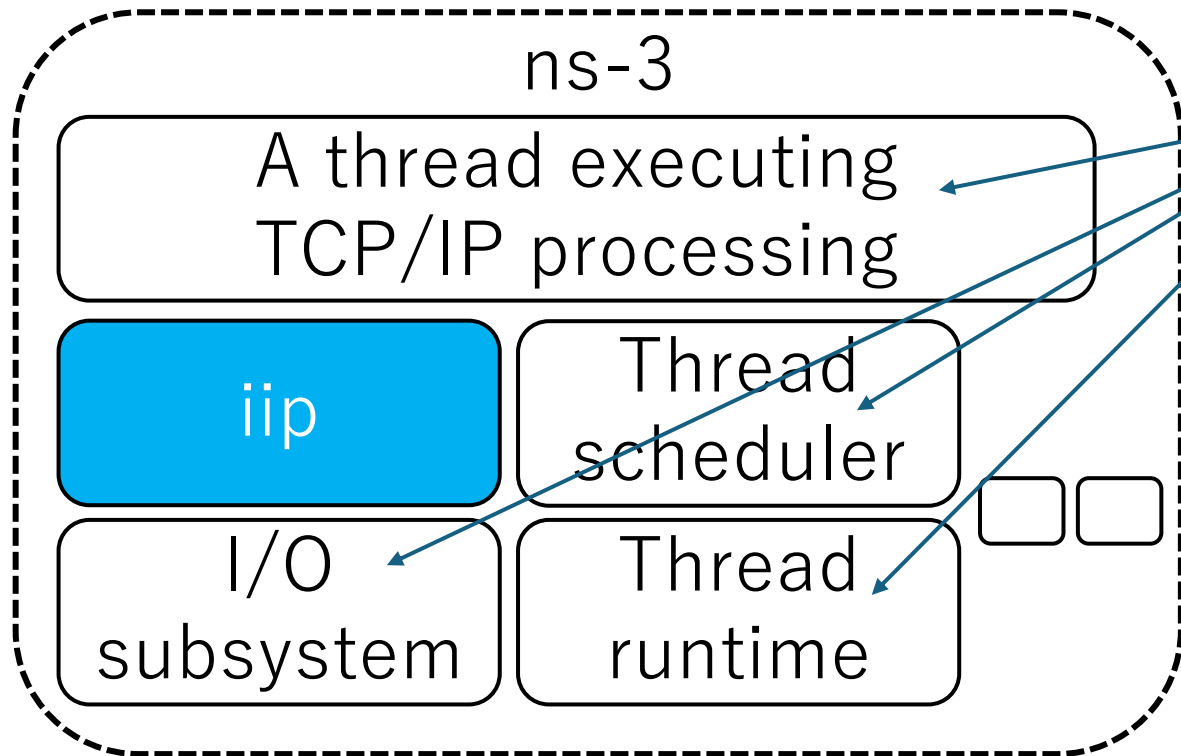
Integratability Benefits

- iip can be integrated into the ns-3 simulator written in C++
 - <https://github.com/yasukata/iip-ns>



Integratability Benefits

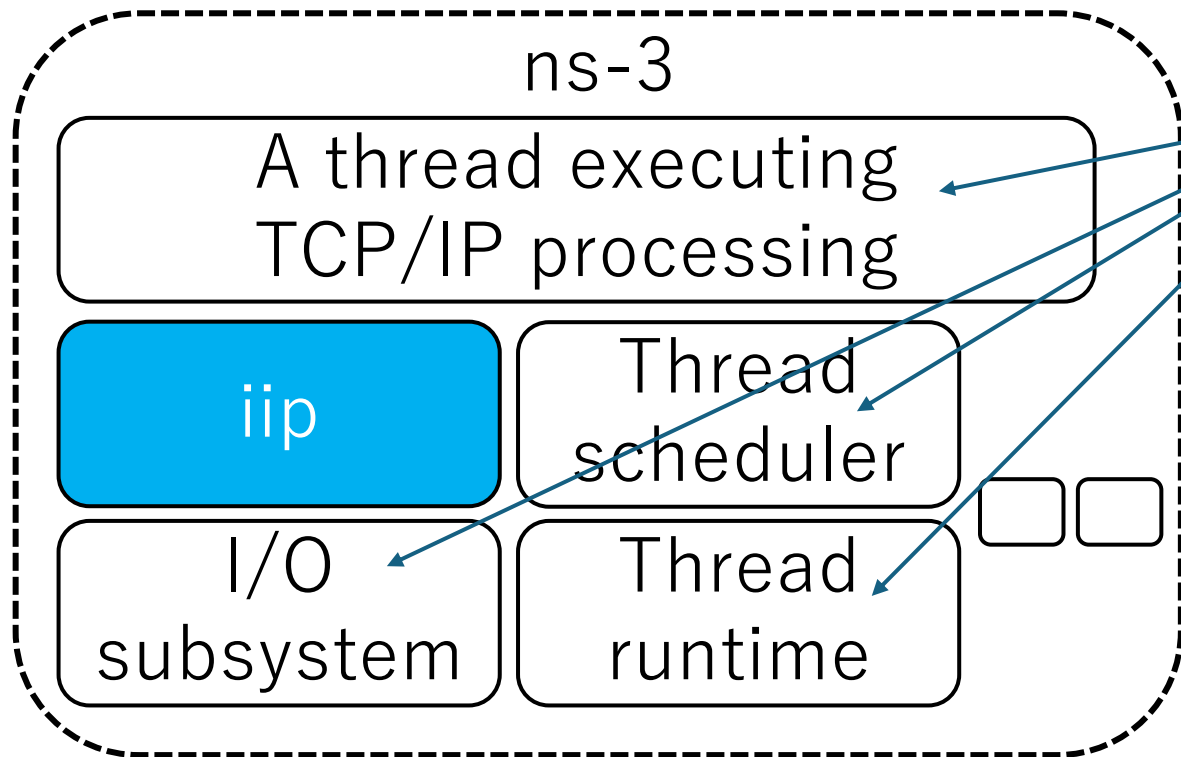
- iip can be integrated into the ns-3 simulator written in C++
 - <https://github.com/yasukata/iip-ns>



These functionalities are specific to ns-3

Integratability Benefits

- iip can be integrated into the ns-3 simulator written in C++
 - <https://github.com/yasukata/iip-ns>



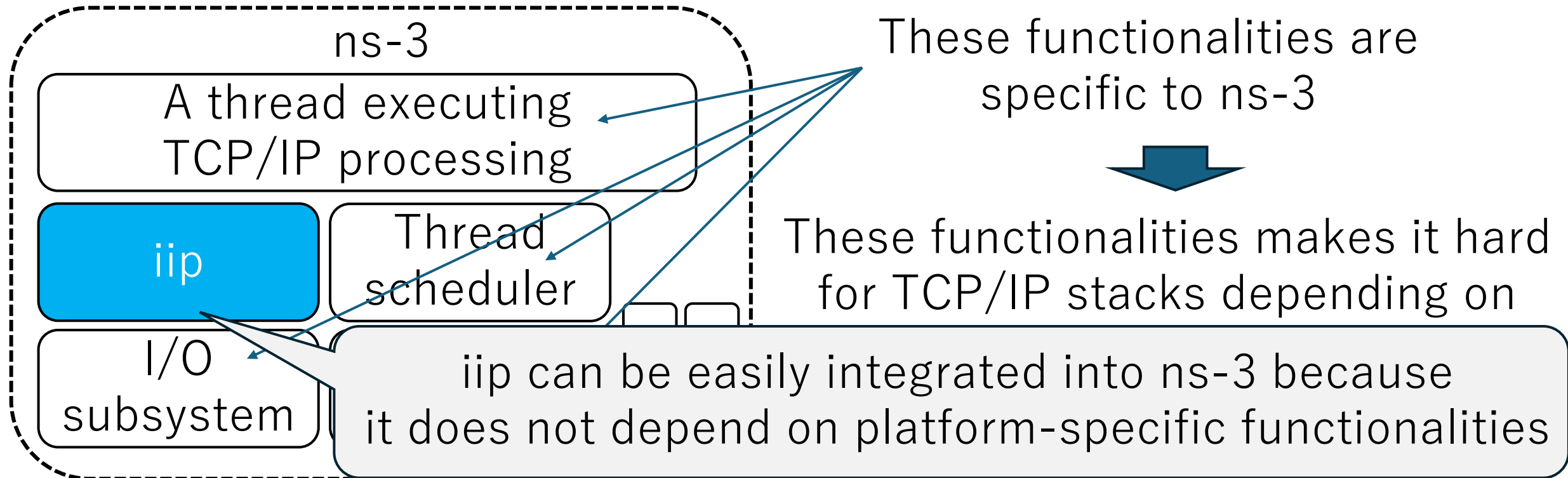
These functionalities are specific to ns-3



These functionalities makes it hard for TCP/IP stacks depending on common OS features to run on ns-3

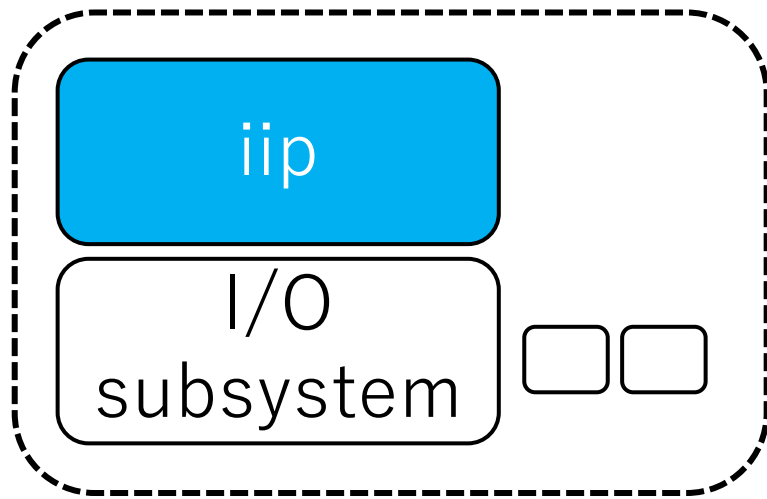
Integratability Benefits

- iip can be integrated into the ns-3 simulator written in C++
 - <https://github.com/yasukata/iip-ns>



Integratability Benefits

- iip runs on various I/O backends, including but not limited to:
 - DPDK (Linux): <https://github.com/yasukata/iip-dpdk>
 - AF_XDP (Linux): https://github.com/yasukata/iip-af_xdp
 - netmap (Linux/FreeBSD): <https://github.com/yasukata/iip-netmap>

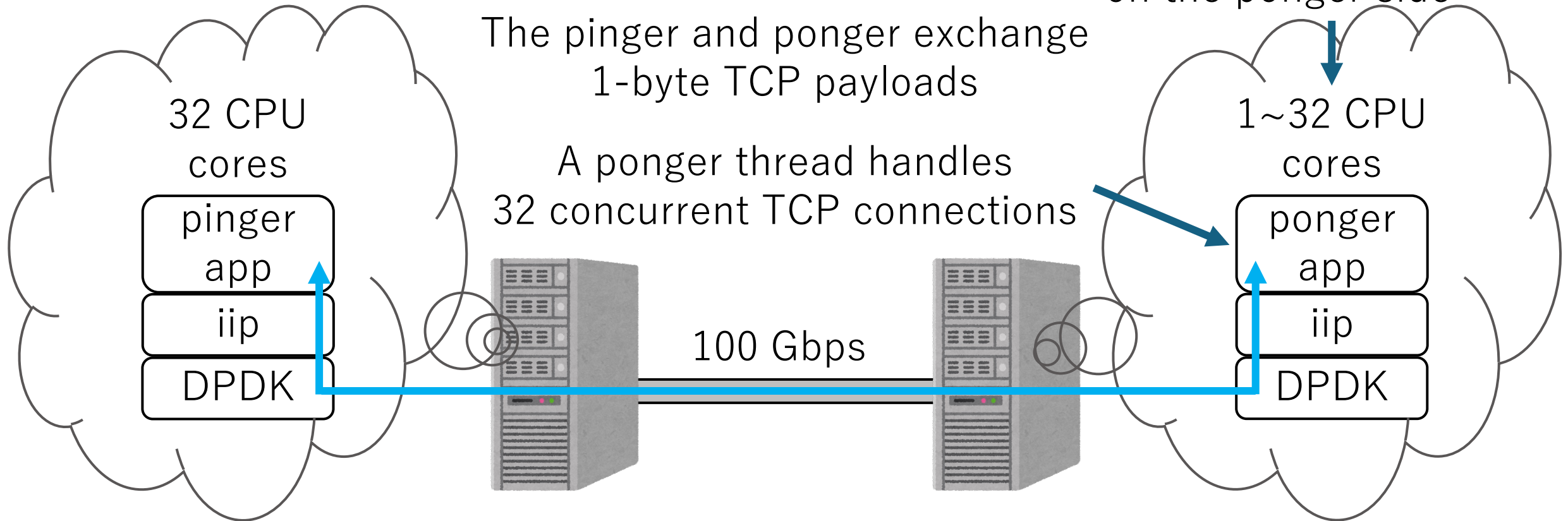


iip does not depend on specific I/O subsystems

Evaluation: Small Message Exchange

- TCP ping-pong workload

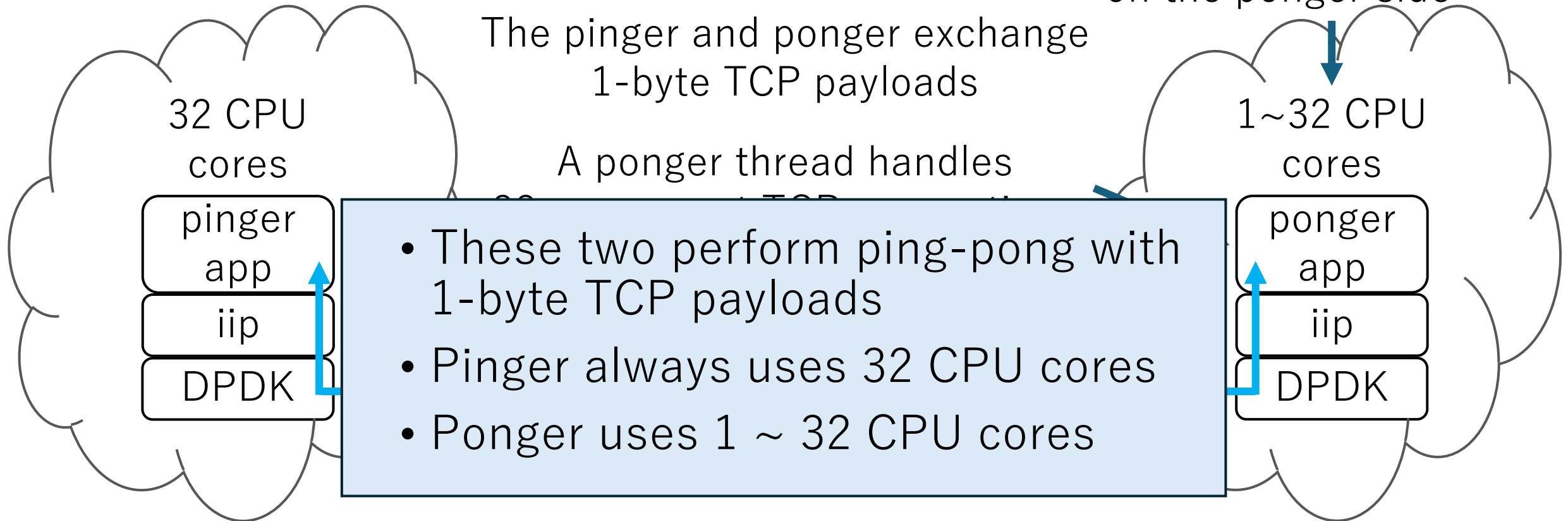
We have tested while changing the number of CPU cores on the ponger side



Evaluation: Small Message Exchange

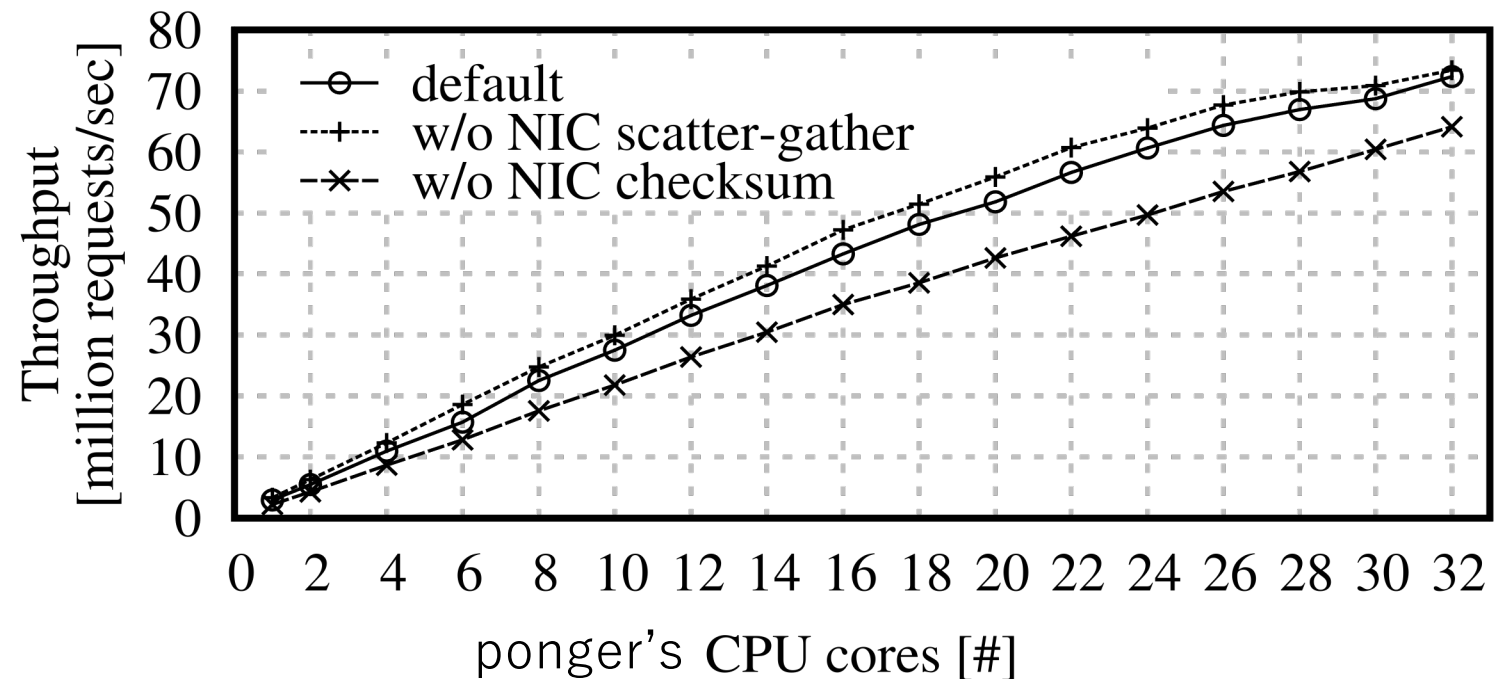
- TCP ping-pong workload

We have tested while changing the number of CPU cores on the ponger side



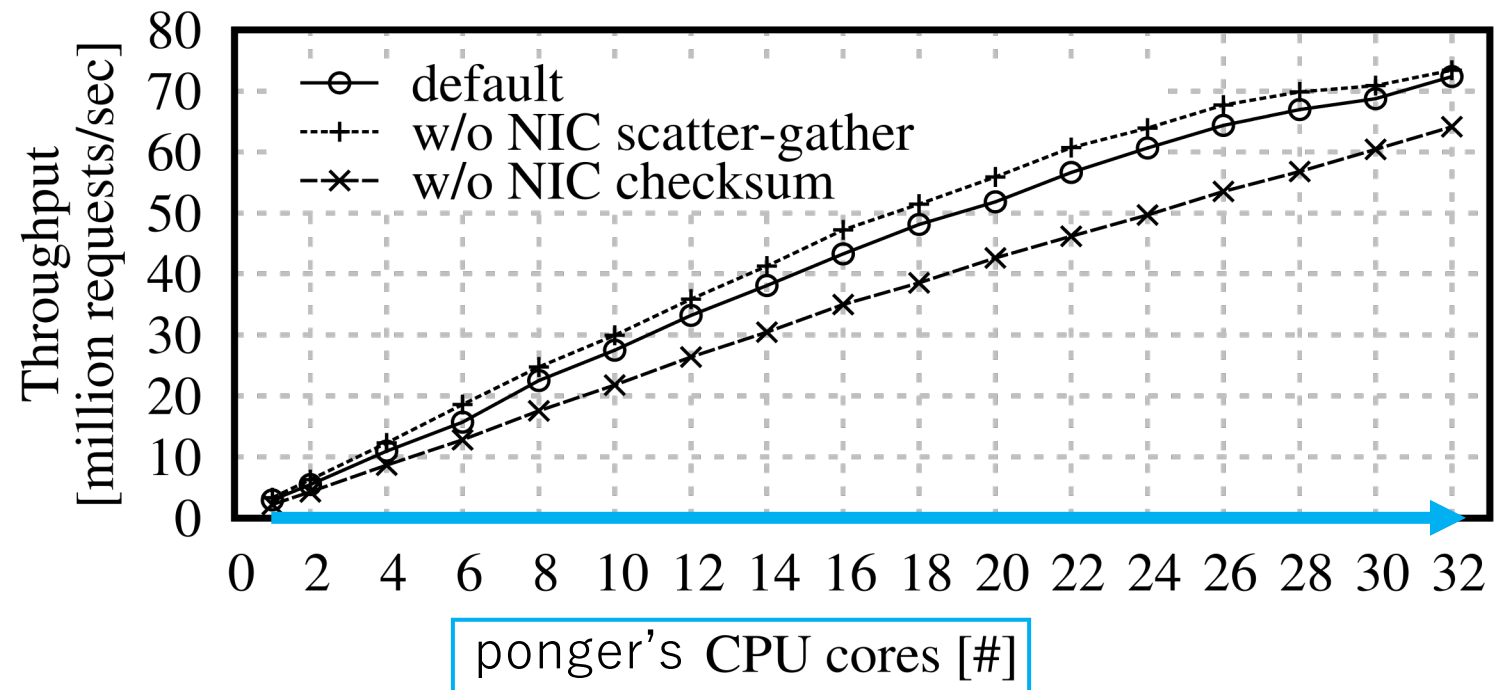
Evaluation: Small Message Exchange

- The pinger and ponger apps exchange 1-byte TCP payloads



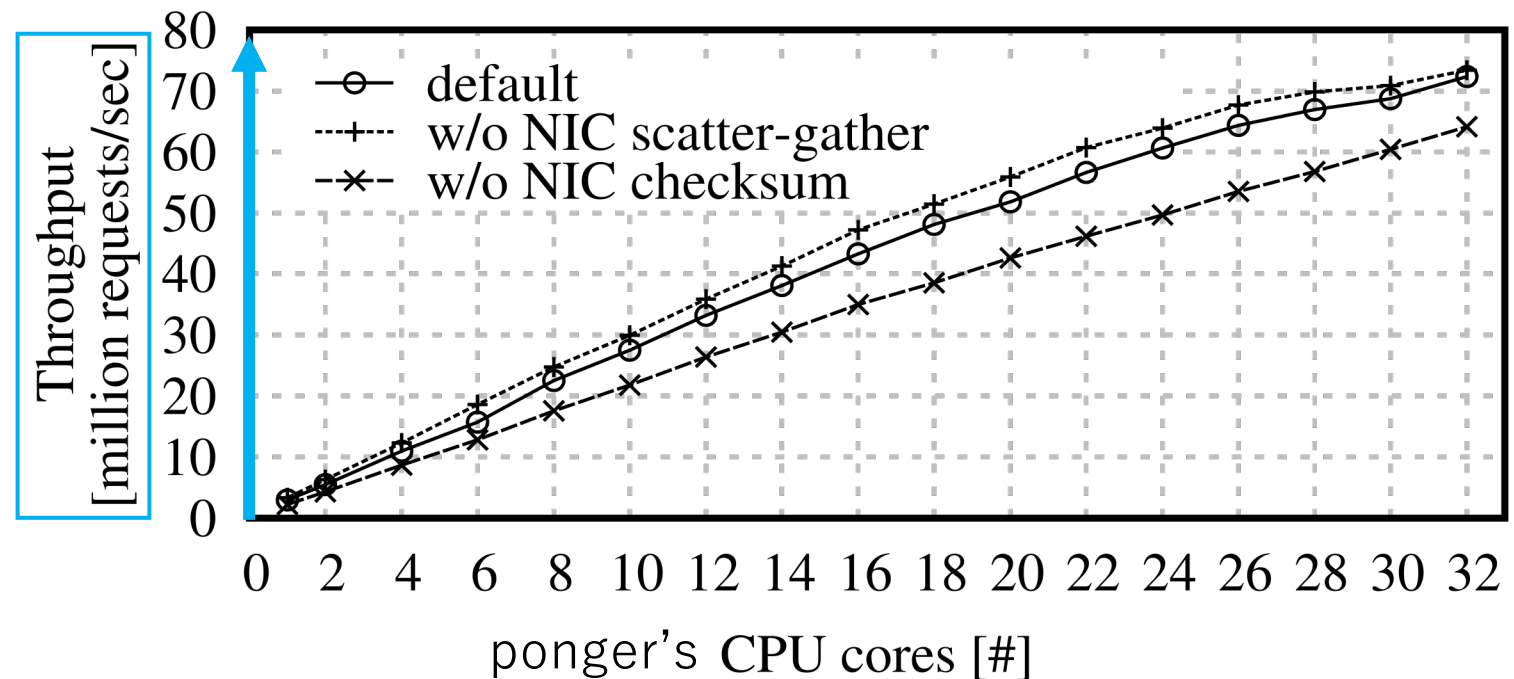
Evaluation: Small Message Exchange

- The pinger and ponger apps exchange 1-byte TCP payloads



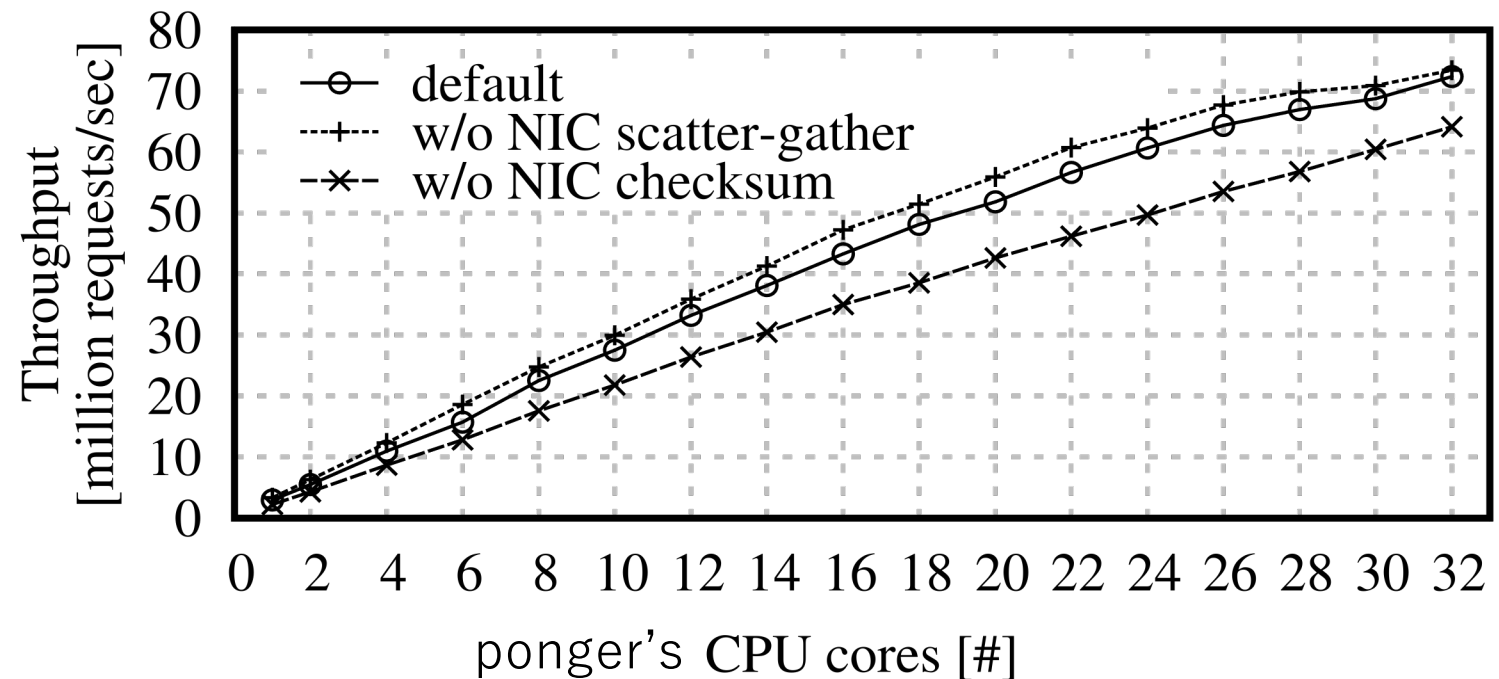
Evaluation: Small Message Exchange

- The pinger and ponger apps exchange 1-byte TCP payloads



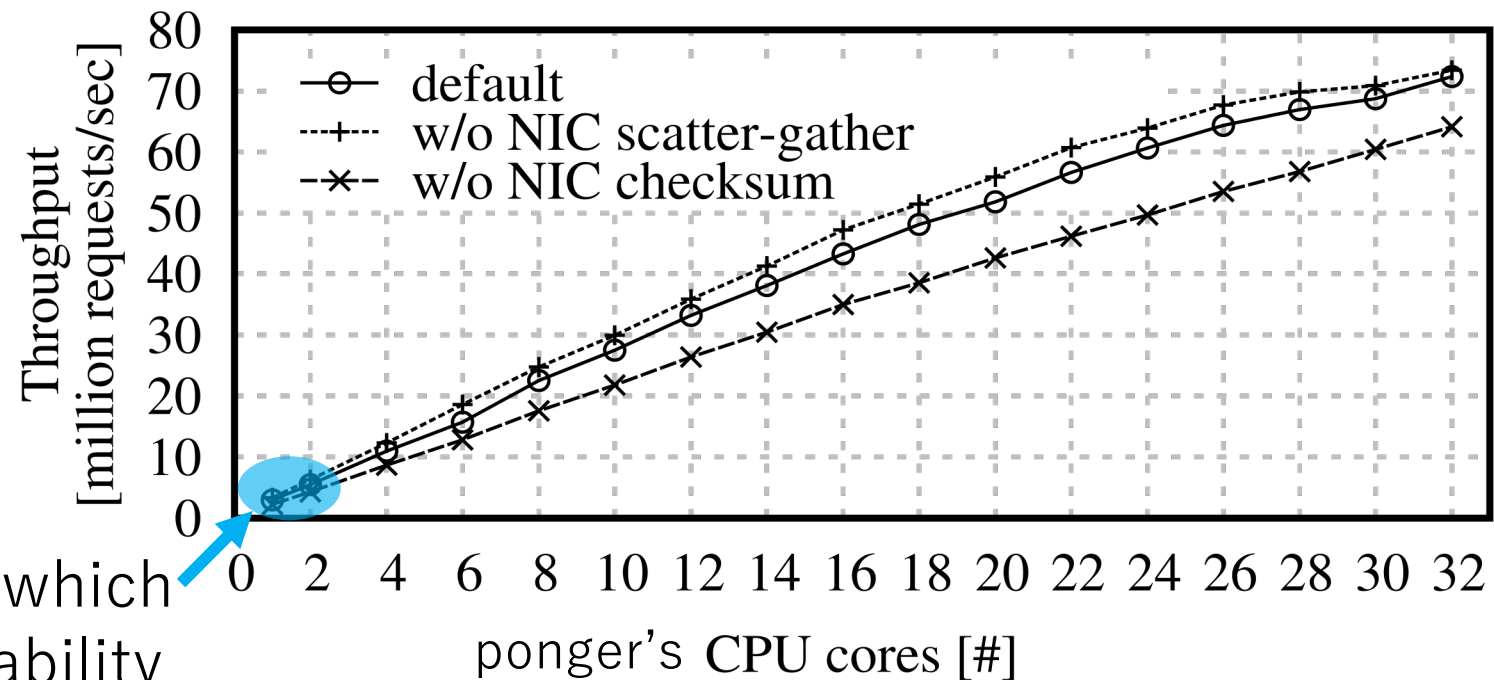
Evaluation: Small Message Exchange

- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability



Evaluation: Small Message Exchange

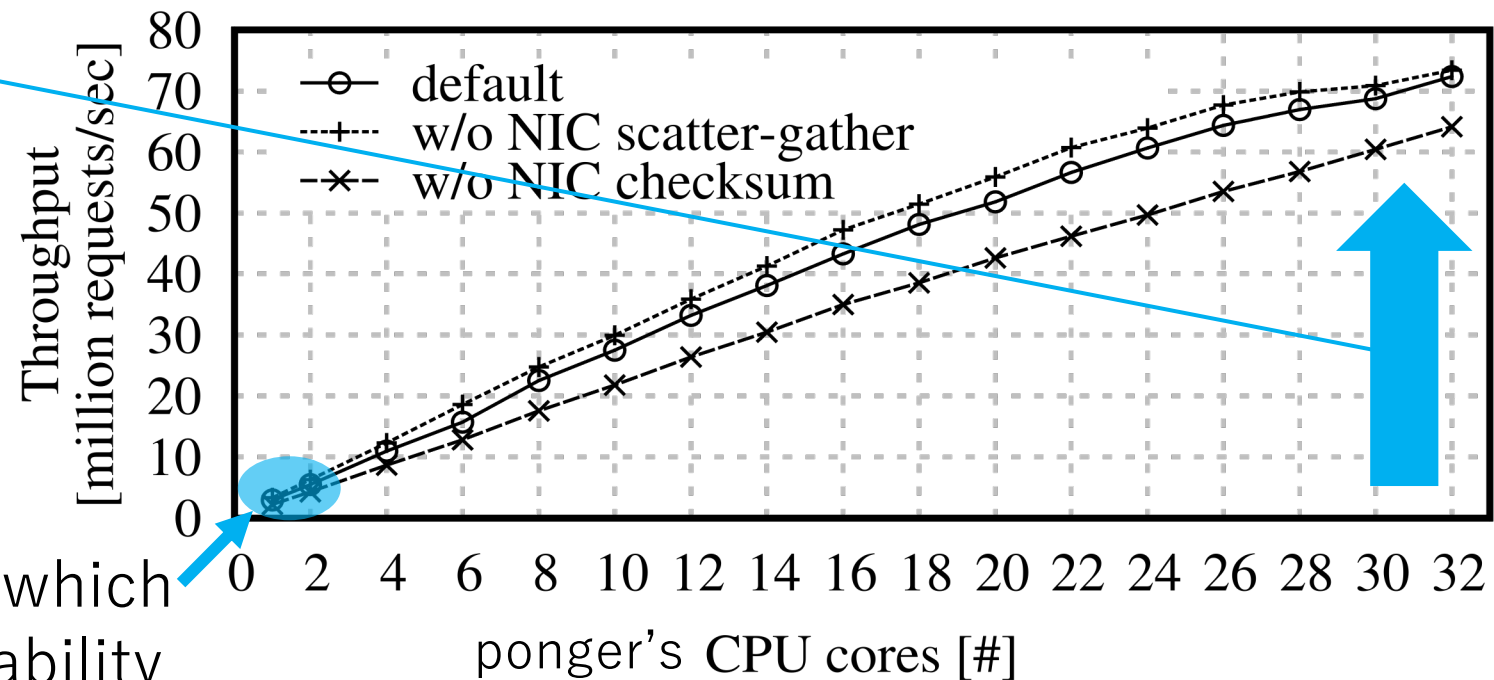
- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability



Performance of TCP/IP stacks which are unaware of multi-core scalability

Evaluation: Small Message Exchange

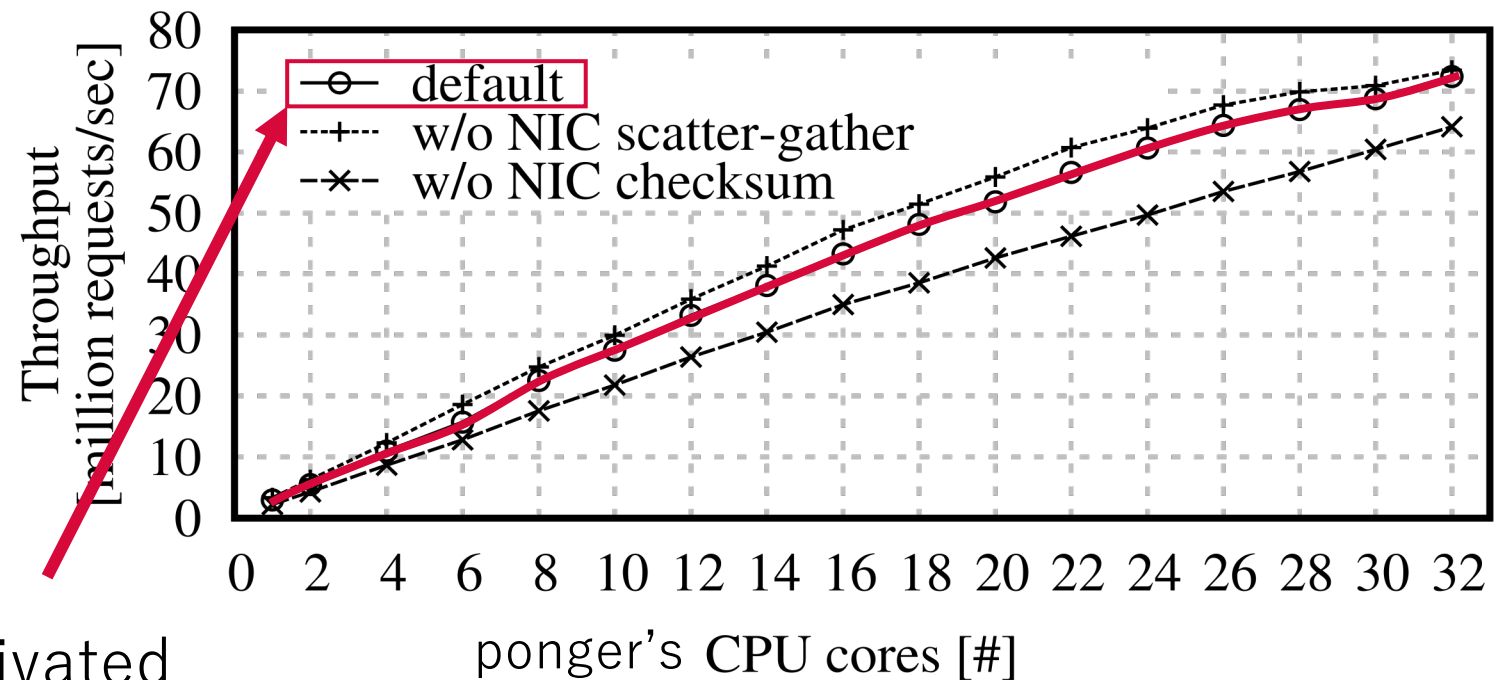
- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability



Performance of TCP/IP stacks which are unaware of multi-core scalability

Evaluation: Small Message Exchange

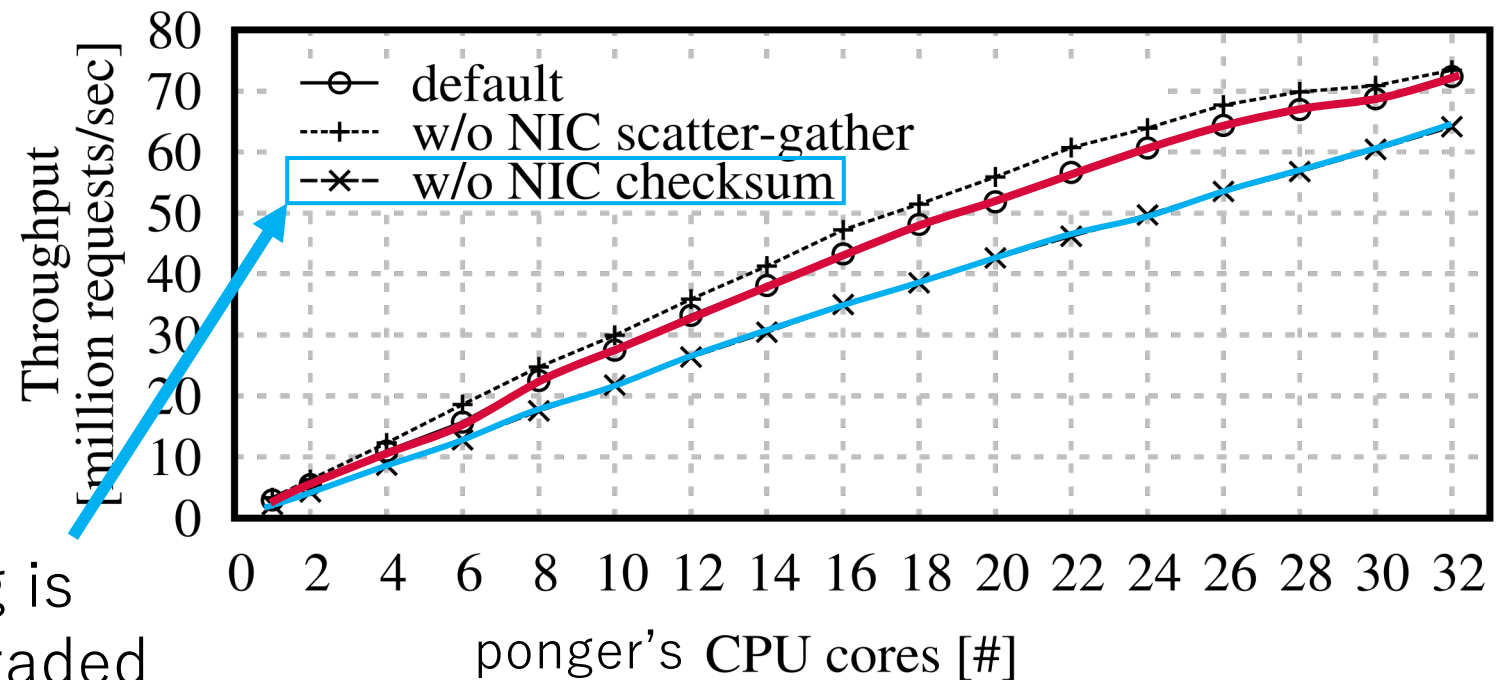
- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability



All NIC offloading features and zero-copy transmission are activated

Evaluation: Small Message Exchange

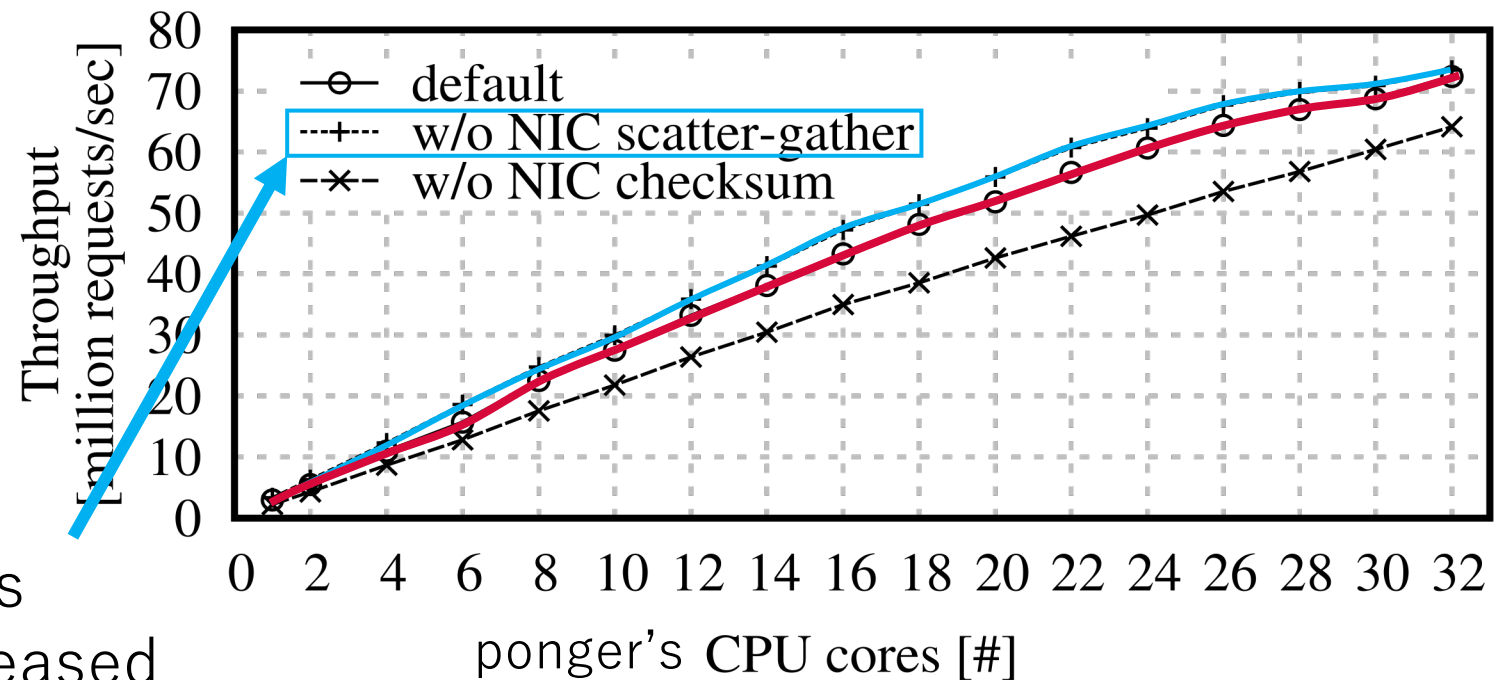
- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability
 - Checksum offloading



When NIC checksum offloading is deactivated, throughput is degraded

Evaluation: Small Message Exchange

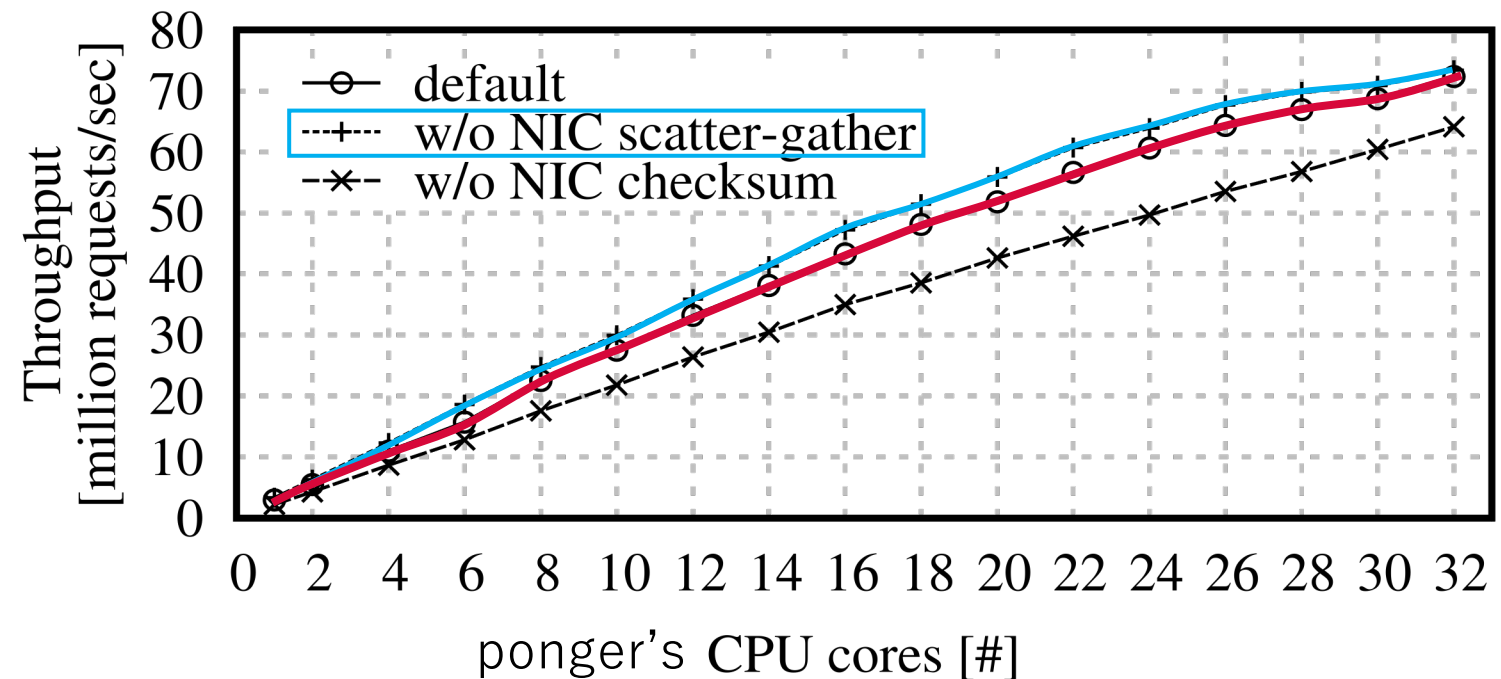
- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability
 - Checksum offloading



When zero-copy transmission is deactivated, throughput is increased

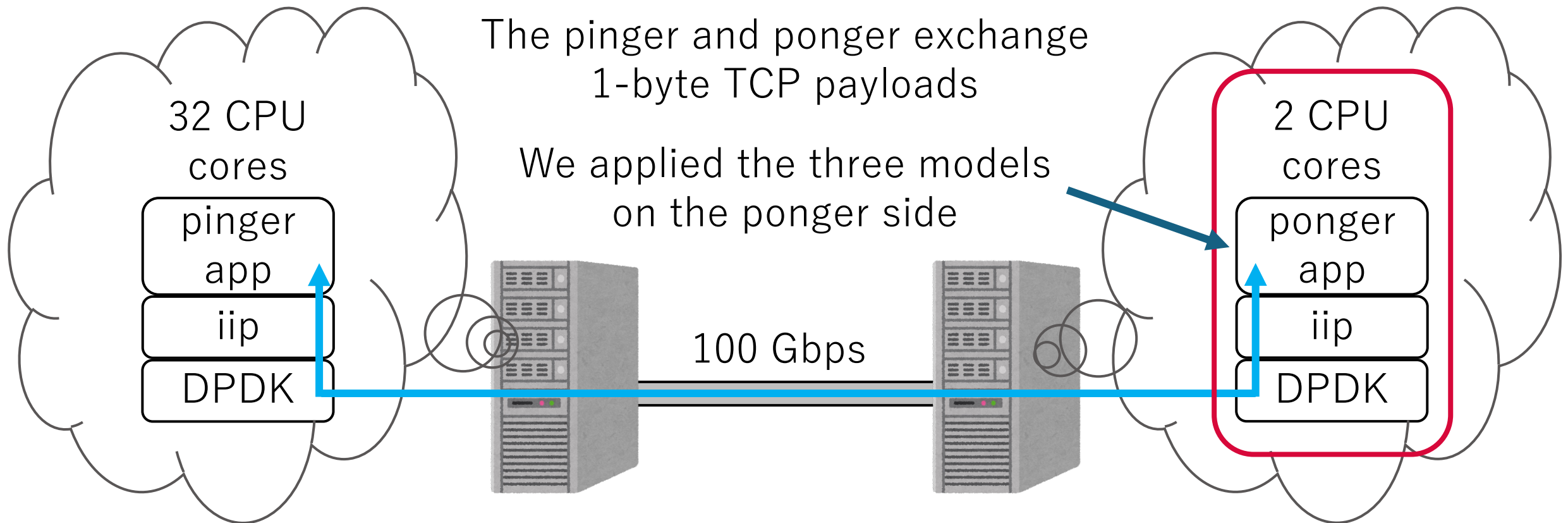
Evaluation: Small Message Exchange

- The pinger and ponger apps exchange 1-byte TCP payloads
- Performance factors
 - Multi-core scalability
 - Checksum offloading
- Tips
 - For small messages, copy is faster than zero-copy transmission



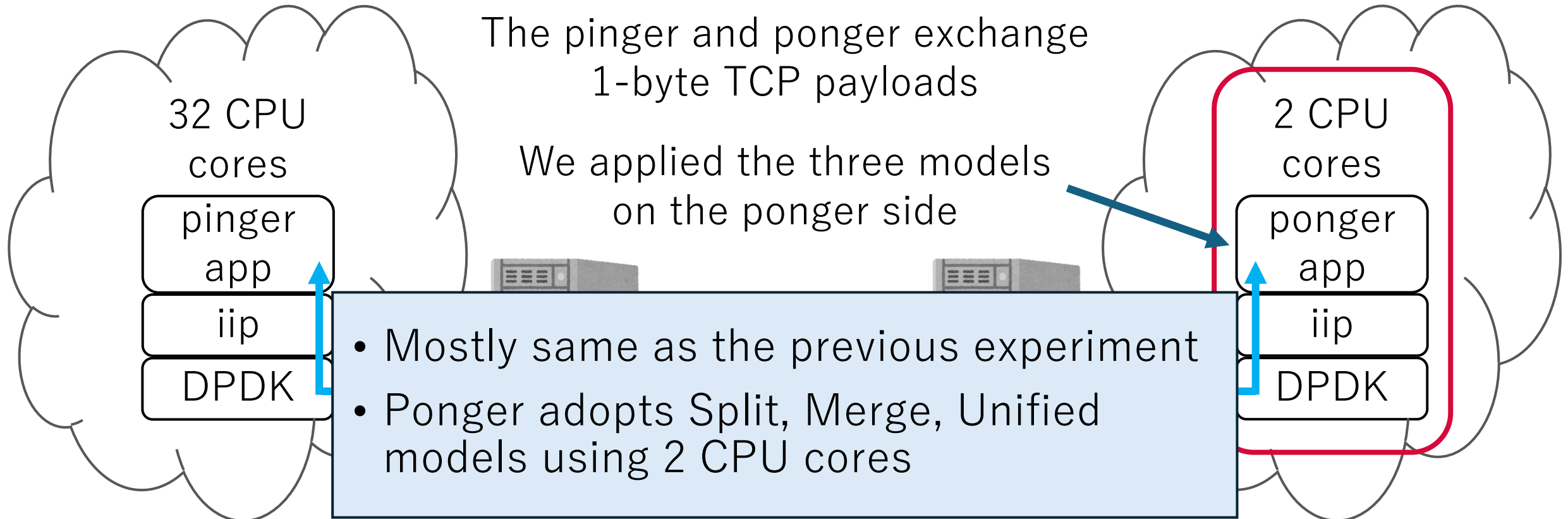
Evaluation: CPU Core Assignment Models

- TCP ping-pong workload



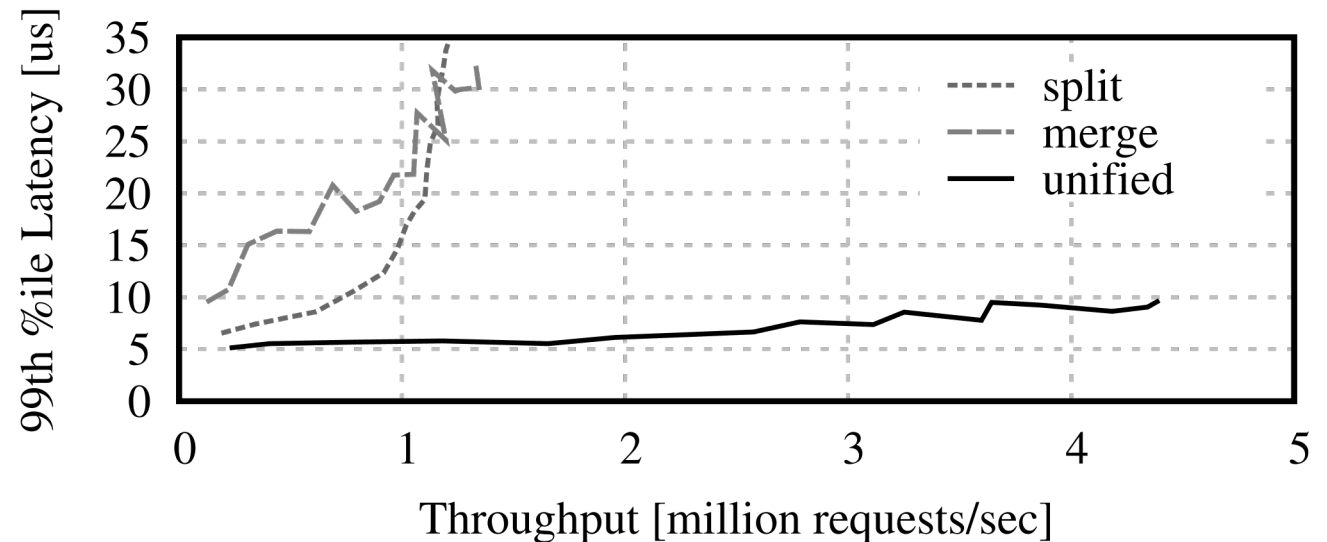
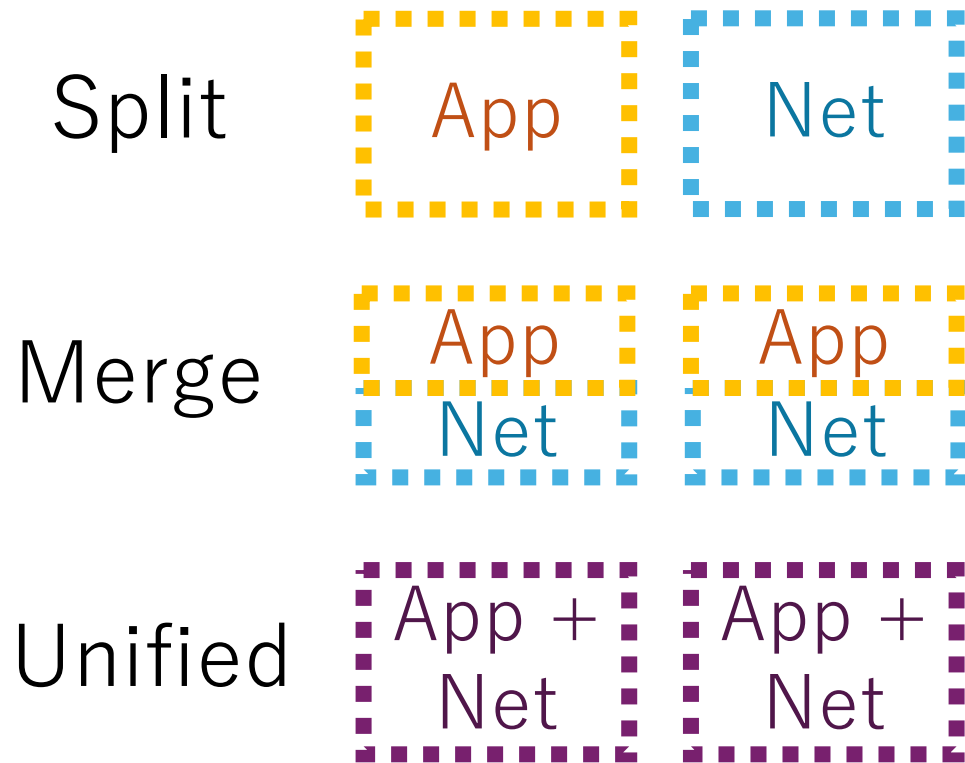
Evaluation: CPU Core Assignment Models

- TCP ping-pong workload



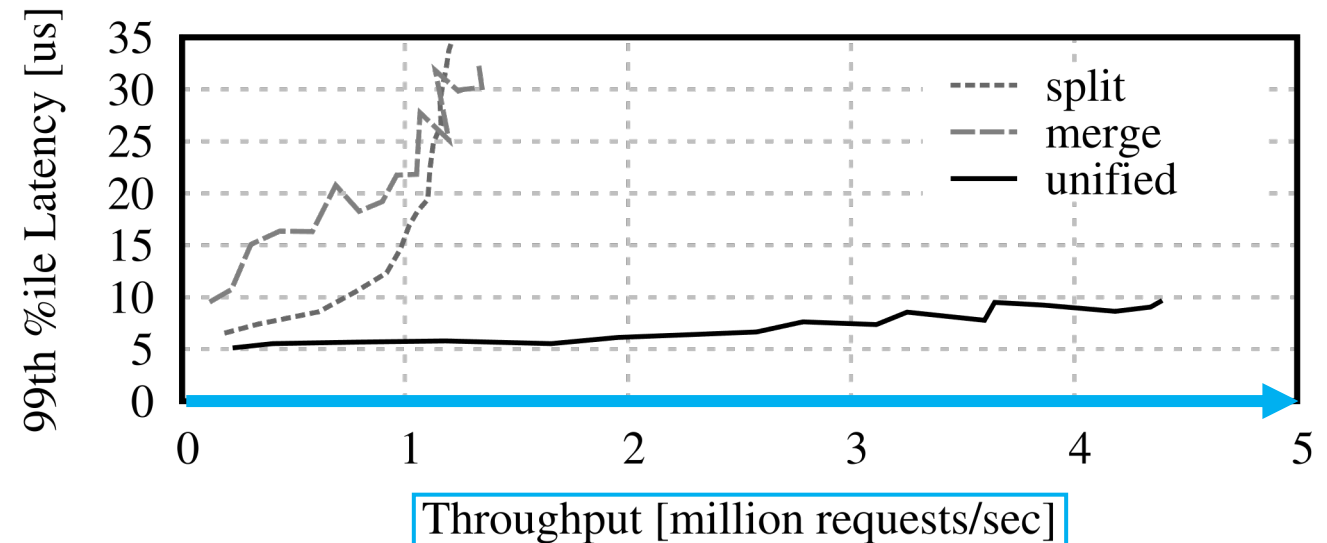
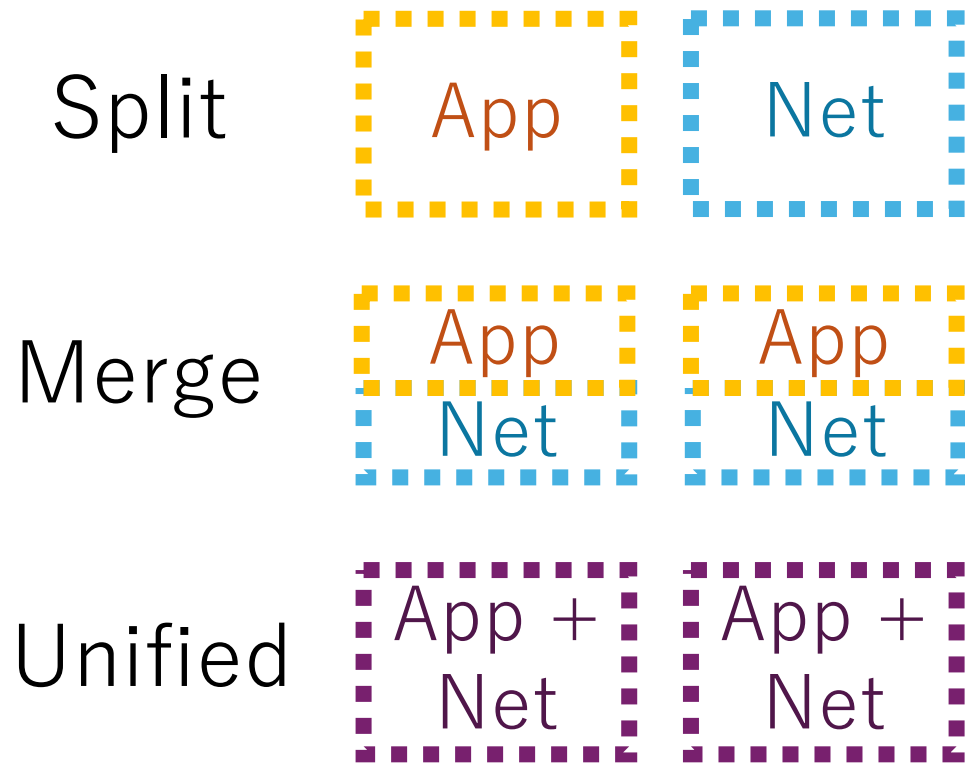
Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



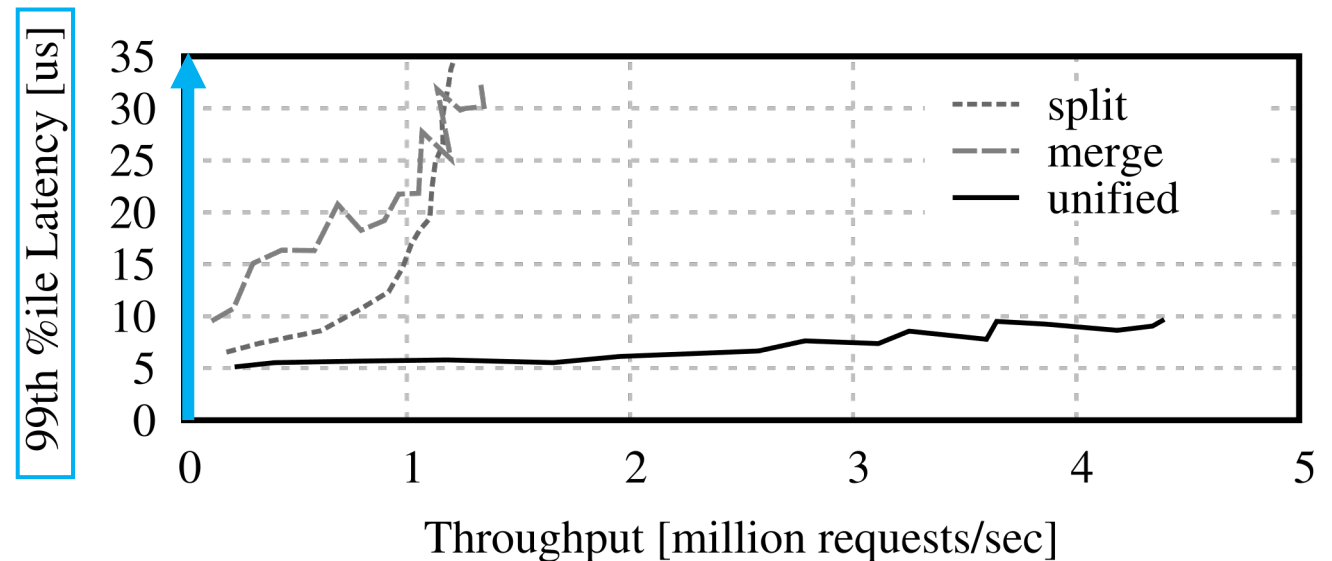
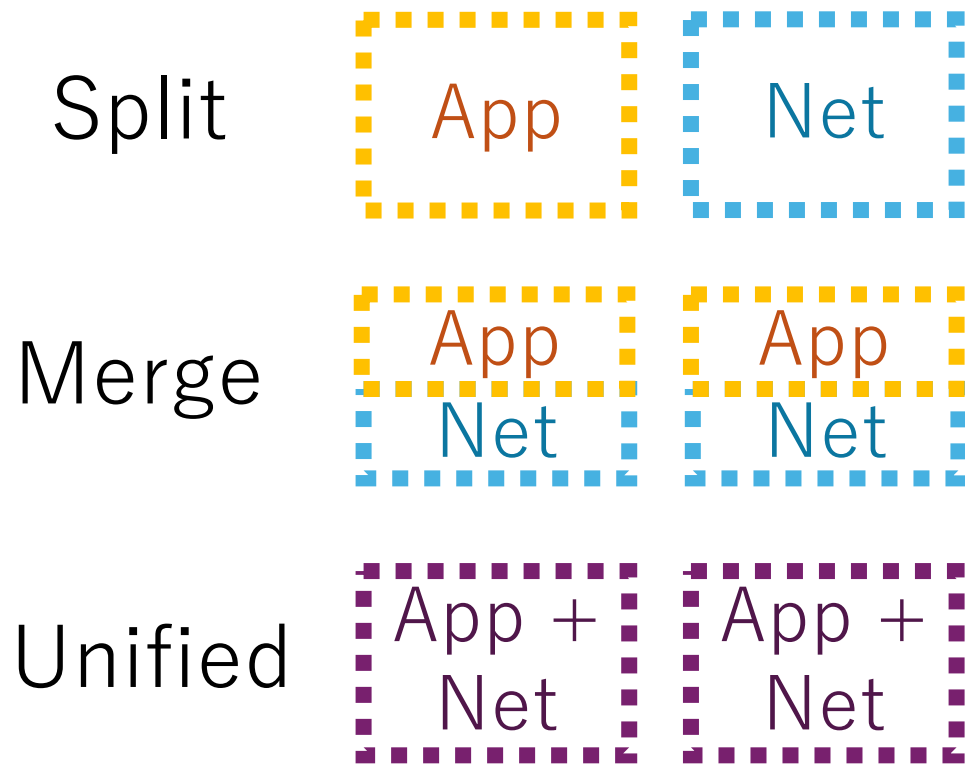
Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



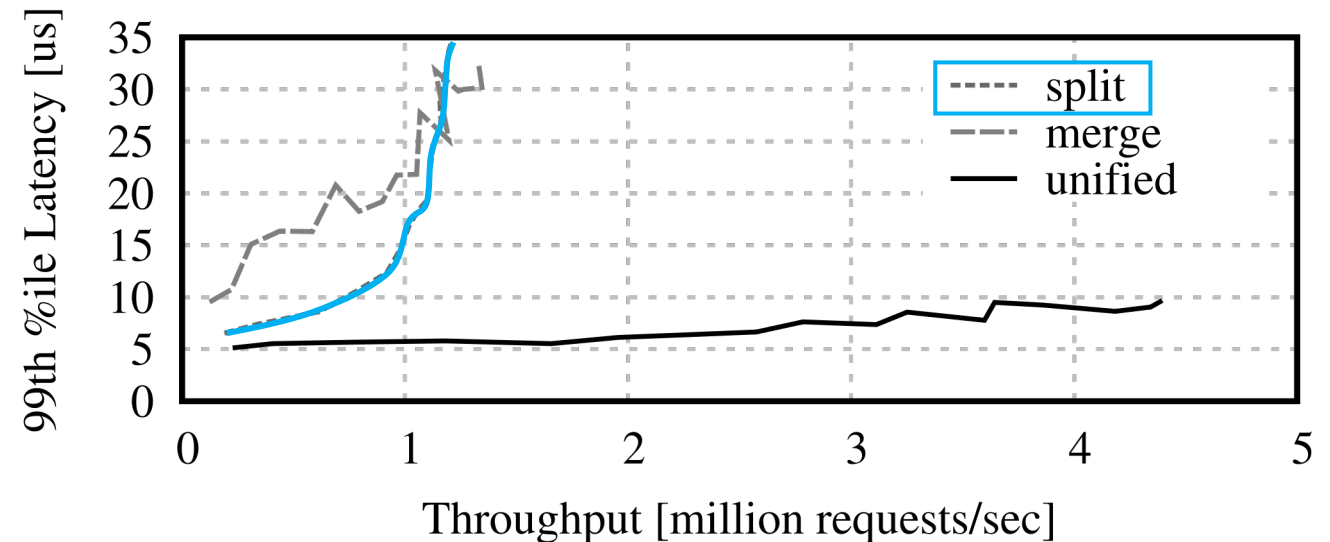
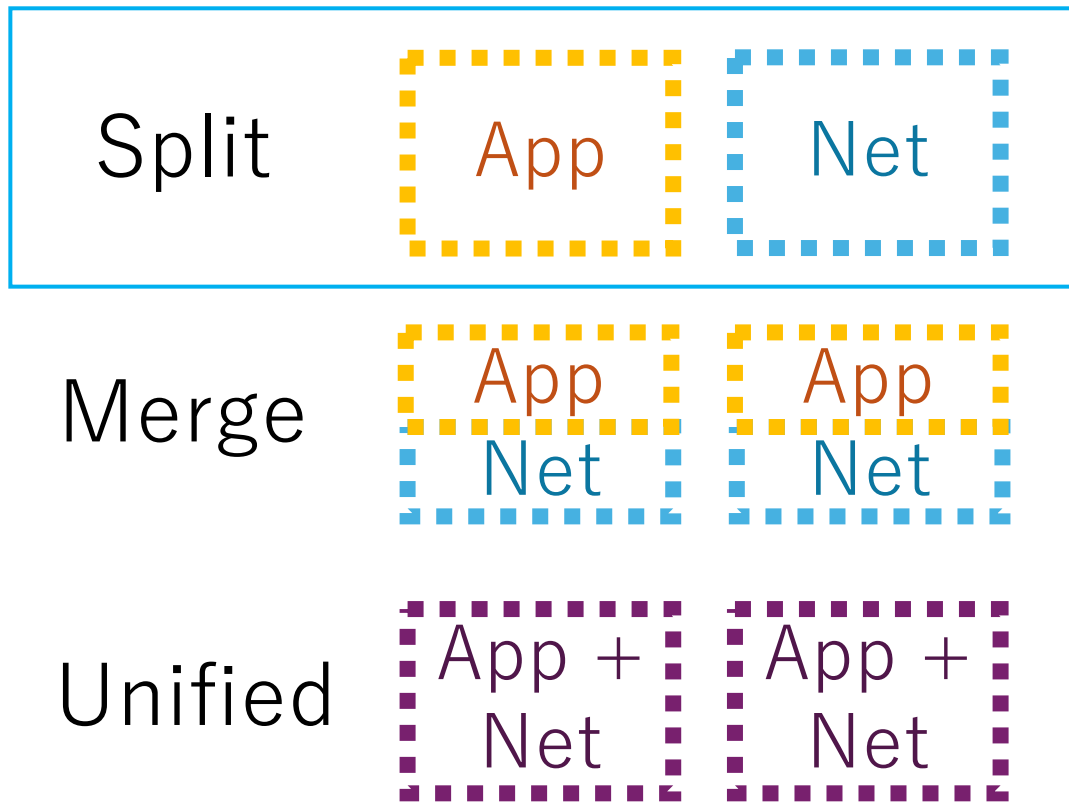
Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



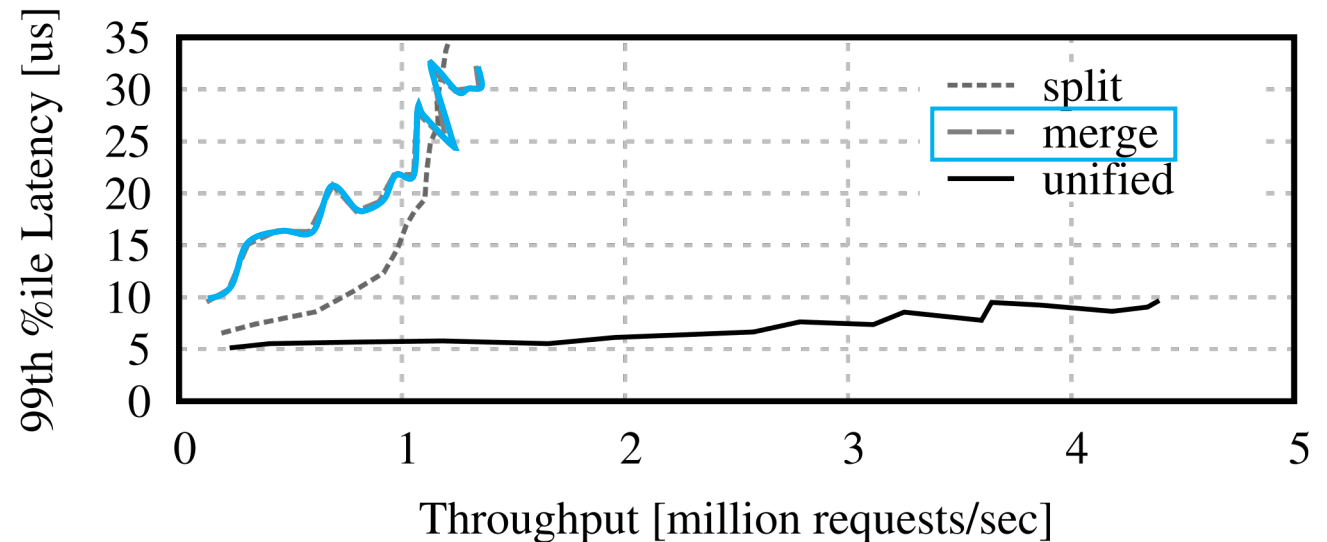
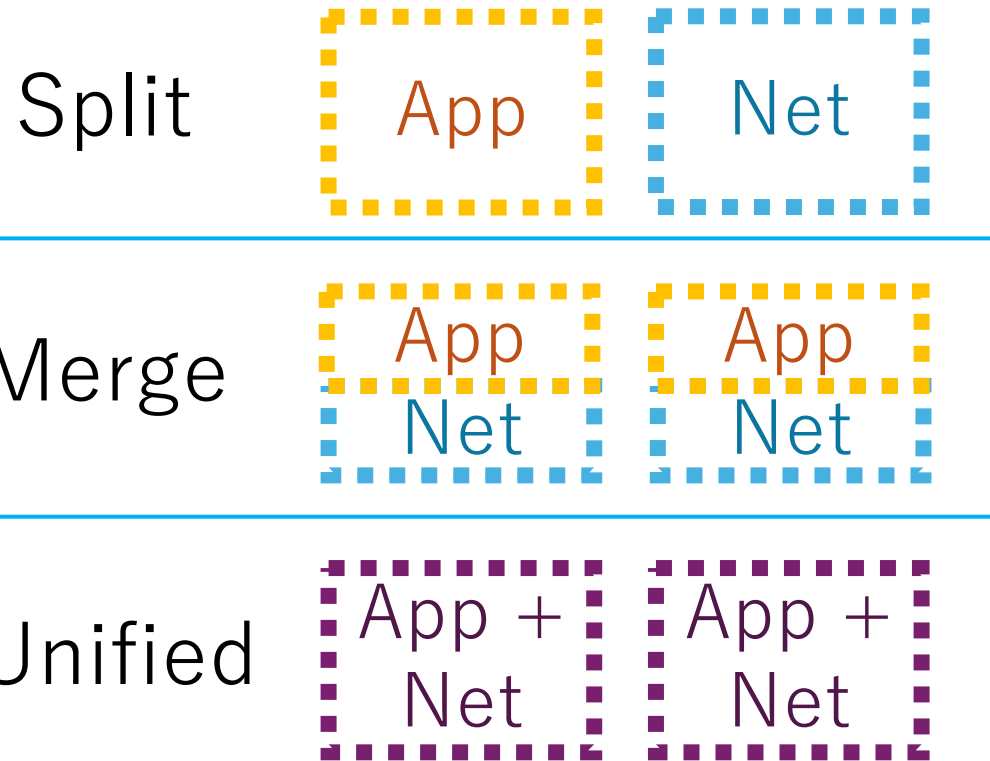
Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



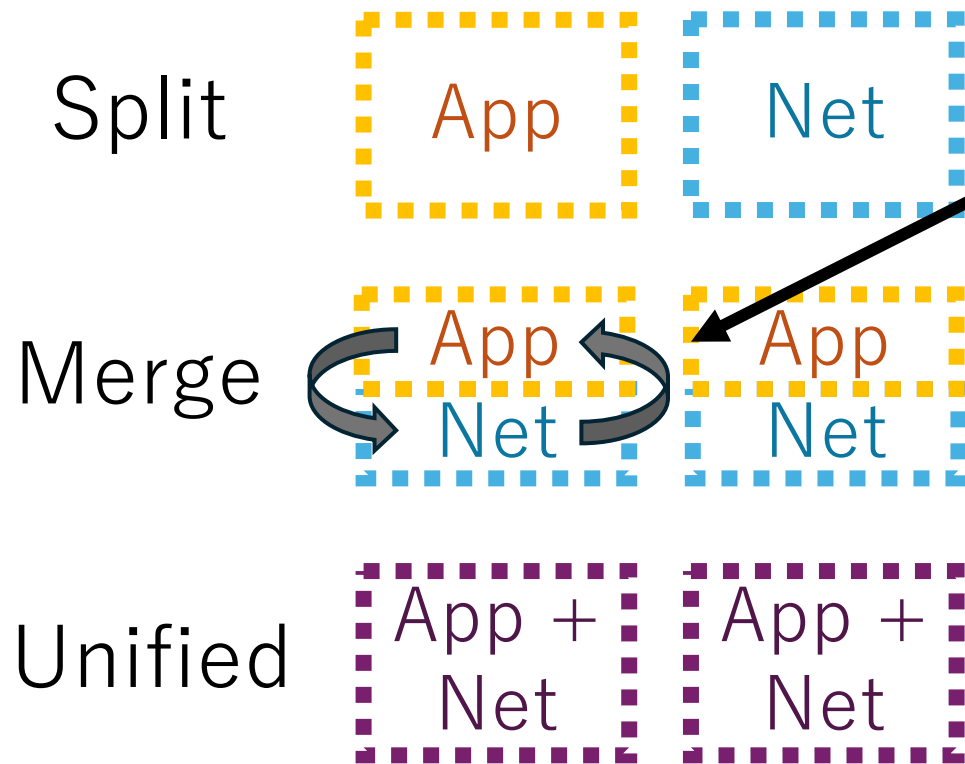
Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads

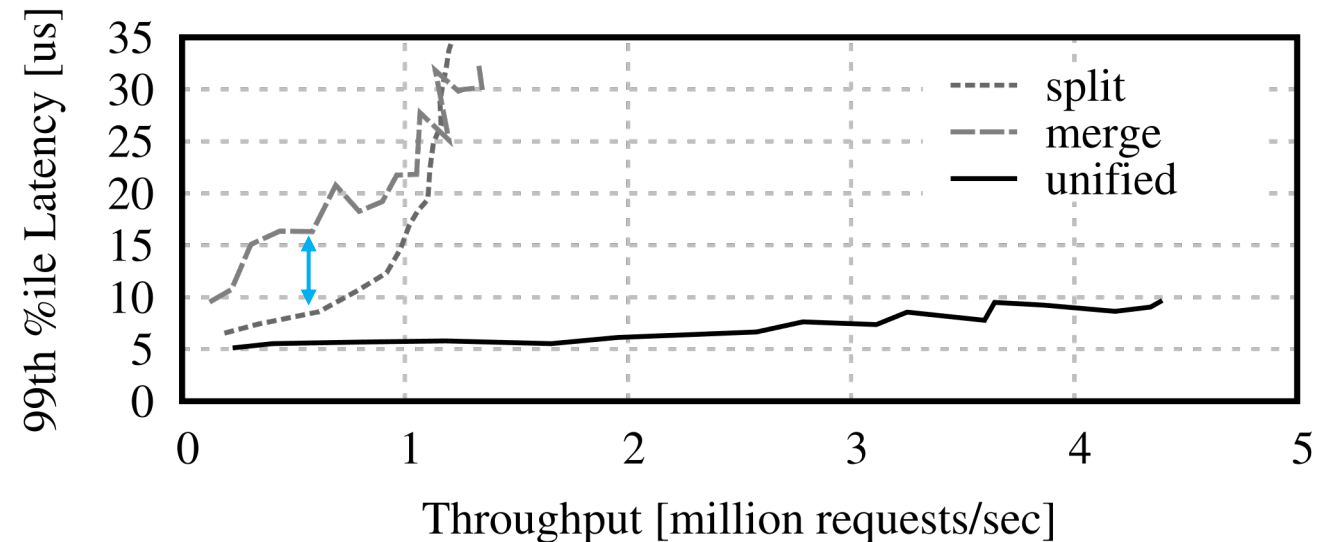


Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads

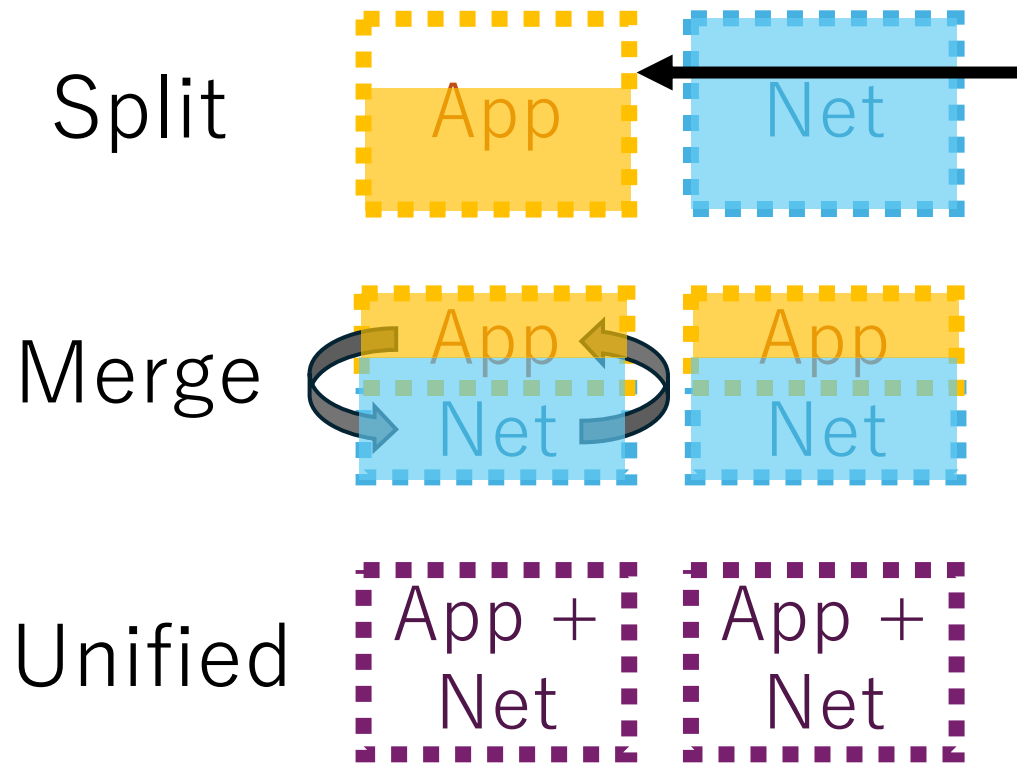


The latency gap comes from the app/net transition cost of the merge model

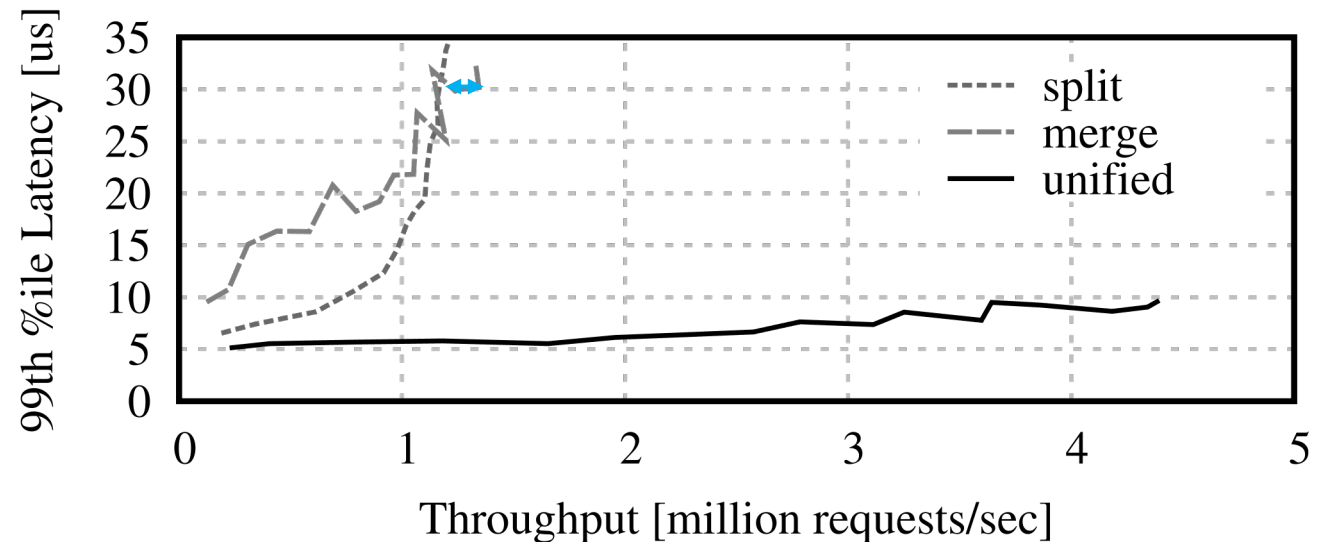


Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



The merge model exhibits higher throughput than the split model because of CPU utilization

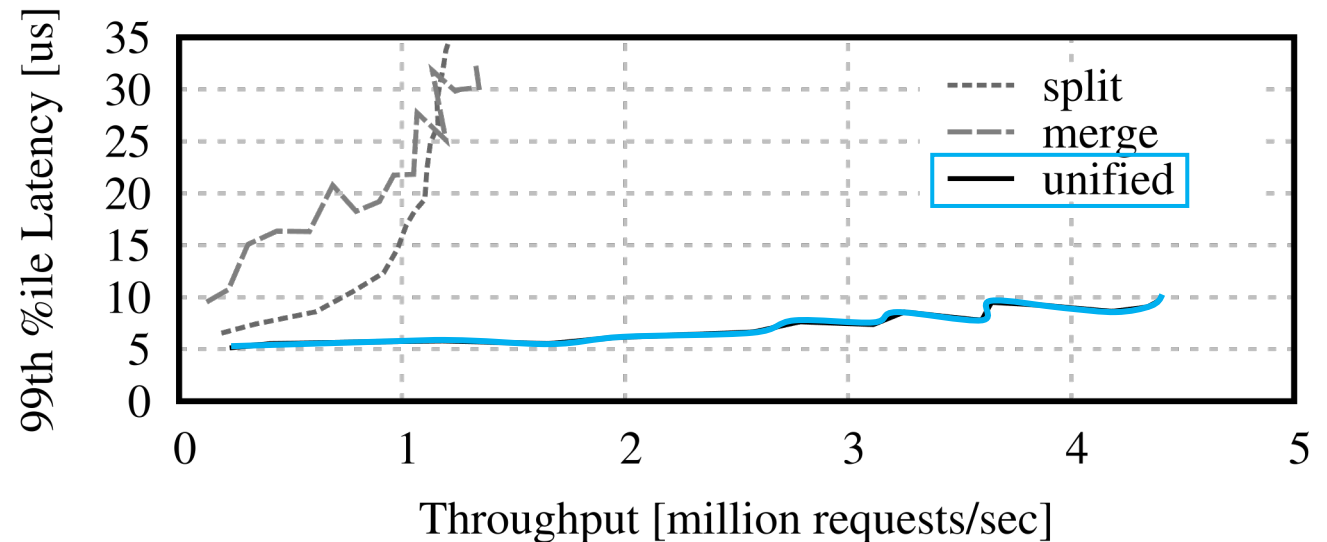


Evaluation: CPU Core Assignment Models

- The pinger and ponger apps exchange 1-byte TCP payloads



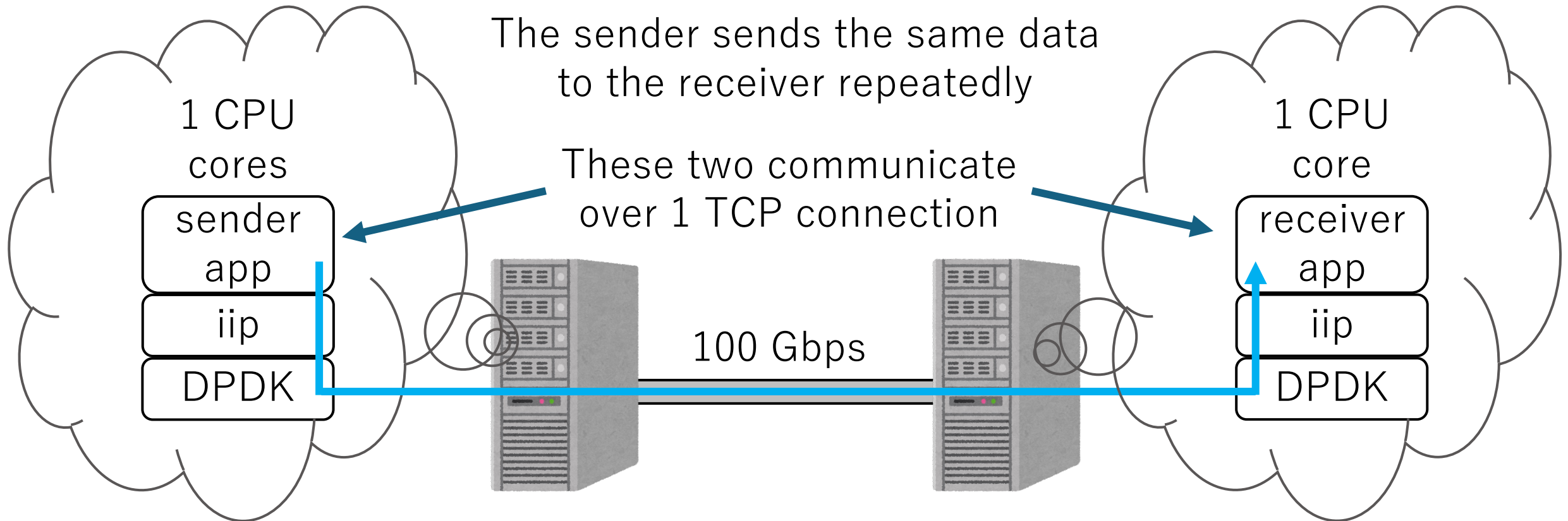
The unified model achieves the best speed because it is free from the issues of the others



Evaluation: Bulk Data Transfer

- Data transfer workload

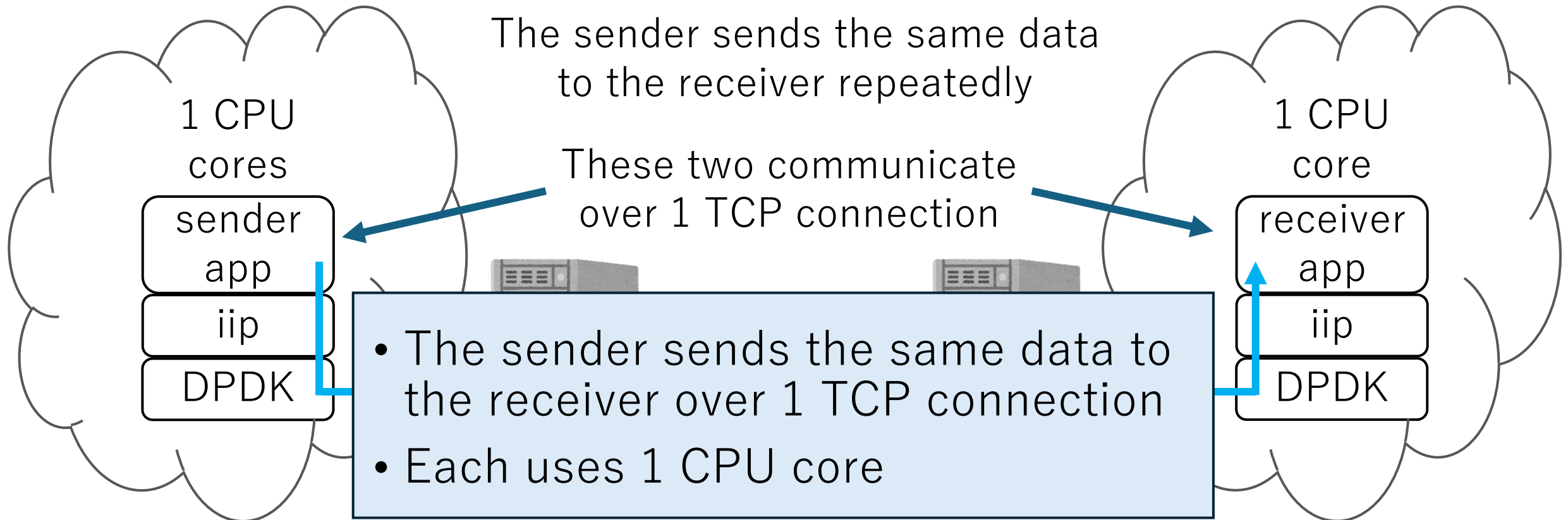
Each uses 1 CPU core



Evaluation: Bulk Data Transfer

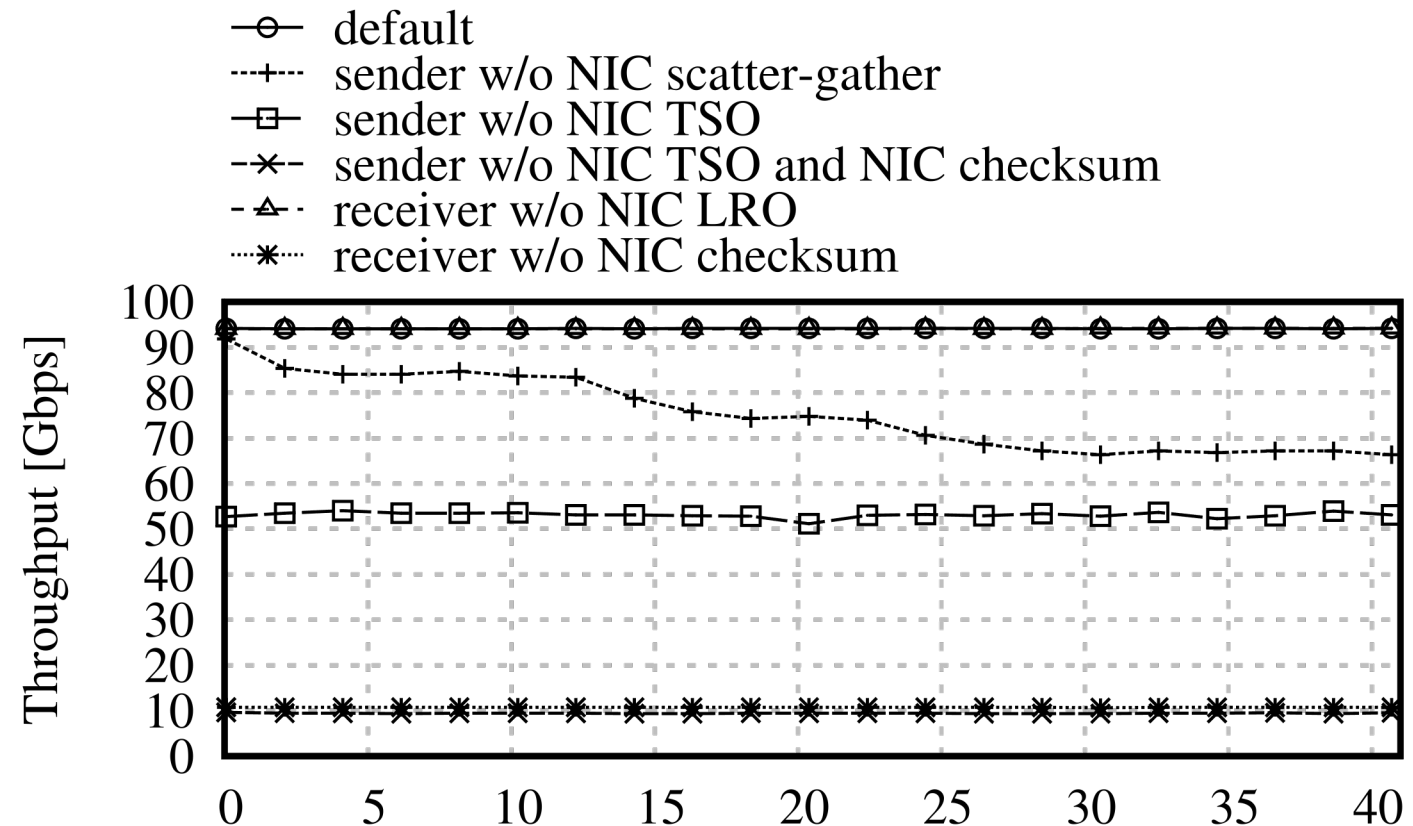
- Data transfer workload

Each uses 1 CPU core



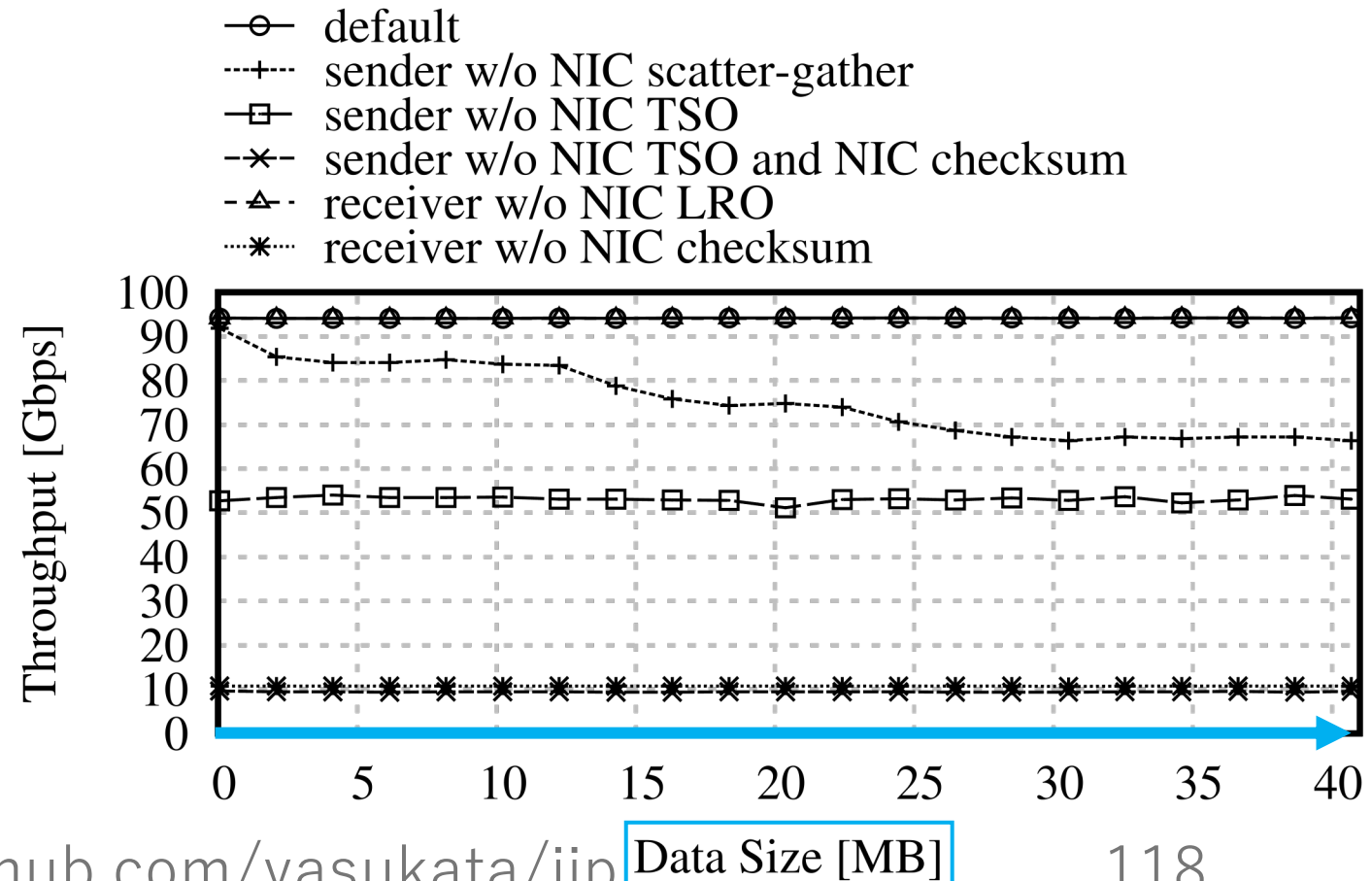
Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver



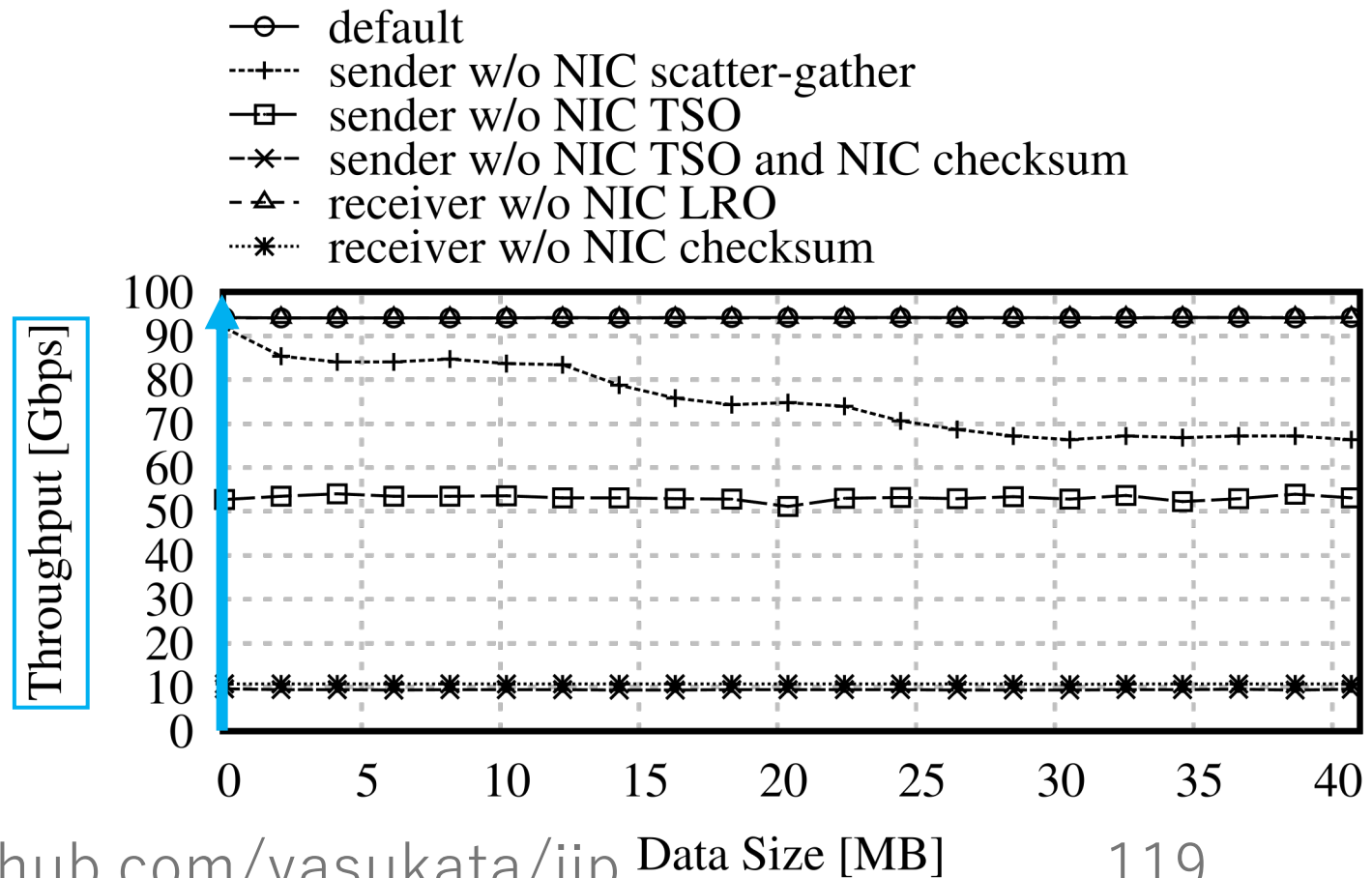
Evaluation: Bulk Data Transfer

- The sender repeatedly sends **the same data** to the receiver



Evaluation: Bulk Data Transfer

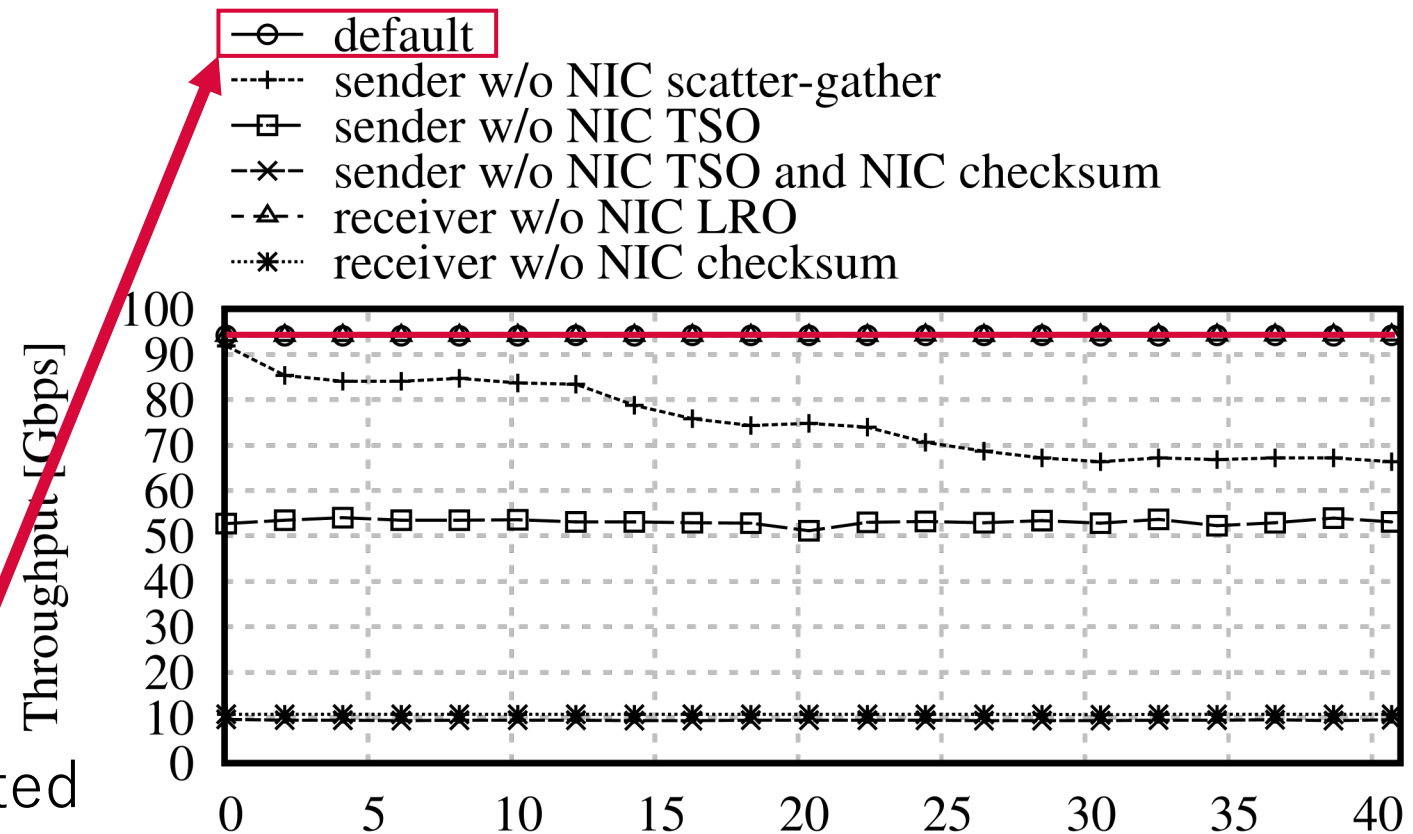
- The sender repeatedly sends the same data to the receiver



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver

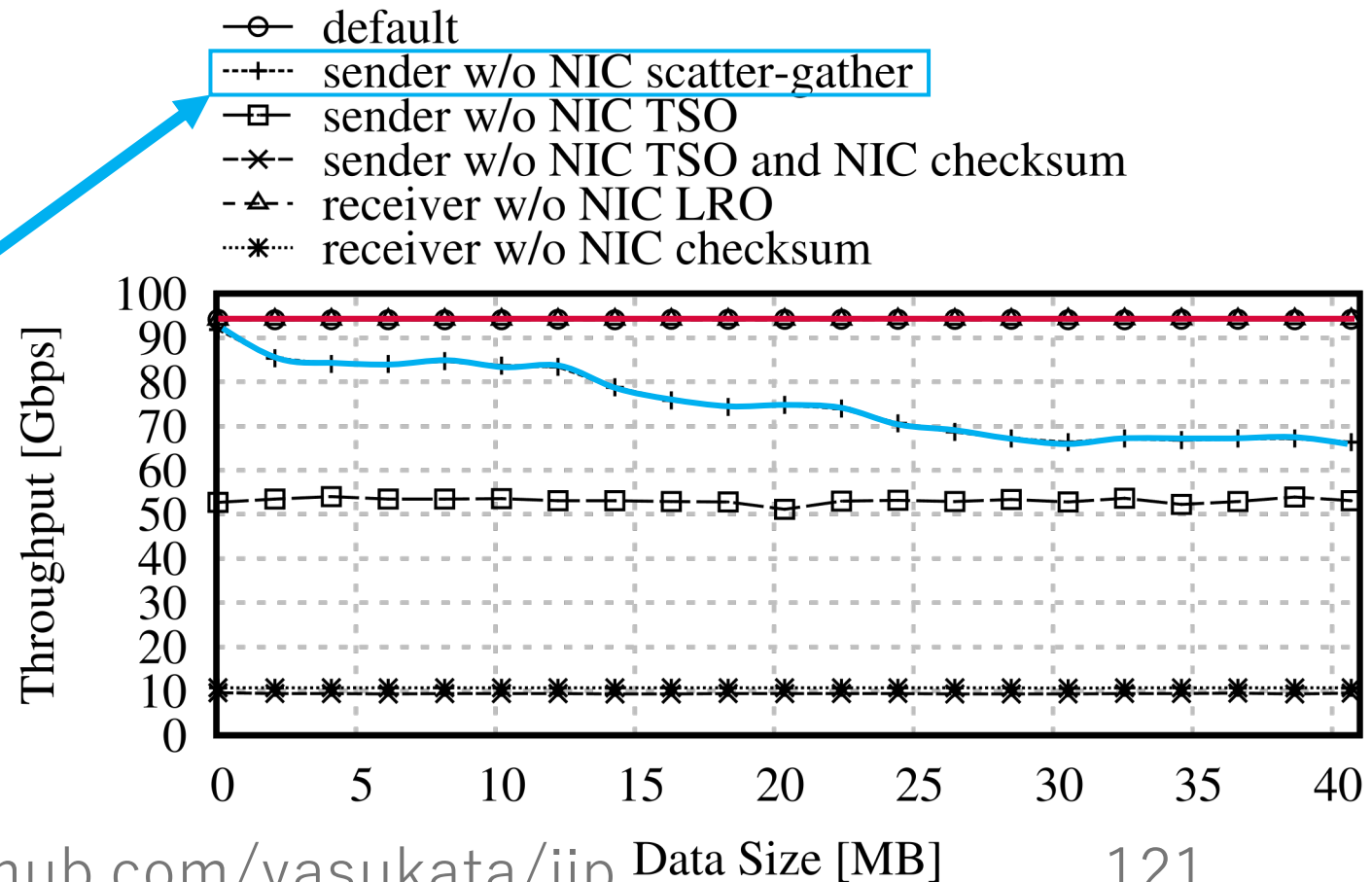
All NIC offloading features and zero-copy transmission are activated



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission

When the sender deactivates zero-copy transmission, throughput is degraded according to the size of the data

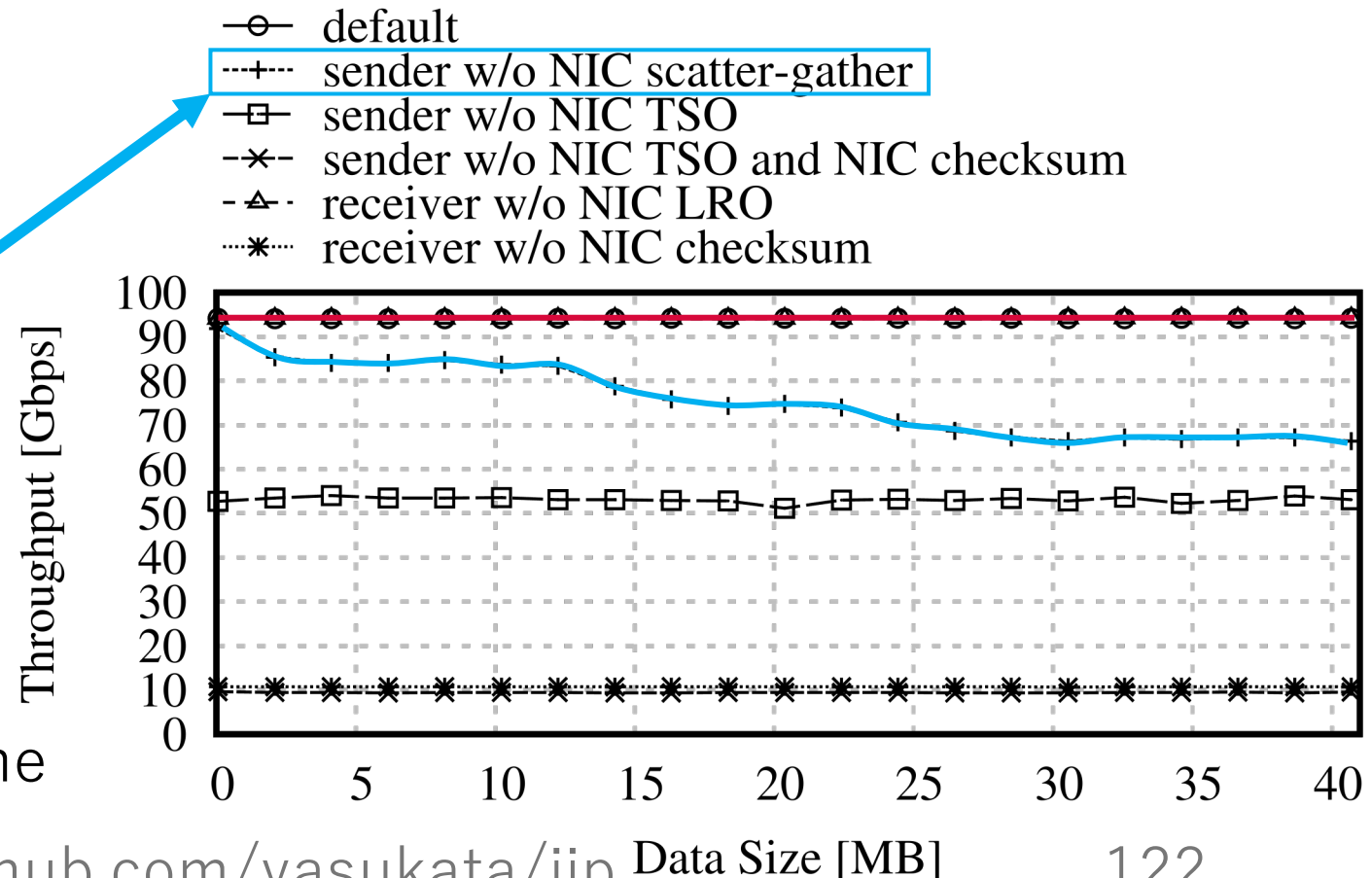


Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission

When the sender deactivates zero-copy transmission, throughput is degraded according to the size of the data

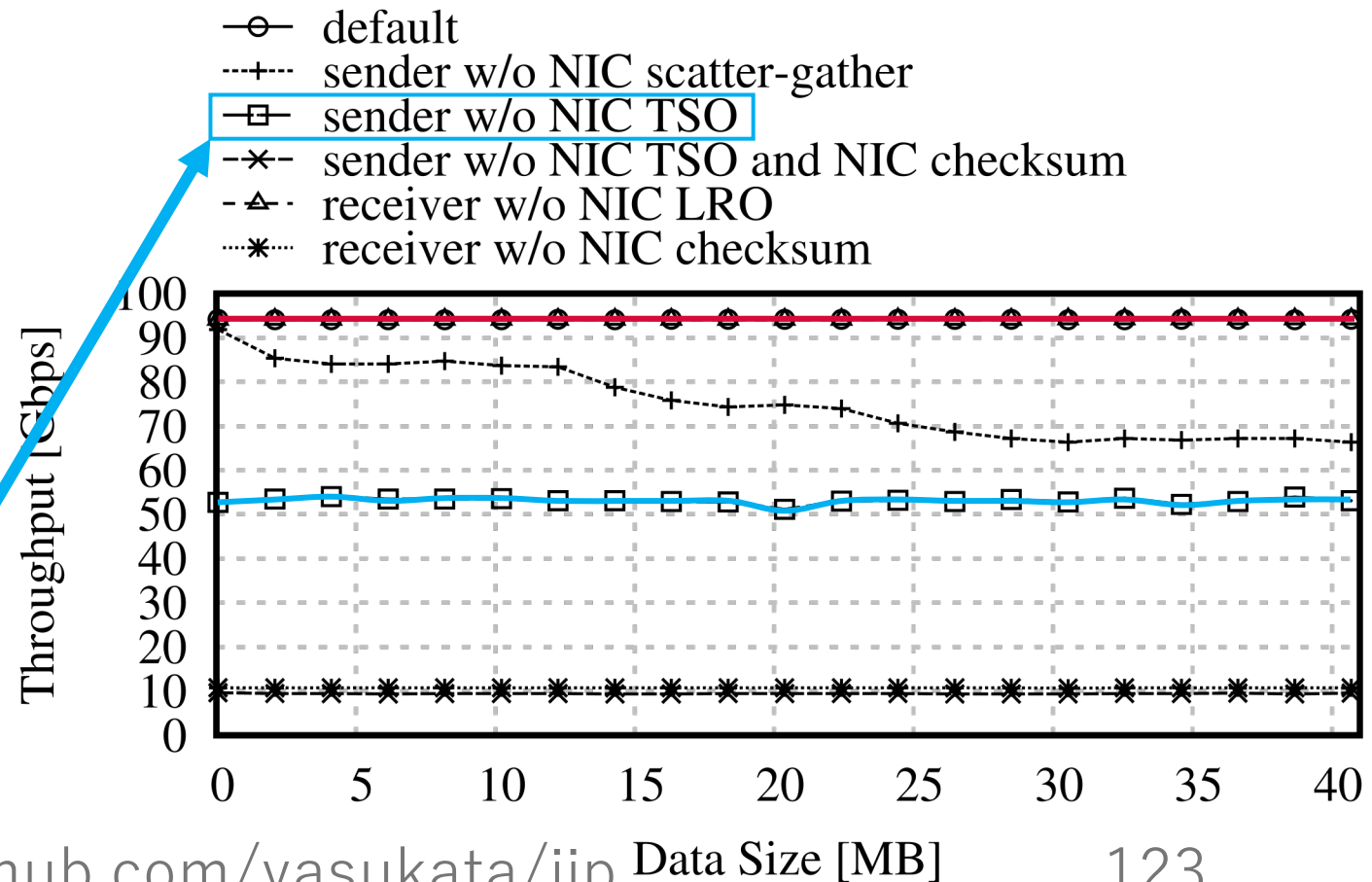
We consider this happens because the payload occupies the CPU cache



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission
 - TSO

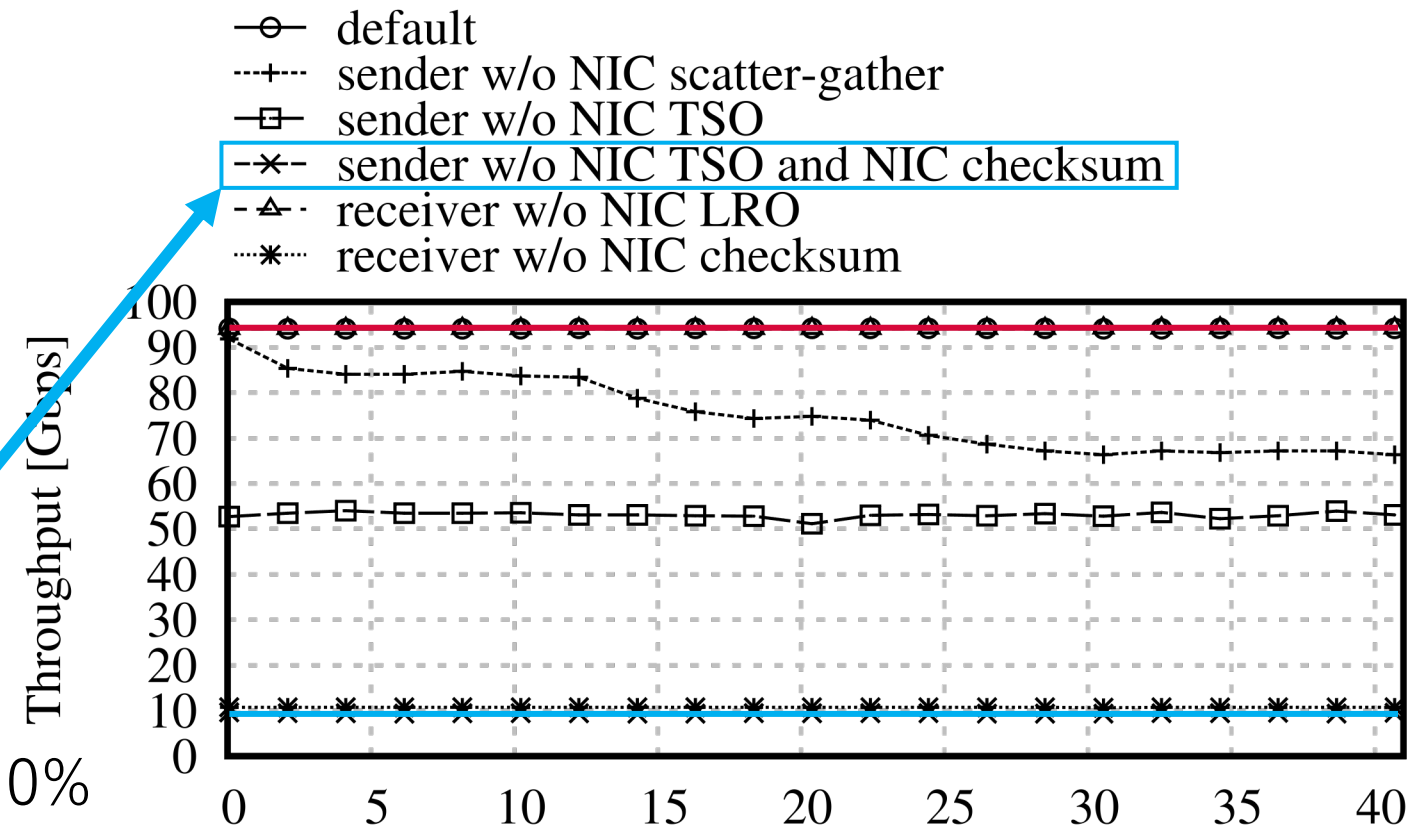
When the sender deactivates TSO, throughput is almost halved



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission
 - TSO
 - Checksum offloading

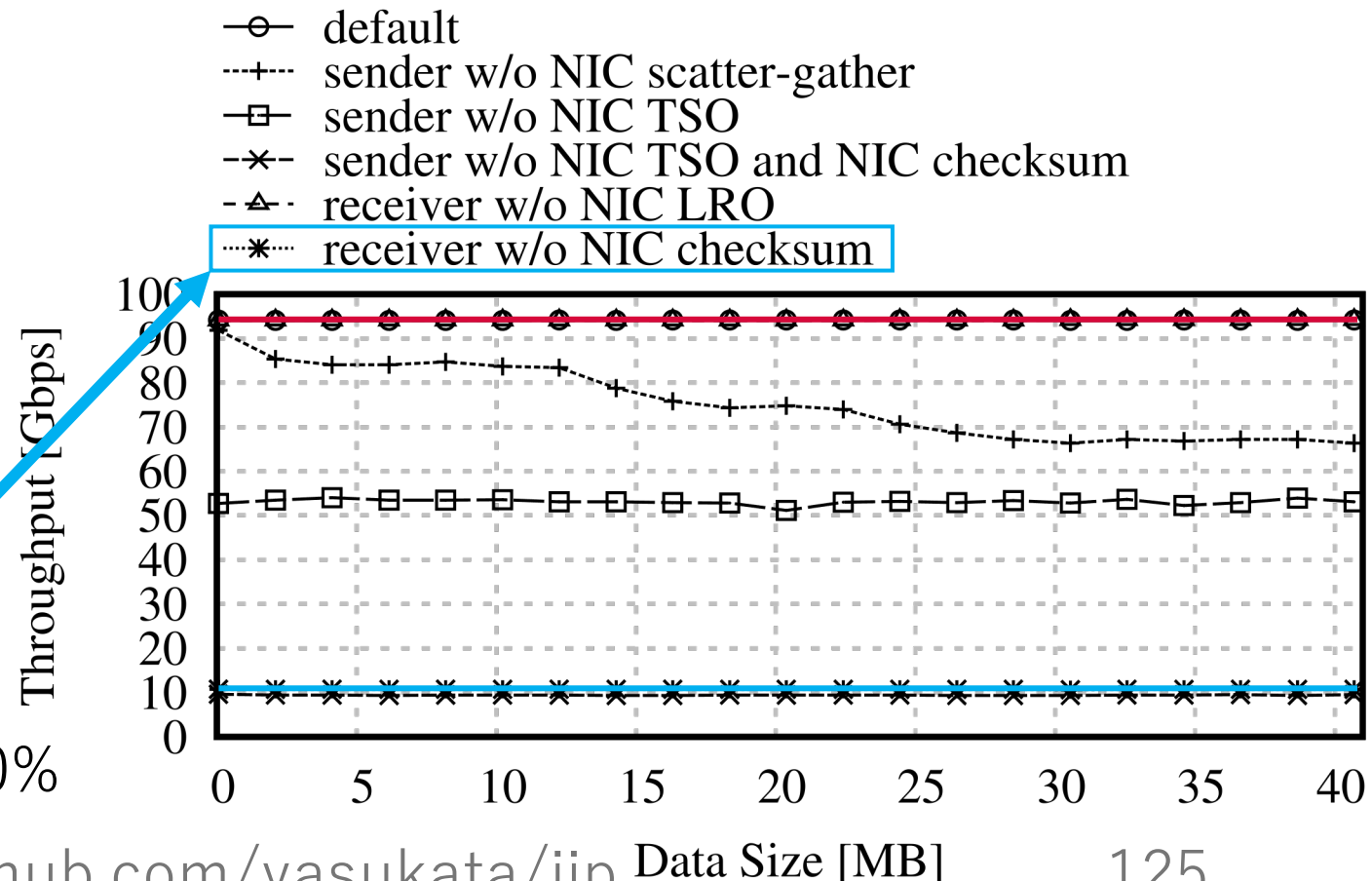
When the sender deactivates TSO and checksum offloading, throughput goes down to around 10%



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission
 - TSO
 - Checksum offloading

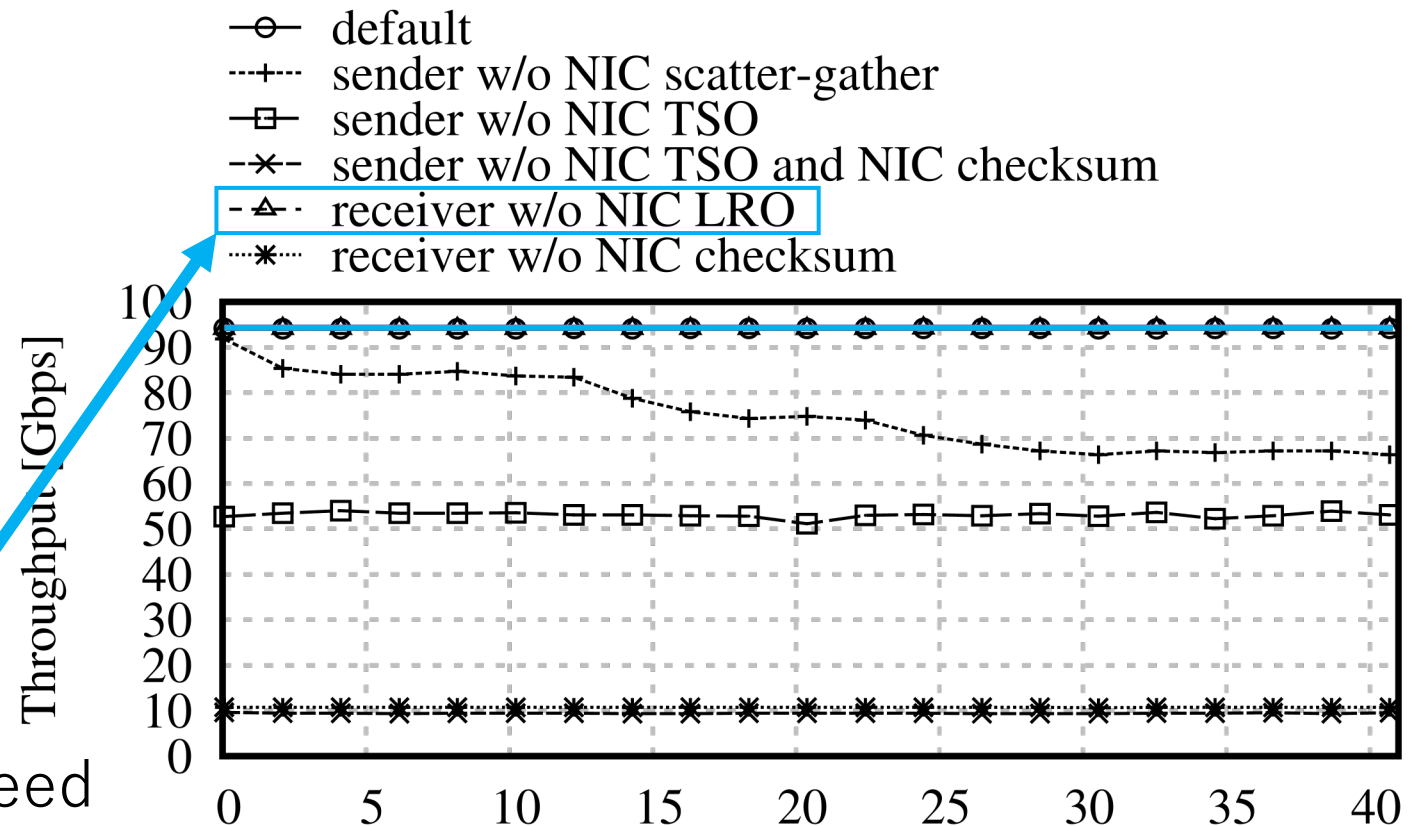
When the receiver deactivates checksum offloading, throughput goes down to around 10%



Evaluation: Bulk Data Transfer

- The sender repeatedly sends the same data to the receiver
- Performance factors
 - Zero-copy transmission
 - TSO
 - Checksum offloading

When the receiver disables LRO,
it can still catch up with the link speed



Evaluation: Bulk Data Transfer

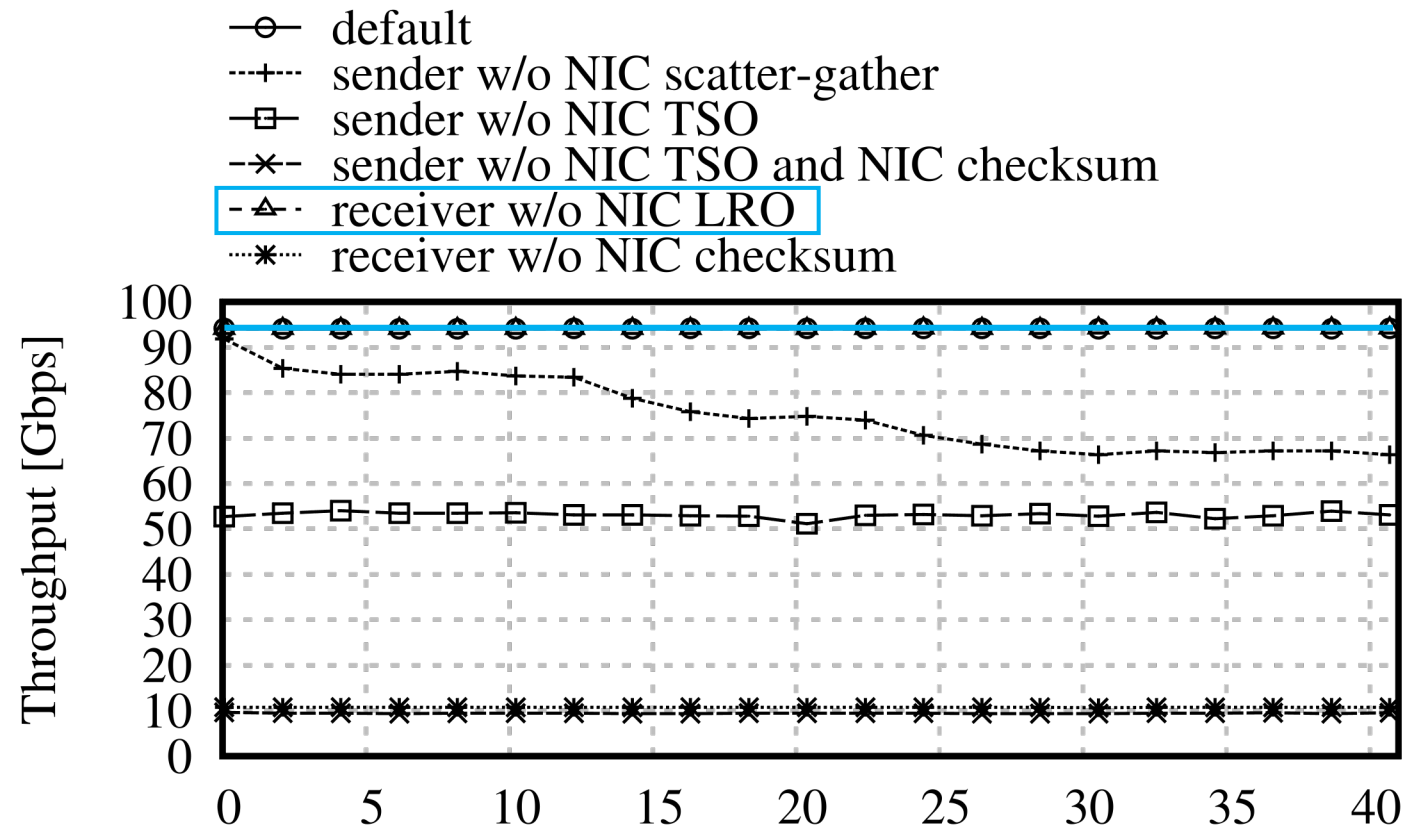
- The sender repeatedly sends the same data to the receiver

- Performance factors

- Zero-copy transmission
- TSO
- Checksum offloading

- Note

- This does not mean LRO is not necessary
 - Just we did not see differences in this workload



Summary

- iip is a TCP/IP stack implementation that aims to allow for easy integration and good performance simultaneously

Please try it if you are interested

- Main page: <https://github.com/yasukata/iip>
- Assets used in the paper: <https://github.com/yasukata/bench-iip/tree/9cf2488ec93ae51f4bd7b18923a5d1a233852f66>